Operations Research, Spring 2016

Pre-lecture Problems for Lecture 7: Integer Programming

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Note. The deadline of submitting the pre-lecture problem is 10:10am, April 7, 2015. Please submit a hard copy of your work in class. Late submissions will not be accepted. Each student must submit her/his individual work. Submit ONLY the problem that counts for grades.

1. (0 point) In this problem, we will use the branch-and-bound algorithm to solve the following IP

$$\begin{array}{ll} \max & 3x_1 + 5x_2 \\ \text{s.t.} & x_1 + x_2 \leq 16 \\ & x_2 \leq 7.5 \\ & x_i \in \mathbb{Z}_+ \quad \forall i = 1, 2. \end{array}$$

- (a) Solve the linear relaxation of the given IP. Show that both x_1 and x_2 are fractional in the optimal solution.
- (b) Branch on x_1 and continue until the IP is fully solved. Depict the branching tree.
- (c) Instead of branch on x_1 , branch on x_2 and continue until the IP is fully solved. Depict the branching tree. Compare the result with branching on x_1 .
- 2. (0 point) Consider the following integer program

$$\max x_1 + x_2$$
s.t.
$$2x_1 + x_2 - 6 \le M_1 z$$

$$x_1 + 2x_2 - 8 \le M_2 (1 - z)$$

$$x_1 \le 10$$

$$x_2 \le 10$$

$$x_i \ge 0 \quad \forall i = 1, 2$$

$$z \in \{0, 1\},$$

where M_1 and M_2 are parameters and x_1 , x_2 , and z are variables. The binary variable z is to select at least one constraint to be satisfied.

- (a) What values of M_1 and M_2 can enable z to do the "at-least-one" selection?
- (b) Depict the feasible region on the (x_1, x_2) plane. Then graphically solve the IP. **Hint.** The feasible region is nonconvex on the (x_1, x_2) plane. To solve this problem graphically, no branch-and-bound is needed.
- 3. (10 points) Use the branch-and-bound algorithm to solve the IP