Linear Algebra and its Applications, Spring 2013 Homework 14

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Note 1. This homework is due 8:30 am, December 24, 2013. Please submit a hard copy into the homework box outside the TAs' lab.

Note 2. "Problem sets" should be found in the textbook (the fourth edition).

- 1. (20 points) In this problem, we derive a necessary condition for local minima of constrained nonlinear programs. We consider a problem of minimizing f(x) over $x \in F \subseteq \mathbb{R}^n$.
 - (a) For a set $F \subseteq \mathbb{R}^n$, an interior point $x \in F$ is a point that does not lie on the boundary. More precisely, if there exists a small enough $\epsilon > 0$ such that $x + \lambda d \in F$ for all unitlength $d \in \mathbb{R}^n$ and $\lambda \in [0, \epsilon)$, then x is an interior point. Such a unit-length direction d is called a feasible direction at x. According to this definition, list three interior points of $S = \{x \in \mathbb{R}^3 | x_1^2 + x_2^2 + x_3^2 \leq 1\}.$
 - (b) For a set $F \subseteq \mathbb{R}^n$, a point $x \in F$ is a boundary point if it is not an interior point. For the S defined in Part (a), what are the boundary points?

Note. An extreme point is a boundary point but the converse is not true.

- (c) Show that if an interior point $\bar{x} \in F$ is a local minimum of $f(\cdot)$, we must have $\nabla f(\bar{x}) = 0$.
- (d) Show that if a boundary point $\bar{x} \in F$ is a local minimum of $f(\cdot)$, we must have $\nabla f(\bar{x})^T d \ge 0$ for all feasible direction d at \bar{x} .

Note. Intuitively, along any feasible direction, the function must "go up".

Note. The general necessary condition is: If a point $\bar{x} \in F$ is a local minimum of $f(\cdot)$, we must have $\nabla f(\bar{x})^T d \geq 0$ for all feasible direction d at \bar{x} . Why we do not need to assume that \bar{x} is a boundary point?

2. (20 points) Consider the nonlinear program

$$\min \quad -(p-c)(a-bp) \\ \text{s.t.} \quad p \ge 0$$

in which a seller chooses a price $p \in \mathbb{R}$ to maximize its profit. a > 0, b > 0, and $c \ge 0$ are parameters.

- (a) (5 points) Show that the objective function is convex over \mathbb{R} . Is it strictly convex?
- (b) (10 points) Given that the feasible region is convex, this is a convex program. Use this fact to solve the program by finding a global minimum p* as a function of a, b, and c.
 Note. Do you need to take boundary points into consideration?
- (c) (5 points) Let a = b = 1 and c = 0, depict the objective function.
- 3. (20 points) Consider the nonlinear program

$$\begin{array}{ll} \min & -(a-2d)^2 d\\ \text{s.t.} & 0 \le d \le \frac{a}{2} \end{array}$$

in which one looks for a length $x \in \mathbb{R}$ to maximize the volume of the box made by folding a paper. a > 0 is a parameter.

(a) (5 points) Show that the objective function is not convex over the feasible region. Note. It is NOT enough to show that the objective function is not convex over \mathbb{R} ! (b) (10 points) Find all the stationary points of the objective function. Are they all feasible? For each of them, determine whether it is a local minimum, local maximum, or a saddle point. Then determine the global minimum d^* .

Note. Do you need to take boundary points into consideration?

- (c) (5 points) Let a = 6, depict the objective function.
- 4. (20 points) Consider the nonlinear program

min
$$||x - x^{0}||$$

s.t. $||x - x^{i}|| \le d \quad \forall i = 1, ..., m.$ (1)

in which one chooses a location $x = (x_1, x_2) \in \mathbb{R}^2$ to locate an airport to minimize the distance to the capital at x_0 while ensuring the distance between the airport and each city is within d. $x^i = (x_1^i, x_2^i) \in \mathbb{R}^2$ for i = 0, ..., m and d > 0 are given parameters.

(a) (5 points) Show that the program is equivalent to

$$\min_{\substack{\|x - x^0\|^2 \\ \text{s.t.} \quad \|x - x^i\|^2 \le d^2 \quad \forall i = 1, ..., m. } }$$
(2)

- (b) (5 points) Show that the objective function of the program in (2) is convex over \mathbb{R}^2 .
- (c) (5 points) Show that the feasible region of the program in (2) is convex.
 - **Note.** Now we know that the program in (2) is a convex program. Though we are unable to analytically solve the program as we did for Problems 2 and 3, a greedy search algorithm can numerically solve it.
- (d) (5 points) Suppose the feasible region is nonempty. Will an optimal solution be a boundary point of the feasible region?

Note. It can be shown that the program in (1) is not a convex program. In general, a nonconvex program can be equivalent to a convex program and solving the latter is much easier.

- 5. (20 points) Solve the following problems:
 - (a) (5 points) Show that $\{(x_1, x_2) \in (0, \infty) \times \mathbb{R} | x_2 \leq \ln x_1\}$ is convex.
 - (b) (5 points) Find a set $F \subseteq \mathbb{R}^3$ over which $f_1(x_1, x_2, x_3) = x_1^3 + 2x_1x_3 + x_2x_3 x_3^2$ is convex.
 - (c) (10 points) For $f_2(x_1, x_2, x_3) = x_1^3 x_1x_2 + 2x_1x_3 + x_2x_3 x_3^2$, find all the stationary points. For each of them, determine whether it is a local minimum.