Operations Research, Spring 2017 Suggested Solution For Homework 2

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1. (a) Let the decision variables be

 $p_1 =$ number of product 1 produced in a day, $p_2 =$ number of product 2 produced in a day, $p_3 =$ number of product 3 produced in a day,

 $\begin{array}{ll} \max & 10p_1 + 15p_2 + 25p_3 \\ \text{s.t.} & 2p_1 + 3p_2 + 5p_3 \leq 300 \\ & 3p_1 + 6p_2 + 7p_3 \leq 400 \\ & 4p_1 + 3p_2 + 9p_3 \leq 500 \\ & p_2 + 10p_3 \leq 600 \\ & p_1 \leq 100 \\ & p_2 \leq 80 \\ & p_3 \leq 50 \\ & p_i \geq 0 \quad \forall i = 1, 2, 3. \end{array}$

(b) The dual program is :

 $\begin{array}{ll} \min & 300y_1 + 400y_2 + 500y_3 + 600y_4 + 100y_5 + 80y_6 + 50y_7 \\ \mathrm{s.t.} & 2y_1 + 3y_2 + 4y_3 + y_5 \geq 10 \\ & 3y_1 + 6y_2 + 3y_3 + y_4 + y_6 \geq 15 \\ & 5y_1 + 7y_2 + 9y_3 + 10y_4 + y_7 \geq 25 \\ & y_i \geq 0 \quad \forall i = 1, ..., 7. \end{array}$

(c) According to the primal optimal solution, the number of material 1 be used is 280, material 2 be used is 400, material 3 be used is 500, material 4 be used is 503¹/₃. y₁, y₄, y₅ and y₆ will be 0 because the primal problem's corresponding constraints are not binding.

Constraints 1,2 and 3 in dual problem will be binding because p_1 , p_2 and p_3 are not 0.

(d) According to 1.(c) we know $y_1 = y_4 = y_5 = y_6 = 0$ and constraints 1, 2, and 3 are binding. Thus, the formulation become :

$$\begin{array}{ll} \min & 400y_2+500y_3+50y_7\\ {\rm s.t.} & 3y_2+4y_3=10\\ & 6y_2+3y_3=15\\ & 7y_2+9y_3+y_7=25\\ & y_i\geq 0 \quad \forall i=1,...,7. \end{array}$$

The optimal solution from y_1 to y_7 is (0, 2, 1, 0, 0, 0, 2), and the optimal objective value is 1400 for both primal and dual problem.

(e) The shadow price of primal problem's constraint 2 is 2 so you should buy material 2 at \$1.5 per unit.

The shadow price of primal problem's constraint 3 is 1 so you should not buy material 3 at \$1.5 per unit.

2. (a) The optimal solution is $x^3 = (0, 7)$ with optimal value $z^* = 14$.



(b)

$$\begin{array}{ll} \max & 10y_1 + 25y_2 \\ \text{s.t.} & y_1 + 3y_2 \leq 3 \\ & 2y_1 + 4y_2 \leq 2 \\ & y_i \geq 0 \quad \forall i = 1,2 \end{array}$$

(c) Solve the dual problem by simplex method then we get $y_1 = 0$ and $y_2 = 0.5$, so the shadow price of constraint 1 is 0 and the shadow price of constraint 2 is 0.5. $y_1 = 0, y_2 = 0.5$, and $z^{**} = 12.5$.

- 3. (a) The node with the highest objective value may have descendants with higher objective value of feasible solution. Then you can omit the other nodes which has lower objective value.
 - (b) You may find the integer solution faster.
- 4. (a) Let the decision variables be

$$x_{ij} = \begin{cases} 1 & \text{if item } i \text{ is chosen and put in bag } j \\ 0 & \text{otherwise.} \end{cases}, i = 1, ..., 12, j = 1, 2.$$

Let the parameters be

 I_i = importance of item i, i = 1, ..., 12, W_i = weight of item i., i = 1, ..., 12.

The IP formulation is

$$\max \sum_{j=1}^{2} \sum_{i=1}^{12} I_i x_{ij}$$
s.t.
$$\sum_{i=1}^{12} W_i x_{ij} \le 20 \quad \forall j = 1, 2$$

$$\sum_{j=1}^{2} x_{ij} \le 1 \quad \forall i = 1, ..., 12$$

$$x_{ij} \in \{0, 1\} \quad \forall i = 1, ..., 12 \quad \forall j = 1, 2.$$

(b) Let the decision variables be

$$x_{ij} = \begin{cases} 1 & \text{if item } i \text{ is chosen and put in bag } j \\ 0 & \text{otherwise.} \end{cases}, i = 1, ..., 12, j = 1, 2.$$

Let the parameters be

$$I_i$$
 = importance of item $i, i = 1, ..., 12,$
 W_i = weight of item $i_i, i = 1, ..., 12.$

The IP formulation is

$$\max \sum_{j=1}^{2} \sum_{i=1}^{12} I_i x_{ij}$$
s.t.
$$\sum_{i=1}^{12} W_i x_{ij} \le 20 \quad \forall j = 1, 2$$

$$\sum_{j=1}^{2} x_{ij} \le 1 \quad \forall i = 1, ..., 12$$

$$x_{2j} + x_{3j} \le 1 \quad \forall j = 1, 2$$

$$x_{4j} + x_{5j} + x_{6j} \le 2 \quad \forall j = 1, 2$$

$$\sum_{j=1}^{2} (x_{8j} + x_{9j} + x_{10,j} + x_{11,j} + x_{12,j}) \ge 2$$

$$\sum_{j=1}^{2} (x_{1j} + x_{2j}) \ge 1 - \sum_{j=1}^{2} x_{3j}$$

$$x_{ij} \in \{0, 1\} \quad \forall i = 1, ..., 12 \quad \forall j = 1, 2.$$

(c) Let the decision variables be

$$x_{ij} = \begin{cases} 1 & \text{if item } i \text{ is chosen and put in bag } j \\ 0 & \text{otherwise.} \end{cases}, i = 1, ..., 12, j = 1, 2.$$

Let the parameters be

$$I_i = \text{importance of item } i, i = 1, ..., 12,$$

 $W_i = \text{weight of item } i., i = 1, ..., 12.$

The IP formulation is

$$\begin{array}{ll} \min & k \\ \text{s.t.} & k \geq \sum_{i=1}^{12} (x_{i1} - x_{i2}) \\ & k \geq \sum_{i=1}^{12} (x_{i2} - x_{i1}) \\ & \sum_{j=1}^{2} \sum_{i=1}^{12} I_i x_{ij} = z^* \\ & \sum_{i=1}^{12} W_i x_{ij} \leq 20 \quad \forall j = 1, 2 \\ & \sum_{j=1}^{2} x_{ij} \leq 1 \quad \forall i = 1, ..., 12 \\ & x_{ij} \in \{0, 1\} \quad \forall i = 1, ..., 12 \quad \forall j = 1, 2. \end{array}$$

5. Let the decision variables be

$$\begin{aligned} x_i &= \begin{cases} 1 & \text{if an ambulance is located at district } i \\ 0 & \text{otherwise.} \end{cases}, i = 1, ..., 8, j = 1, ..., 8, \\ y_{ij} &= \begin{cases} 1 & \text{if an ambulance is located at district } i \text{ and is sent to district } j \\ 0 & \text{otherwise.} \end{cases}, i = 1, ..., 8, j = 1, ..., 8. \end{aligned}$$

Let the parameters be

$$T_{ij}$$
 = traveling time between district *i* and *j*, *i* = 1, ..., 8, *j* = 1, ..., 8, P_j = population in district *j*, *j* = 1, ..., 8.

The IP formulation is

$$\begin{array}{ll} \min & M \\ \text{s.t.} & M \geq y_{ij}T_{ij}P_j \\ & y_{ij} \leq x_i \quad \forall i=1,...,8 \quad \forall j=1,...,8 \\ & \displaystyle \sum_{i=1}^8 x_i = 2 \\ & \displaystyle \sum_{i=1}^8 y_{ij} = 1 \quad \forall j=1,...,8 \\ & x_i, y_{ij} \in \{0,1\} \quad \forall i=1,...,8 \quad \forall j=1,...,8. \end{array}$$

6. (a) The model of problem 4.b is

```
param I; #12
param B; #2
param Importance{i in 1..I};
param Weight{i in 1..I};
var x{i in 1..I, j in 1..B} binary;
```

```
maximize importance:
sum{j in 1..B}sum{i in 1..I}(Importance[i]*x[i,j]);
subject to WeightConstraint{j in 1..B}: #one bag can only afford 20 kilograms.
sum{i in 1..I}(Weight[i]*x[i, j])<=20 ;
subject to ItemConstraint{i in 1..I}: #one item can only be put in one bag.
sum{j in 1..B}(x[i, j])<=1 ;
subject to ChooseConstraint1{j in 1..B}: #items 2 and 3 cannot be put in the same bag.
x[2,j]+x[3,j]<=1 ;
subject to ChooseConstraint2{j in 1..B}: #items 4, 5 and 6 cannot be put in the same bag.
x[4,j]+x[5,j]+x[6,j]<= 2 ;
subject to ChooseConstraint3: #at least two of items 8 to 12 must be carried.
sum{j in 1..B}(x[8,j]+x[9,j]+x[10,j]+x[11,j]+x[12,j])>=2 ;
subject to ChooseConstraint4: #at least one of items 1 and 2 must be carried if item 3 is n
sum{j in 1..B}(x[1,j]+x[2,j])>=1-(sum{j in 1..B}(x[3,j])) ;
```

The optimal objective value of 4.b is 44, and the value of x are shown in Table 1. We should put item 1, 2, 5, 8 in bag 1, item 4, 9, 10, 11, 12 in bag 2. Moreover, there are three answers which are also correct. They are $\{(1, 2, 4, 5)(8, 9, 10, 11, 12)\}, \{(1, 2, 4, 9)(5, 8, 10, 11, 12)\}$ and $\{(1, 7, 8)(4, 9, 10, 11, 12)\}$.

	x[*][1]	x[*][2]
1	1	0
2	1	0
3	0	0
4	0	1
5	1	0
6	0	0
7	0	0
8	1	0
9	0	1
10	0	1
11	0	1
12	0	1

Table 1: Solution of our model

(b) The model of problem 5 is

```
param L; #8
param Time{i in 1..L, j in 1..L};
param Population{j in 1..L};
var x{i in 1..L} binary; #whether district i has an ambulance or not
var y{i in 1..L, j in 1..L} binary; #whether an ambulance can be sent from district i to j
var T; #maximum population-weighted firefighting time
minimize cost:
T;
```

```
subject to Constraint{i in 1..L, j in 1..L}:
T >= y[i,j]*Time[i,j]*Population[j];
subject to AvailableConstraint{i in 1..L, j in 1..L}:
y[i,j]<=x[i];
subject to ItemConstraint: #There are only two anbulances.
sum{i in 1..L}(x[i])=2 ;
subject to ChooseConstraint1{j in 1..L}: #One district can only be served by one ambulance.
sum{i in 1..L}(y[i,j])=1 ;
```

The optimal objective value of 5 is 135. We should locate the ambulances at district 1 and 6, and the separation plan is shown in Table 2. We should send an ambulance from district 1 to district 1, 2, 4 and send send an ambulance from district 6 to district 3, 5, 6, 7 and 8. Moreover, following answers are correct. They are $\{(1, 2, 5)(3, 4, 6, 7, 8)\}$ and $\{(1, 2)(3, 4, 5, 6, 7, 8)\}$.

	to 1	to 2	to 3	to 4	to 5	to 6	to 7	to 8
from 1	1	1	0	1	0	0	0	0
from 2	0	0	0	0	0	0	0	0
from 3	0	0	0	0	0	0	0	0
from 4	0	0	0	0	0	0	0	0
from 5	0	0	0	0	0	0	0	0
from 6	0	0	1	0	1	1	1	1
from 7	0	0	0	0	0	0	0	0
from 8	0	0	0	0	0	0	0	0

Table 2: Separation plan