Operations Research, Spring 2016 Suggested Solution for Pre-lecture Problems for Lecture 8

Solution providers: Kiwi Liu and Johnny Chen Department of Information Management National Taiwan University

1. (a) Let the demand set I and the location set J both be the six towns , who locate at $\{(0, 60), (20, 50), (30, 20), (40, 80), (50, 50), (90, 60)\}$. Suppose that there are 52 weeks in a year, the 5-year demand size is $h_i \in \{2600000, 3900000, 3120000, 8000, 2080000, 780000\}$, at demand $i \in I$. The distance between location $j \in J$ and demand $i \in I$ is d_{ij} . The fixed construction cost at location $j \in J$ is $f_j \in \{200000, 180000, 160000, 190000, 150000, 200000\}$. Let $x_j = 1$ if a DC is built at location $j \in J$ or 0 otherwise. Let $y_{ij} = 1$ is demand $i \in I$ is served by DC at location $j \in J$ or 0 otherwise.

$$\min \quad \sum_{i \in I} \sum_{j \in J} \frac{h_{ij}}{500} d_{ij} y_{ij} + \sum_{j \in J} f_j x_j$$
s.t. $y_i j \leq x_j \quad \forall i \in I, j \in J$

$$\sum_{j \in J} y_{ij} = 1 \quad \forall i \in I$$

$$x_3 = 1$$

$$x_j \in \{0, 1\} \quad \forall j \in J$$

$$y_i \in \{0, 1\} \quad \forall i \in I, j \in J.$$

- (b) Reset the location set J to {(0, 60), (20, 50), (30, 20), (40, 80), (50, 50), (90, 60), (0, 20), (20, 40), (40, 30), (60, 40)} and everything follows.
- 2. (a) For each job $j \in J = \{1, 2, ..., 10\}$, the processing time is p_j and the due time is d_j , where $p_j \in \{6, 9, 3, 5, 10, 6, 3, 9, 7, 10\}$ and $d_j \in \{50, 53, 55, 56, 59, 60, 62, 67, 68, 70\}$. let C_j be the completion time of job $j \in J$. Let $z_{ij} = 1$ if job j is before job i or 0 otherwise, $i \in J, j \in J, i < j$. Let $M = \sum_{j \in J} p_j$.

$$\begin{array}{ll} \min & \sum_{j \in J} T_j \\ \text{s.t.} & T_j \geq C_j - d_j & \forall j \in J \\ & C_i + p_j - C_j \leq M z_{ij} & \forall i \in J, j \in J, i < j \\ & C_j + p_i - C_i \leq M(1 - z_{ij}) & \forall i \in J, j \in J, i < j \\ & T_j \geq 0, C_j \geq 0 & \forall j \in J \\ & z_{ij} \in \{0, 1\} & \forall i \in J, j \in J, i < j. \end{array}$$

(b)

$$\begin{array}{ll} \min & \sum_{j \in J} T_j \\ \text{s.t.} & T_j \geq C_j - d_j & \forall j \in J \\ & C_i + p_j - C_j \leq M z_{ij} & \forall i \in J, j \in J, i < j \\ & C_j + p_i - C_i \leq M(1 - z_{ij}) & \forall i \in J, j \in J, i < j \\ & T_j \geq 0, C_j \geq 0 & \forall j \in J \\ & z_{15} = 0 \\ & z_{34} = 0 \\ & z_{57} = 0 \\ & z_{67} = 0 \\ & z_{69} = 0 \\ & z_{6(10)} = 0 \\ & z_{ij} \in \{0, 1\} & \forall i \in J, j \in J, i < j. \end{array}$$

3. (a) Following the notation in Problem 1, let $a_{ij} = 1$ if $d_{ij} < 40$ or 0 otherwise, $i \in I, j \in J$.

$$\begin{array}{ll} \min & \sum_{j \in J} x_j \\ \text{s.t.} & \sum_{j \in J} a_{ij} x_j \geq 1 \quad \forall i \in I \\ & x_3 = 1 \\ & x_j \in \{0,1\} \quad \forall j \in J. \end{array}$$

(b) Let V be the 6 towns.

Let E be all the paths between 6 towns.

Let d_{ij} be the distance of path $(i, j) \in E$. Let $x_{ij} = 1$ if the path between the town *i* and the town *j* is selected, $(i, j) \in E$.

$$\begin{array}{ll} \min & \sum\limits_{\substack{(i,j)\in E \\ s.t. \\ \sum\limits_{i\in V, i\neq k} x_{ik} = 1 \\ \sum\limits_{j\in V, j\neq k} x_{kj} = 1 \\ \sum\limits_{j\in S, j\in S, i\neq j} x_{ij} \leq |S| - 1 \\ x_{ij} \in \{0,1\} \\ \end{array} \quad \forall k \in V \\ \forall k \in V \\$$