# Linear Algebra and its Applications, Spring 2013 Homework 14 

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Note 1. This homework is due 8:30 am, December 24, 2013. Please submit a hard copy into the homework box outside the TAs' lab.
Note 2. "Problem sets" should be found in the textbook (the fourth edition).

1. (20 points) In this problem, we derive a necessary condition for local minima of constrained nonlinear programs. We consider a problem of minimizing $f(x)$ over $x \in F \subseteq \mathbb{R}^{n}$.
(a) For a set $F \subseteq \mathbb{R}^{n}$, an interior point $x \in F$ is a point that does not lie on the boundary. More precisely, if there exists a small enough $\epsilon>0$ such that $x+\lambda d \in F$ for all unitlength $d \in \mathbb{R}^{n}$ and $\lambda \in[0, \epsilon)$, then $x$ is an interior point. Such a unit-length direction $d$ is called a feasible direction at $x$. According to this definition, list three interior points of $S=\left\{x \in \mathbb{R}^{3} \mid x_{1}^{2}+x_{2}^{2}+x_{3}^{2} \leq 1\right\}$.
(b) For a set $F \subseteq \mathbb{R}^{n}$, a point $x \in F$ is a boundary point if it is not an interior point. For the $S$ defined in Part (a), what are the boundary points?
Note. An extreme point is a boundary point but the converse is not true.
(c) Show that if an interior point $\bar{x} \in F$ is a local minimum of $f(\cdot)$, we must have $\nabla f(\bar{x})=0$.
(d) Show that if a boundary point $\bar{x} \in F$ is a local minimum of $f(\cdot)$, we must have $\nabla f(\bar{x})^{T} d \geq 0$ for all feasible direction $d$ at $\bar{x}$.
Note. Intuitively, along any feasible direction, the function must "go up".
Note. The general necessary condition is: If a point $\bar{x} \in F$ is a local minimum of $f(\cdot)$, we must have $\nabla f(\bar{x})^{T} d \geq 0$ for all feasible direction $d$ at $\bar{x}$. Why we do not need to assume that $\bar{x}$ is a boundary point?
2. (20 points) Consider the nonlinear program

$$
\begin{aligned}
\min & -(p-c)(a-b p) \\
\text { s.t. } & p \geq 0
\end{aligned}
$$

in which a seller chooses a price $p \in \mathbb{R}$ to maximize its profit. $a>0, b>0$, and $c \geq 0$ are parameters.
(a) (5 points) Show that the objective function is convex over $\mathbb{R}$. Is it strictly convex?
(b) (10 points) Given that the feasible region is convex, this is a convex program. Use this fact to solve the program by finding a global minimum $p^{*}$ as a function of $a, b$, and $c$.
Note. Do you need to take boundary points into consideration?
(c) (5 points) Let $a=b=1$ and $c=0$, depict the objective function.
3. (20 points) Consider the nonlinear program

$$
\begin{array}{ll}
\min & -(a-2 d)^{2} d \\
\text { s.t. } & 0 \leq d \leq \frac{a}{2}
\end{array}
$$

in which one looks for a length $x \in \mathbb{R}$ to maximize the volume of the box made by folding a paper. $a>0$ is a parameter.
(a) (5 points) Show that the objective function is not convex over the feasible region.

Note. It is NOT enough to show that the objective function is not convex over $\mathbb{R}$ !
(b) (10 points) Find all the stationary points of the objective function. Are they all feasible? For each of them, determine whether it is a local minimum, local maximum, or a saddle point. Then determine the global minimum $d^{*}$.
Note. Do you need to take boundary points into consideration?
(c) (5 points) Let $a=6$, depict the objective function.
4. (20 points) Consider the nonlinear program

$$
\begin{align*}
\min & \left\|x-x^{0}\right\| \\
\text { s.t. } & \left\|x-x^{i}\right\| \leq d \quad \forall i=1, \ldots, m \tag{1}
\end{align*}
$$

in which one chooses a location $x=\left(x_{1}, x_{2}\right) \in \mathbb{R}^{2}$ to locate an airport to minimize the distance to the capital at $x_{0}$ while ensuring the distance between the airport and each city is within $d$. $x^{i}=\left(x_{1}^{i}, x_{2}^{i}\right) \in \mathbb{R}^{2}$ for $i=0, \ldots, m$ and $d>0$ are given parameters.
(a) (5 points) Show that the program is equivalent to

$$
\begin{align*}
\min & \left\|x-x^{0}\right\|^{2} \\
\text { s.t. } & \left\|x-x^{i}\right\|^{2} \leq d^{2} \quad \forall i=1, \ldots, m \tag{2}
\end{align*}
$$

(b) (5 points) Show that the objective function of the program in (2) is convex over $\mathbb{R}^{2}$.
(c) (5 points) Show that the feasible region of the program in (2) is convex.

Note. Now we know that the program in (2) is a convex program. Though we are unable to analytically solve the program as we did for Problems 2 and 3, a greedy search algorithm can numerically solve it.
(d) (5 points) Suppose the feasible region is nonempty. Will an optimal solution be a boundary point of the feasible region?

Note. It can be shown that the program in (1) is not a convex program. In general, a nonconvex program can be equivalent to a convex program and solving the latter is much easier.
5. (20 points) Solve the following problems:
(a) (5 points) Show that $\left\{\left(x_{1}, x_{2}\right) \in(0, \infty) \times \mathbb{R} \mid x_{2} \leq \ln x_{1}\right\}$ is convex.
(b) (5 points) Find a set $F \subseteq \mathbb{R}^{3}$ over which $f_{1}\left(x_{1}, x_{2}, x_{3}\right)=x_{1}^{3}+2 x_{1} x_{3}+x_{2} x_{3}-x_{3}^{2}$ is convex.
(c) (10 points) For $f_{2}\left(x_{1}, x_{2}, x_{3}\right)=x_{1}^{3}-x_{1} x_{2}+2 x_{1} x_{3}+x_{2} x_{3}-x_{3}^{2}$, find all the stationary points. For each of them, determine whether it is a local minimum.

