# IM2010: Operations Research Linear Programming Formulation (Chapter 3)

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## Introduction

- It is important to learn how to model a practical situation as a linear program.
- This process is typically called linear programming formulation or modeling.
- ▶ We will introduce three types of LP problems, demonstrate how to formulate them, and discuss some important issues.
  - There are certainly many other types of LP problems.
- ► For large-scale problems, **compact formulations** are used.

# Road map

- ▶ Resource allocation.
- Materials blending.
- ▶ Production and inventory.
- Compact formulations.

## **Resource allocation**

- We produce products to sell.
- ► Each product requires some resources. **Resources are limited**.
- ▶ We want to maximize the total sales revenue while ensuring resources are enough.

# Resource allocation: the problem

- We may produce desks and tables.
  - Producing a desk requires four units of wood, one hour of labor, and 30 minutes of machine time.
  - Producing a table requires five units of wood, two hours of labor, and 20 minutes of machine time.
- We may sell everything we produce.
- ▶ For each day, we have
  - ▶ Two workers, each works for eight hours.
  - One machine that can run for eight hours.
  - A supply of 36 units of wood.

▶ Desks and tables are sold at \$800 and \$600 per unit, respectively.

- When we define decision variables, try to answer "what are the decisions to make?"
- ▶ In this example, the decision we want to make is the production quantities of desks and tables.
- ▶ Therefore, we define our decision variables as follows:
- ► Let

 $x_1$  = number of desks produced in a day and  $x_2$  = number of tables produced in a day.

# Formulation: objective function

- ► In the objective function, we write down the quantity that we want to **maximize** or **minimize**.
- ▶ In this example, we want to maximize the total sales revenue.
  - Given our decision variables, may we determine the sales revenue?
  - The sales revenue is  $800x_1 + 600x_2$ .
- ▶ The objective function is thus

max  $800x_1 + 600x_2$ .

## Formulation: constraints

- ► For each **restriction** or **limitation**, we write a constraint.
- ▶ Summarizing data into a table typically helps:

Resource	Consum	Total supply	
nesource	Desk	Table	Total supply
Wood	4 units	5 units	36 units
Labor hour	1 hour	2 hours	16 hours
Machine time	30 minutes	20 minutes	8 hours

## Formulation: constraints

▶ The supply of wood is limited:

 $4x_1 + 5x_2 \le 36.$ 

• The number of labor hours is limited:

 $x_1 + 2x_2 \le 16.$ 

▶ The amount of machine time is limited:

 $30x_1 + 20x_2 \le 240.$ 

▶ Use the same unit of measurement!

▶ Production quantities are **nonnegative**:  $x_i \ge 0$   $\forall i = 1, 2$ .

## Formulation: the complete formulation

#### ▶ The complete formulation is

$\max$	$800x_{1}$	+	$600x_2$			
s.t.	$4x_1$	+	$5x_2$	$\leq$	36	(wood)
	$x_1$	+	$2x_2$	$\leq$	16	(labor)
	$30x_{1}$	+	$20x_{2}$	$\leq$	240	(machine)
	$x_i \ge 0$	d C	i = 1, 2			

- ▶ **Clearly** define decision variables **in front of** your formulation.
- ▶ Write **comments** after the objective function and constraints.
- ▶ Do not forget nonnegativity constraints.

### Formulation: the complete formulation

#### • We may simplify the formulation:

• Once we find an optimal solution, please use the **original** objective function in calculating the associated objective value.

# Fractional and integer variables

- ▶ The optimal solution of this linear program is to produce 6.86 desks or 1.71 tables. Can we?
- ▶ Indeed we cannot. Then why linear programming?
  - ▶ It always **supports** our decisions. E.g., we may round down to get a feasible solution that is near optimal.
  - ▶ In practice, people use mathematical programming typically when the quantities are **large**. Rounding 6.86 may deviate a lot but rounding 68600.86 may be much more acceptable.
  - When it is necessary, we should impose integer constraints on variables and apply **integer programming** (to be covered later in the semester).
- ▶ If it is not specified in the problem, using LP is enough.

# Road map

- ▶ Resource allocation.
- ► Materials blending.
- ▶ Production and inventory.
- Compact formulations.

# Material blending

- ► In some situations, we need to determine not only products to produce but also materials to input.
- ▶ This is because we have some **flexibility** in making the products.
- ▶ For example, in making orange juice, we may use orange, sugar, water, etc. Different ways of **blending** these materials results in different qualities of juice.
- ► The goal is to save money (lower the proportion of expensive materials) while maintaining quality.
- ▶ This is introduced in Section 3.7 of the textbook.

# Material blending: the problem

- ▶ We blend materials 1, 2, and 3 to make products 1 and 2.
- ▶ The quality of a product, which depends on the proportions of these three materials, must meet the standard:
  - ▶ Product 1: at least 40% of material 1; at least 20% of material 2.
  - ▶ Product 2: at least 50% of material 1; at most 30% of material 3.
- At most 100 kg of product 1 and 150 kg of product 2 can be sold.
- ▶ Prices for products 1 and 2 are \$10 and \$15 per kg, respectively.
- ▶ Costs for materials 1 to 3 are \$8, \$4, and \$3 per kg, respectively.
- ▶ Amount of a product made equals the amount of materials input.
- We want to maximize the total profit.

▶ Probably our first attempt is to define the following: Let

 $x_1 = \text{kg of product 1 produced},$   $x_2 = \text{kg of product 2 produced},$   $y_1 = \text{kg of material 1 produced},$   $y_2 = \text{kg of material 2 produced, and}$  $y_3 = \text{kg of material 3 produced}.$ 

- ▶ May we express the quality of each product? No!
- ▶ We need to specify the amount of material 1 used for product 1, the amount of material 1 used for product 2, etc.
- ▶ So we need to **redefine** our decision variables.

▶ How about this: Let

 $x_1 = \text{kg}$  of material 1 used for product 1,  $x_2 = \text{kg}$  of material 1 used for product 2,  $x_3 = \text{kg}$  of material 2 used for product 1,  $x_4 = \text{kg}$  of material 2 used for product 2,  $x_5 = \text{kg}$  of material 3 used for product 1, and  $x_6 = \text{kg}$  of material 3 used for product 2.

▶ The definition is correct and precise, but **not easy to use**.

Similar to computer programming: give your variables reasonable names that allow people to know what they are.

▶ How about this: Let

 $x_{11} = \text{kg of material 1 used for product 1},$   $x_{12} = \text{kg of material 1 used for product 2},$   $x_{21} = \text{kg of material 2 used for product 1},$   $x_{22} = \text{kg of material 2 used for product 2},$   $x_{31} = \text{kg of material 3 used for product 1, and}$  $x_{32} = \text{kg of material 3 used for product 2}.$ 

▶ Much better.

- ▶ How to find the production quantities of products and the purchasing quantities of materials?
- Let's summarize the variables into a table:

	Product 1	Product 2
Material 1	$x_{11}$	$x_{12}$
Material 2 Material 3	$egin{array}{c} x_{21} \ x_{31} \end{array}$	$egin{array}{c} x_{22} \ x_{32} \end{array}$

▶ What are the desired quantities?

#### ▶ The desired quantities:

	Product 1	Product 2	Purchasing quantity
Material 1	$x_{11}$	$x_{12}$	$x_{11} + x_{12}$
Material 2	$x_{21}$	$x_{22}$	$x_{21} + x_{22}$
Material 3	$x_{31}$	$x_{32}$	$x_{31} + x_{32}$
Production quantity	$x_{11} + x_{21} + x_{31}$	$x_{12} + x_{22} + x_{32}$	

## Formulation: objective function

- ▶ Let's write down the total profit.
- ▶ Sales revenues depend on the amount of products we sell.
  - How much product 1 may we sell?  $x_{11} + x_{21} + x_{31}$ .
  - Similarly, we have  $x_{12} + x_{22} + x_{32}$  kg of product 2.
- ▶ Material costs depend on the amount of materials we purchase.
  - ► Similarly, we need to buy x<sub>11</sub> + x<sub>12</sub> kg of material 1, x<sub>21</sub> + x<sub>22</sub> kg of material 2 and x<sub>31</sub> + x<sub>32</sub> kg of material 3.
- ▶ The objective function is

$$\max 10(x_{11} + x_{21} + x_{31}) + 15(x_{12} + x_{22} + x_{32}) - 8(x_{11} + x_{12}) - 4(x_{21} + x_{22}) - 3(x_{31} + x_{32}) = \max 2x_{11} + 7x_{12} + 6x_{21} + 11x_{22} + 7x_{31} + x_{32}.$$

### Formulation: quality constraints

▶ In product 1, how to guarantee at least 40% are material 1?

$$\frac{x_{11}}{x_{11} + x_{21} + x_{31}} \ge 0.4.$$

- ▶ It is conceptually correct. However, it is **nonlinear**!
- ▶ Let's fix the nonlinearity by taking the denominator to the RHS:

$$x_{11} \ge 0.4(x_{11} + x_{21} + x_{31}).$$

Though equivalent, they are just different.

• We may (but is not required to) choose other format, such as  $0.6x_{11} - 0.4x_{21} - 0.4x_{31} \ge 0$  or  $3x_{11} - 2x_{21} - 2x_{31} \ge 0$ .

### Formulation: constraints

▶ In total we have four quality constraints:

- $x_{11} \ge 0.4(x_{11} + x_{21} + x_{31}).$
- $x_{21} \ge 0.2(x_{11} + x_{21} + x_{31}).$
- $x_{12} \ge 0.5(x_{12} + x_{22} + x_{32}).$
- $x_{13} \le 0.3(x_{12} + x_{22} + x_{32}).$

#### Formulation: constraints

▶ The demands are limited:

$$x_{11} + x_{21} + x_{31} \le 100$$

and

$$x_{12} + x_{22} + x_{32} \le 150.$$

▶ The quantities are nonnegative:

$$x_{ij} \ge 0 \quad \forall i = 1, ..., 3, j = 1, 2.$$

### Formulation: the complete formulation

▶ The complete formulation is

$$\begin{split} \max & 10(x_{11}+x_{21}+x_{31})+15(x_{12}+x_{22}+x_{32}) \\ & -8(x_{11}+x_{12})-4(x_{21}+x_{22})-3(x_{31}+x_{32}) \\ \text{s.t.} & x_{11} \geq 0.4(x_{11}+x_{21}+x_{31}) \\ & x_{21} \geq 0.2(x_{11}+x_{21}+x_{31}) \\ & x_{12} \geq 0.5(x_{12}+x_{22}+x_{32}) \\ & x_{13} \leq 0.3(x_{12}+x_{22}+x_{32}) \\ & x_{11}+x_{21}+x_{31} \leq 100 \\ & x_{12}+x_{22}+x_{32} \leq 150 \\ & x_{ij} \geq 0 \quad \forall i=1,...,3, j=1,2. \end{split}$$

# Remarks

- ▶ We may need to redefine decision variables when we find they are not enough.
- ▶ We may from time to time use multi-dimensional variables.
- ▶ We need to remove nonlinear constraints or objective functions, even if we just replace them with equivalent linear ones.

# Road map

- ▶ Resource allocation.
- Materials blending.
- ▶ Production and inventory.
- Compact formulations.

# **Production and inventory**

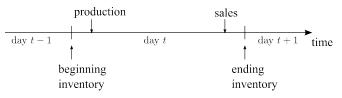
- ▶ When we are making decisions, we may need to consider what will happen in the **future**.
- ► This creates **multi-period** problems.
- In particular, in many cases products produced today may be stored and then sold in the future.
  - Maybe production is cheaper today.
  - Maybe the price is higher in the future.
- So the production decision must be jointly considered with the inventory decision.
- ▶ Introduced in Section 3.10 of the textbook.

# Production and inventory: the problem

- ▶ Suppose we are going to produce and sell a product in four days.
- ► For each day, there are different amounts of demands to fulfill.
  - ▶ Days 1, 2, 3, and 4: 100, 150, 200, and 170 units, respectively.
- ▶ The unit production costs are different for different days:
  - ▶ Days 1, 2, 3, and 4: \$9, \$12, \$10, and \$12 per unit, respectively.
- ▶ The prices are all **fixed**. So maximizing profits is the same as minimizing costs.

# Production and inventory: the problem

- We may store a product and sell it later.
  - ► The **inventory cost** is \$1 per unit per day.
  - ► E.g., producing 620 units on day 1 to fulfill all demands costs  $9 \times 620 + 1 \times 150 + 2 \times 200 + 3 \times 170 = 6640$  dollars.
- ► Timing:



- Beginning inventory + production sales = ending inventory.
- ► Inventory costs are assessed according to **ending inventory**.

▶ We need to determine the production quantities: Let

 $x_t$  = production quantity of day t, t = 1, ..., 4.

- ▶ Is that information enough?
  - ▶ E.g., given a plan (450, 0, 170, 0), we do not know whether the demand on day 4 is fulfilled with the productions on day 1 or 3.
- ▶ So we also need to determine the inventory quantities: Let

 $y_t$  = ending inventory of day t, t = 1, ..., 4.

▶ It is important to specify "ending"!

# Formulation: objective function

▶ We have production costs:

$$9x_1 + 12x_2 + 10x_3 + 12x_4.$$

▶ We also have inventory costs:

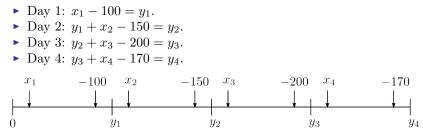
$$1(y_1 + y_2 + y_3 + y_4).$$

▶ So the objective function is

min  $9x_1 + 12x_2 + 10x_3 + 12x_4 + y_1 + y_2 + y_3 + y_4$ .

#### Formulation: constraints

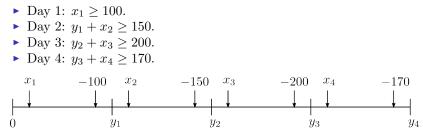
▶ We need to relate adjacent periods through ending inventories:



► This is typically called the **inventory balancing** constraint.

### Formulation: constraints

▶ We must satisfy all the demands at the moment of sales:



▶ Finally, all quantities must be nonnegative.

### Formulation: the complete formulation

#### ▶ The complete formulation is

 $9x_1 + 12x_2 + 10x_3 + 12x_4 + y_1 + y_2 + y_3 + y_4$ min s.t.  $x_1 - 100 = y_1$  $y_1 + x_2 - 150 = y_2$  $y_3 + x_3 - 200 = u_3$  $y_3 + x_4 - 170 = y_4$  $x_1 > 100$  $y_1 + x_2 > 150$  $y_2 + x_3 > 200$  $y_3 + x_4 > 170$  $x_t, y_t \ge 0 \quad \forall t = 1, \dots, 4.$ 

# Simplifying the formulation

- ▶ Let's look at the demand fulfillment constraints again.
- The first one is  $x_1 \ge 100$ .
  - But we have the first inventory balancing constraint  $x_1 100 = y_1$ and the nonnegativity constraint  $y_1 \ge 0$ . They together imply  $x_1 \ge 100$ .
  - Similarly,  $y_1 + x_2 150 = y_2$  and  $y_2 \ge 0$  imply  $y_1 + x_2 \ge 150$ .
- ▶ So all demand fulfillment constraints can be removed.

### Formulation: the simplified formulation

▶ The simplified formulation is

 $\begin{array}{ll} \min & 9x_1 + 12x_2 + 10x_3 + 12x_4 + y_1 + y_2 + y_3 + y_4 \\ \text{s.t.} & x_1 - 100 = y_1 \\ & y_1 + x_2 - 150 = y_2 \\ & y_3 + x_3 - 200 = y_3 \\ & y_3 + x_4 - 170 = y_4 \\ & x_t, y_t \geq 0 \quad \forall t = 1, ..., 4. \end{array}$ 

## Remarks

- The main idea is to use inventory variables to connect multiple periods. Otherwise periods will be unconnected.
- ▶ From time to time, we may first write some constraints and then find they are redundant.
- There are other ways of formulating this problem. For example, for the production lot on day t, we may split it into those for day t, those for day t + 1, etc.

# Road map

- ▶ Resource allocation.
- Materials blending.
- ▶ Production and inventory.
- ► Compact formulations.

## **Compact formulations**

- ▶ Most problems in practice are of **large scales**.
  - The number of variables and constraints are huge.
- ▶ Many variables can be grouped together:
  - E.g.,  $x_t$  = production quantity of day t, t = 1, ..., 4.
- ▶ Many constraints can be grouped together:
  - E.g.,  $x_t \ge 0$  for all t = 1, ..., 4.
- ► In modeling large-scale problems, we must use compact formulations to enhance readability and efficiency.

## **Compact formulations**

▶ In general, we may use the following three instruments:

- Indices  $(i, j, k, \ldots)$ .
- Summation  $(\sum)$ .
- ▶ For all  $(\forall)$ .
- ▶ For the joint production-inventory problem, let's write a compact formulation.

# **Production and inventory**

- ► The problem:
  - ▶ We have four periods.
  - ▶ In each period, we first produce and then sell.
  - ▶ Unsold products become ending inventories.
  - Want to minimize the total cost.
- ► Indices:
  - ▶ Because things will **repeat in each period**, it is natural to use an index for periods. Let  $t \in \{1, ..., 4\}$  be the index of periods.
- ▶ Now let's make the LP formulation compact.

## Compacting the objective function

- ▶ The original objective function:
  - min  $9x_1 + 12x_2 + 10x_3 + 12x_4 + y_1 + y_2 + y_3 + y_4$ .
- We may combine the last four terms:
  - min  $9x_1 + 12x_2 + 10x_3 + 12x_4 + \sum_{t=1}^4 y_t$ .
- ► To combine the first four terms, we may need to create a **parameter set**.
  - ▶ Denote  $C = [9 \ 12 \ 10 \ 12]$  as the **production cost vector** where  $C_t$  is the unit price on day t, t = 1, ..., 4.

• min 
$$\sum_{t=1}^{4} C_t x_t + \sum_{t=1}^{4} y_t$$
.

• min  $\sum_{t=1}^{4} (C_t x_t + y_t).$ 

## Compacting the constraints

- ▶ The original constraints:
  - ►  $x_1 100 = y_1$ ,
  - ►  $y_1 + x_2 150 = y_2$ ,
  - $y_2 + x_3 200 = y_3$ , and
  - ►  $y_3 + x_4 170 = y_4$ .
- ▶ Again, let's create a parameter set and group these constraints.
  - ▶ Denote  $D = [100 \ 150 \ 200 \ 170]$  as the demand vector where  $D_t$  is the demand on day t, t = 1, ..., 4.
- For day  $t, t = 2, ..., 4 : y_{t-1} + x_t D_t = y_t$ .
  - We cannot apply this to day 1 as  $y_0$  is undefined!
  - ▶ How may we group the four constraints together?

▶ Let's define  $y_0$ : Let

 $y_t$  = ending inventory of day t, t = 0, ..., 4.

- ▶ The ending inventory of day 0, by definition, should be the initial inventory of day 1.
- ▶ Then we may write

$$y_{t-1} + x_t - D_t = y_t \quad \forall t = 1, ..., 4$$

as the set of inventory balancing constraints.

• Certainly we need to set up the initial inventory:  $y_0 = 0$ .

### The complete compact formulation

▶ The compact formulation is

min 
$$\sum_{t=1}^{4} (C_t x_t + y_t)$$
  
s.t.  $y_{t-1} + x_t - D_t = y_t \quad \forall t = 1, ..., 4$   
 $y_0 = 0$   
 $x_t, y_t \ge 0 \quad \forall t = 1, ..., 4.$ 

- ▶ **Do not forget** " $\forall t = 1, ..., 4$ "! Without that, the formulation is just wrong.
- ► Nonnegativity constraints for multiple sets of variables can be combined to save some "≥ 0".

#### Parameters v.s. variables

- ▶ We need to define decision variables.
  - Let (Define)  $x_t$  = production quantity on day t, t = 1, ..., 4.
- ▶ We need to create parameter sets.
  - ▶ **Denote**  $C = [9 \ 12 \ 10 \ 12]$  **as** the production cost vector where  $C_t$  is the unit production cost on day t, t = 1, ..., 4.
- ▶ For parameters, we just define their names. We do not define parameters. They exist before we give them names!
- ▶ Variables do not exist before we define them.
- One convention is to:
  - Use **lowercase** letters for variables (e.g.,  $x_t$ ).
  - Use **uppercase** letters for parameters (e.g.,  $C_t$ ).

#### Parameters v.s. variables

▶ When creating parameter sets, it is fine to write only:

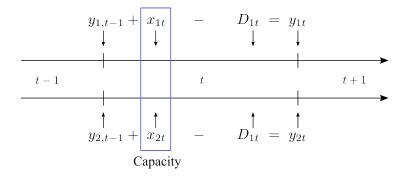
Denote  $C = [9 \ 12 \ 10 \ 12]$  as the production cost vector.

- $C_t$  is naturally its  $t^{\text{th}}$  element and has no ambiguity.
- ▶ The **values** should be indicated when defining the name.
- ▶ It is also fine to write

Denote  $C_t$  as the unit production cost on day t, t = 1, ..., 4.

- ▶ Do not need to specify values.
- ▶ Need to specify **range** through **indices**.
- ▶ In either case, we should indicate the **physical meaning**.

- Suppose we will produce and sell N products in T periods.
- ▶ We are given
  - ▶ The unit production cost of each product in each period,
  - ▶ The demand of each product in each period,
  - ▶ The holding cost of each product per period,
  - ▶ The machine time for producing one unit of each product, and
  - The capacity (measured in total machine time) of each day.
- ▶ How to write an LP that can minimize the total cost?



- Let  $N = \{1, 2, ..., 10\}, T = \{1, 2, ..., 100\}, \text{ and } T_0 = T \cup \{0\}.$
- ▶ For variable, let

 $x_{it}$  = production quantity of product *i* in period  $t, i \in N, t \in T$ , and  $y_{it}$  = ending of product *i* in period  $t, i \in N, t \in T_0$ .

▶ For parameters, denote

 $C_{it}$  as the unit production cost of product i in period  $t, i \in N, t \in T$ ,  $H_i$  as the unit inventory cost per period  $i \in N$ ,  $D_{it}$  as the demand of product i in period  $t, i \in N, t \in T$ ,  $P_i$  as the machine time required for product  $i, i \in N$ , and  $K_t$  as the machine time capacity in period  $t, t \in T$ .

▶ The problem can then be formulated as

$$\min \quad \sum_{i \in N} \sum_{t \in T} C_{it} X_{it} + \sum_{i \in N} H_i \sum_{t \in T} y_{it}$$
s.t. 
$$y_{i,t-1} + x_{i,t-1} - D_{i,t-1} = y_{it} \quad \forall i \in N, t \in T$$

$$y_{i0} = 0 \quad \forall i \in N$$

$$\sum_{i \in N} P_{it} x_{it} \leq K_t \quad \forall t \in T$$

$$x_{it}, y_{it} \geq 0 \quad \forall i \in N, t \in T.$$