

Part II: Policing

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Outline

- GCRA
- Leaky bucket
- Token bucket

Policing

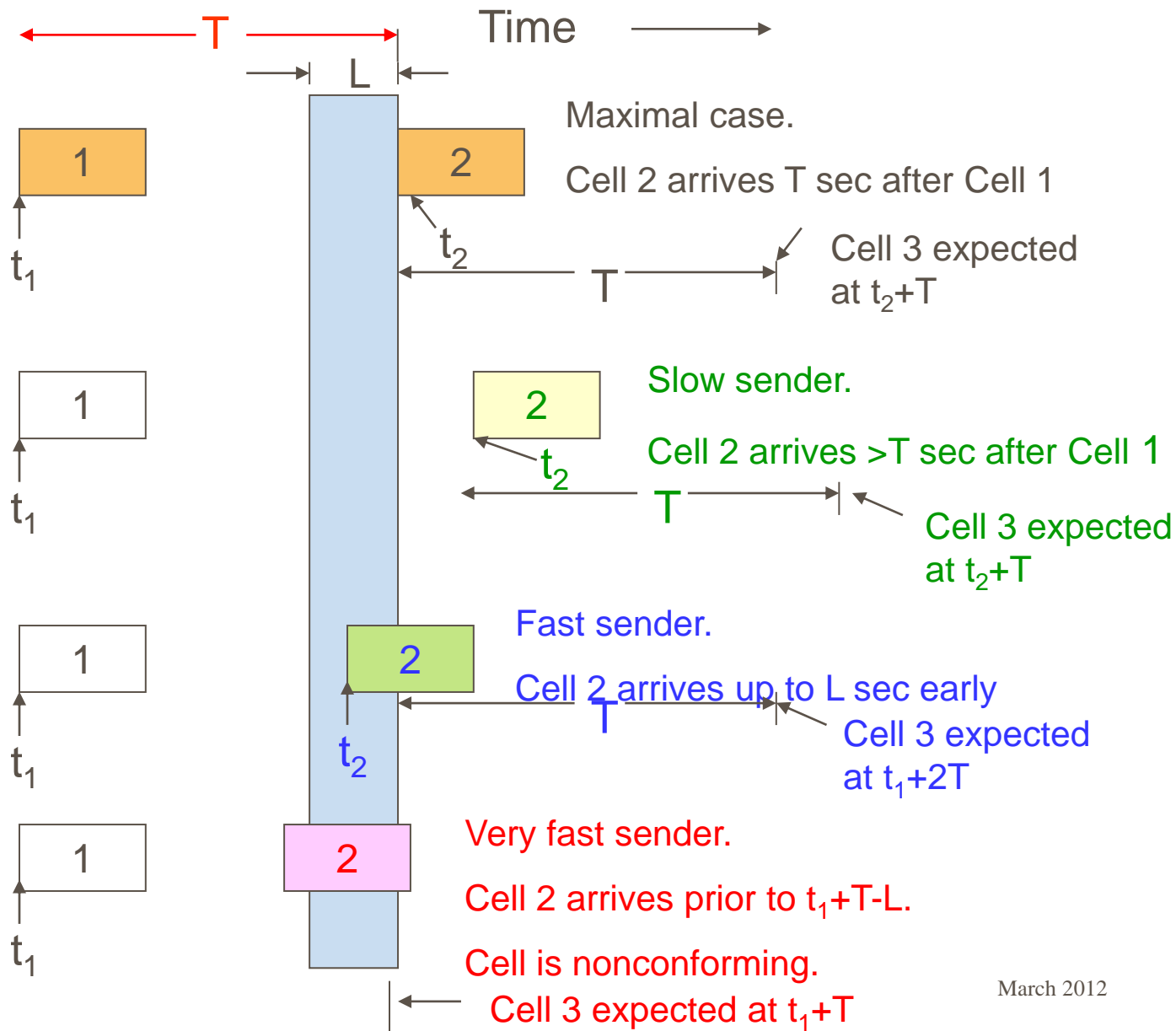
- The goal is to hold sources that are actually used to *committed* resources.
- Mechanisms are used to constraint traffic arrivals injected into the network
- Mechanisms
 - Generic Cell Rate Algorithm
 - Leaky bucket
 - Token bucket

Generic Cell Rate Algorithm (GCRA)

- Aka. Virtual Scheduling
- Used for conformance check
- Two parameters
 - T - Increment
 - L - Limit
- Regulated arrivals (T) with occasional *earlier* arrivals but not earlier than the limit (L)

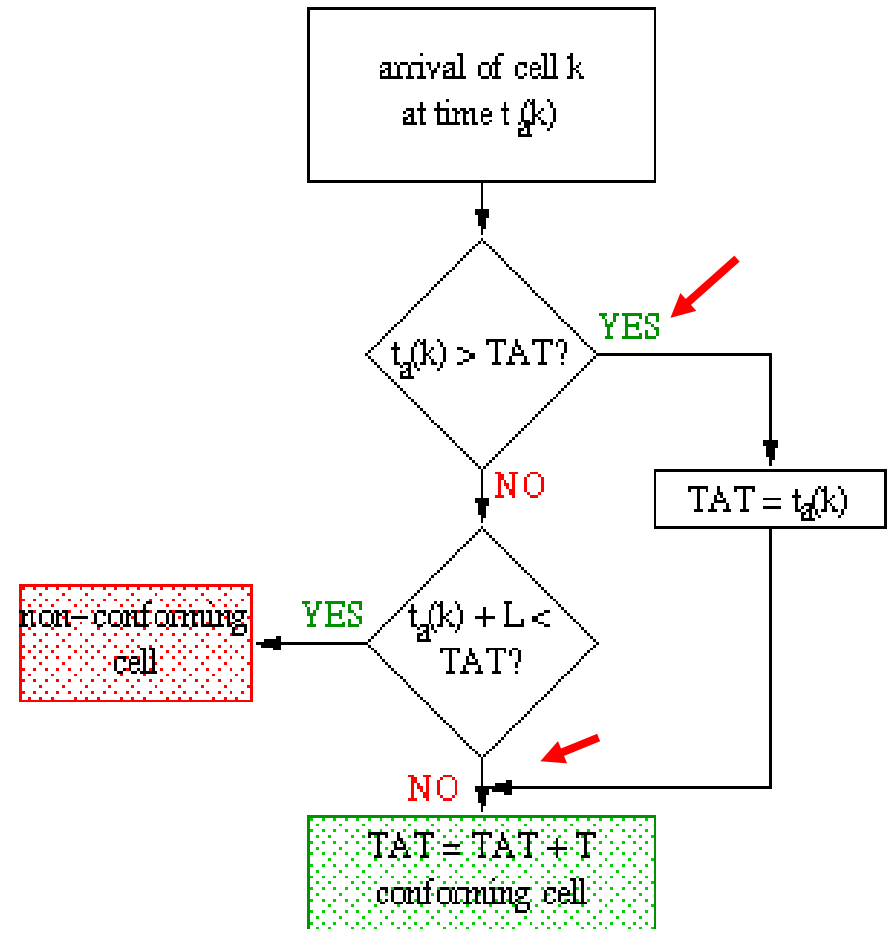
Conformance Check: cell arrival variation control

e.g., PCR
($T=1/PCR$)

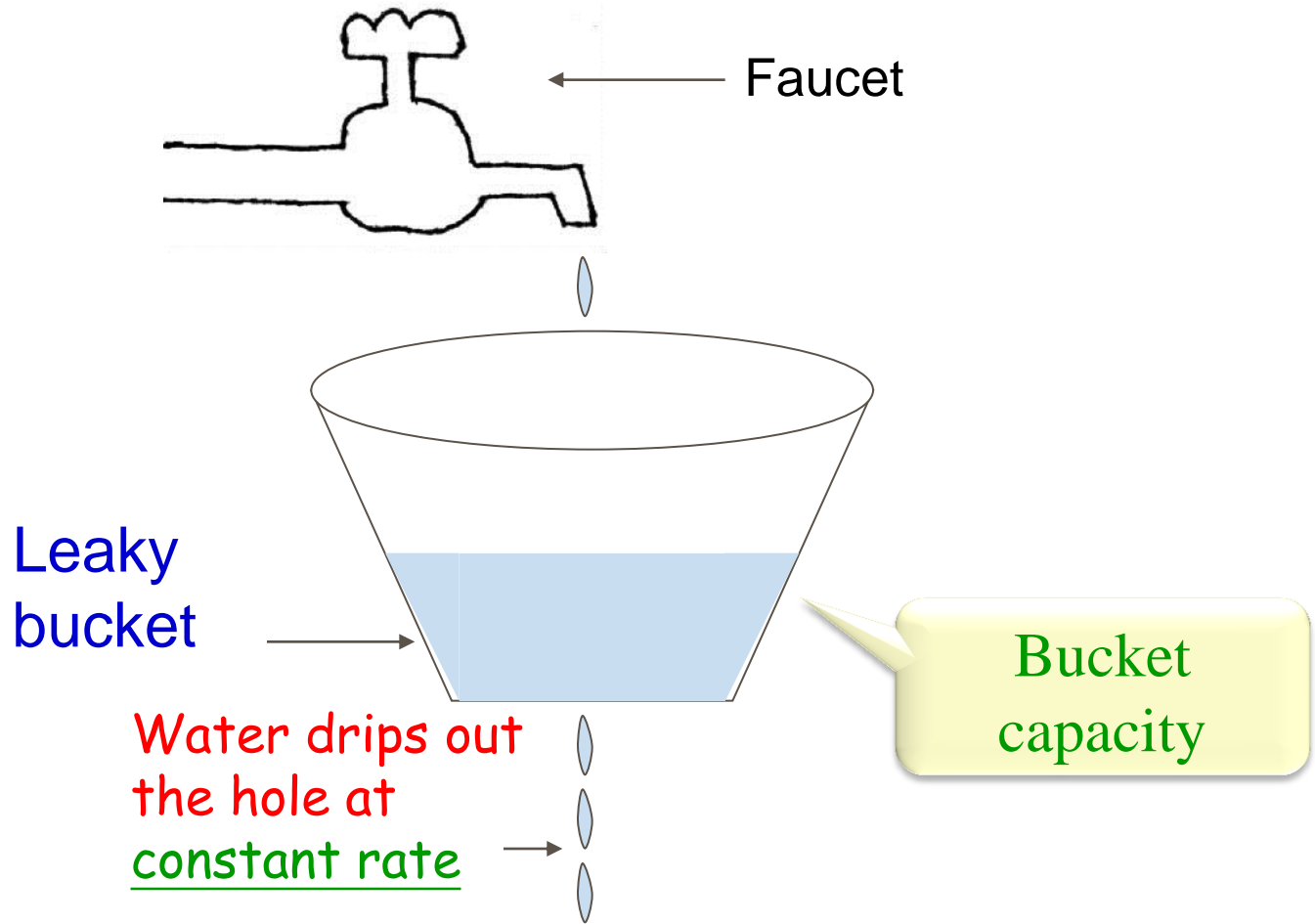


GCRA Algorithm

- Cell arrival time vs. **TAT** (Theoretical Arrival Time)
 - (1) If cell arrival time $>$ curnt_TAT ,
cell is *conforming* and
 $\text{next_TAT} = \text{cell arrival time} + T$ // slow arrival
 - (2) If cell arrival time $<$ curnt_TAT but *not* earlier than limit L ,
cell is *conforming*;
 $\text{next_TAT} = \text{curnt_TAT} + T$ // early arrival but ok
 - (3) Otherwise, cell arrives *too early* than the limit L and *non-conforming*



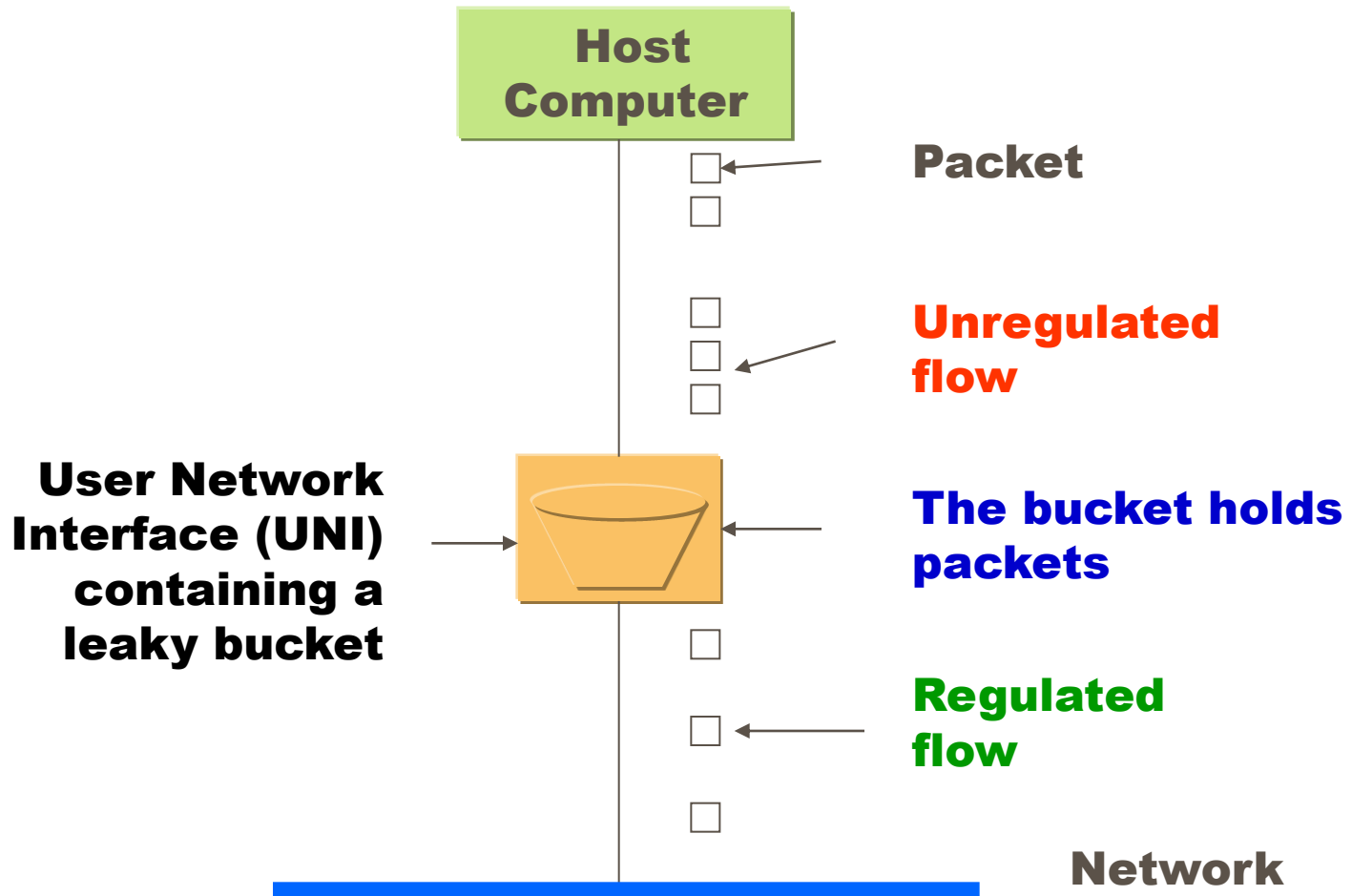
A "Leaky Bucket" with Water

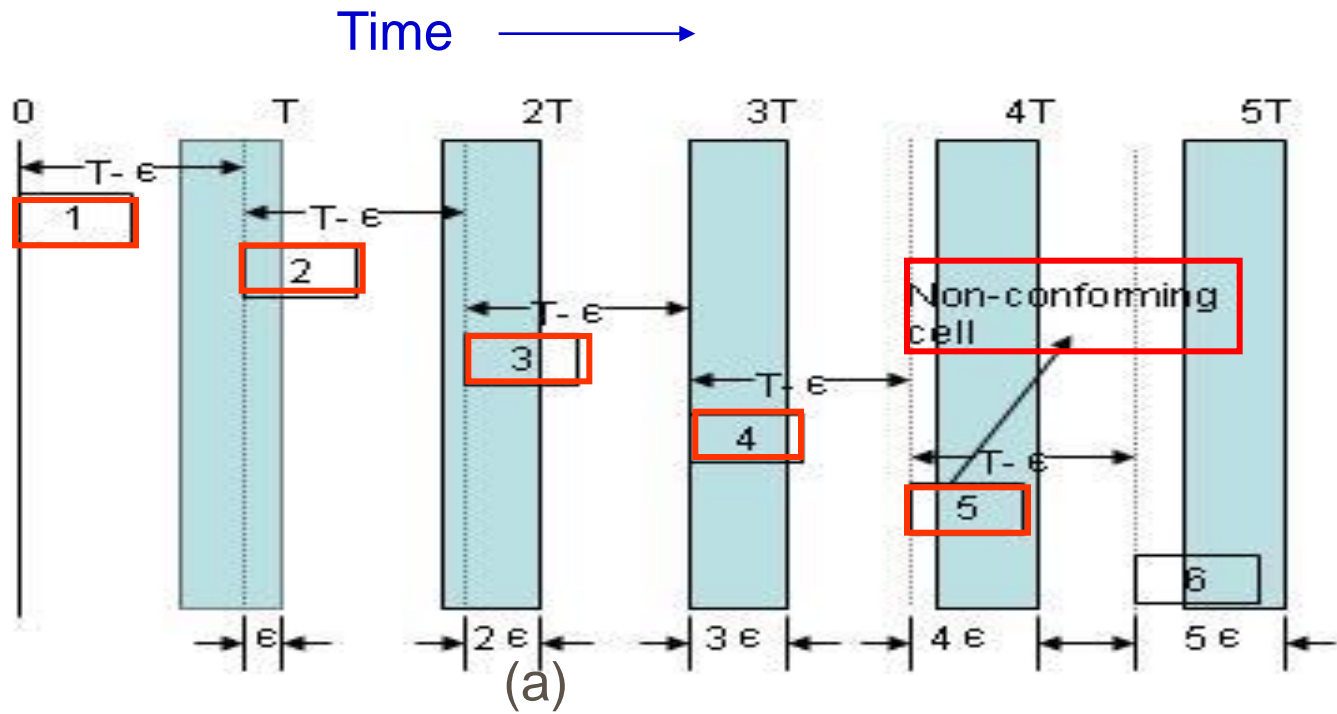


"Continuous-State" Leaky Bucket

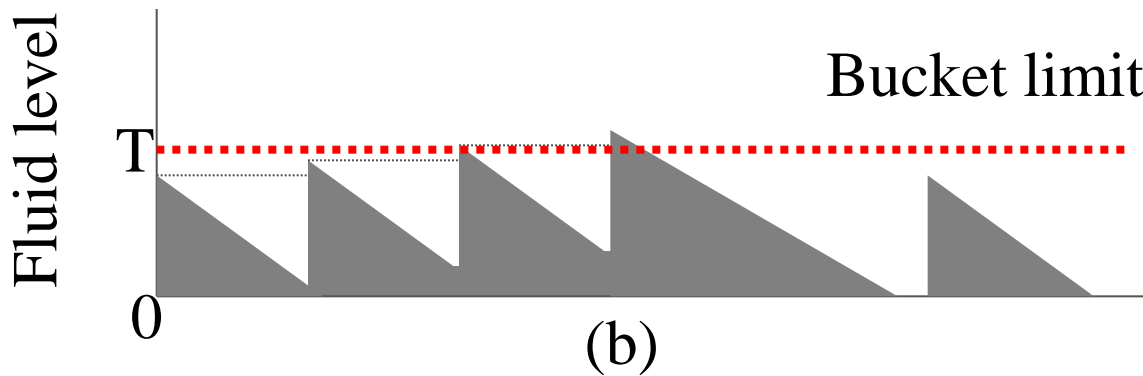
- **Draining rate** - drain-out at a continuous rate of one unit of content per time unit
- Content increased by the increment T for each conforming cell (i.e. work load brought by a cell)
- Follow GCRA
- **Bucket capacity is $T+L$**

A Leaky Bucket with "Packet"





(a) Packet arrivals - a sender trying to cheat



(b) Viewed in terms of leaky bucket

Maximum Burst Tolerance (MBS)

- The maximum number of cells that can be transmitted **back-to-back** given the peak rate (PCR) (bits/sec) constraint
 - L: allowable amount of early arrivals (**bucket capacity = $T+L$**)
 - δ : cell transmission time (sec/cell)
 - T: 1/PCR (sec/cell)
- Assume drainage does not start until the first cell has been entirely transmitted.
- Used in traffic conformance check (i.e., cell variation delay tolerance (CVDT))

- N: the number of conforming cells arrived back-to-back at the peak cell rate.
- $\text{Total_in} - \text{Total_out} = \text{left_over_in_bucket}$

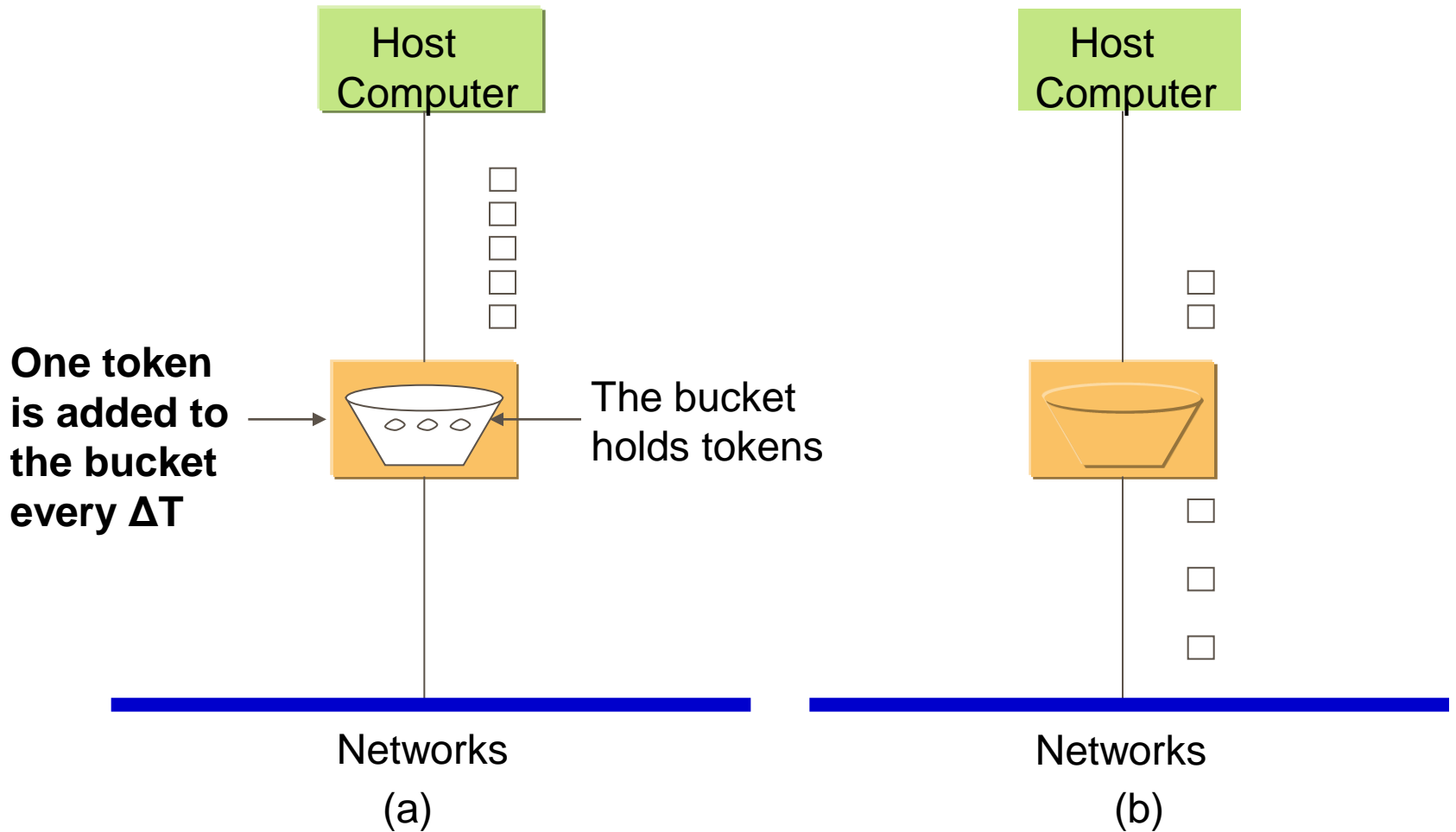
$$NT - (N - 1)\delta = T + L$$

$$N = 1 + \frac{L}{T - \delta}$$

SCR: sustainable cell (packet) rate

- An upper bound of “average rate”
- Average rate = total cells transmitted / duration of the connection

Token Bucket (1/2)



Token Bucket (2/2)

- The bucket holds tokens
- Tokens are generated by a clock at the rate of one token every ΔT sec.
- For a packet to be transmitted, it must capture and destroy a token.
- The bucket can save up tokens to allow **burst arrivals and burst departure**.
 - The maximum number of tokens saved up is the **bucket size**.
 - The maximum number of packets can be sent at once is the **bucket size**.
- Tokens are thrown away when the bucket fills up.
- Packets are never discarded.

Token Bucket: the length of the maximum rate burst

- Let
 - M : the maximum output rate of the token bucket
 - ρ : token generation rate
 - C: bucket size
 - S: duration to output all data
- **Traffic Input = Traffic output** (the balance equation)
 $C \text{ (bits)} + \rho S \text{ (bits/sec} \times \text{sec)} = MS \text{ (bits/sec} \times \text{sec)}$

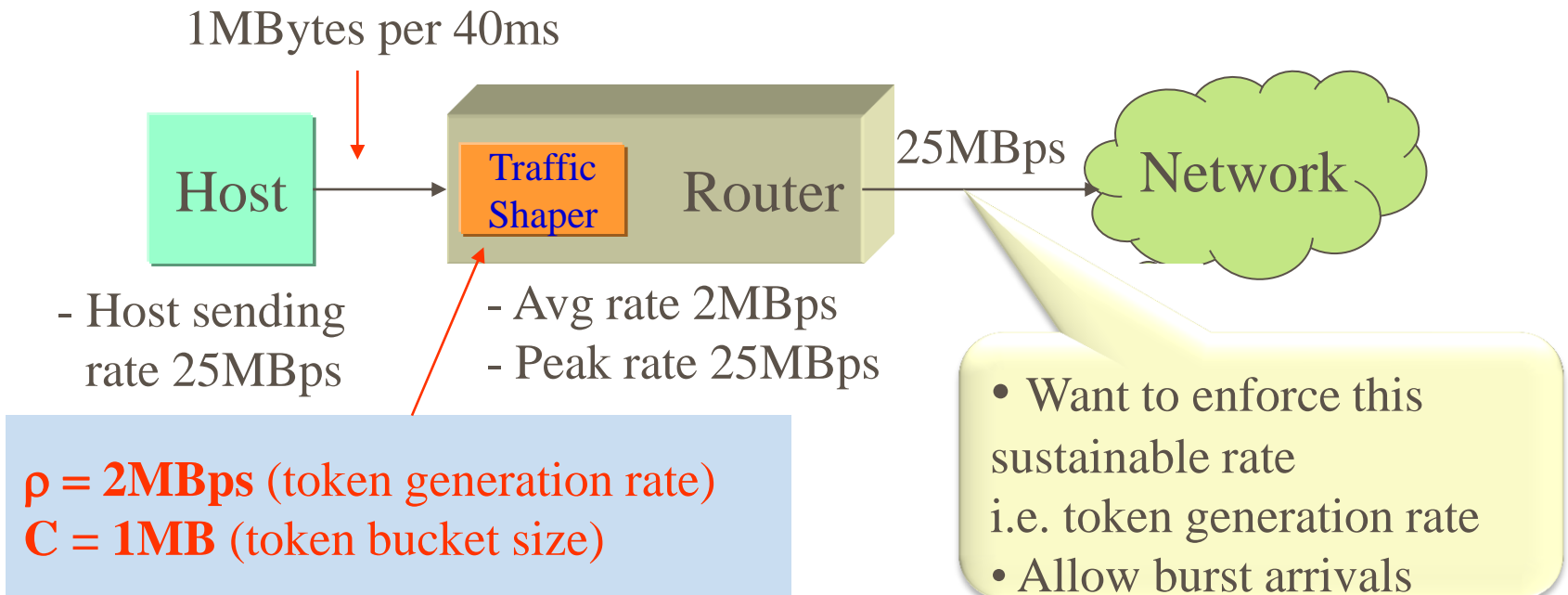
$$\mathbf{S = C / (M - \rho)}$$

- the time duration having the maximum output rate

- Example: $S = 250\text{KB} / (25\text{MB} - 2\text{MB})$
 $= 10.8\text{ms}$

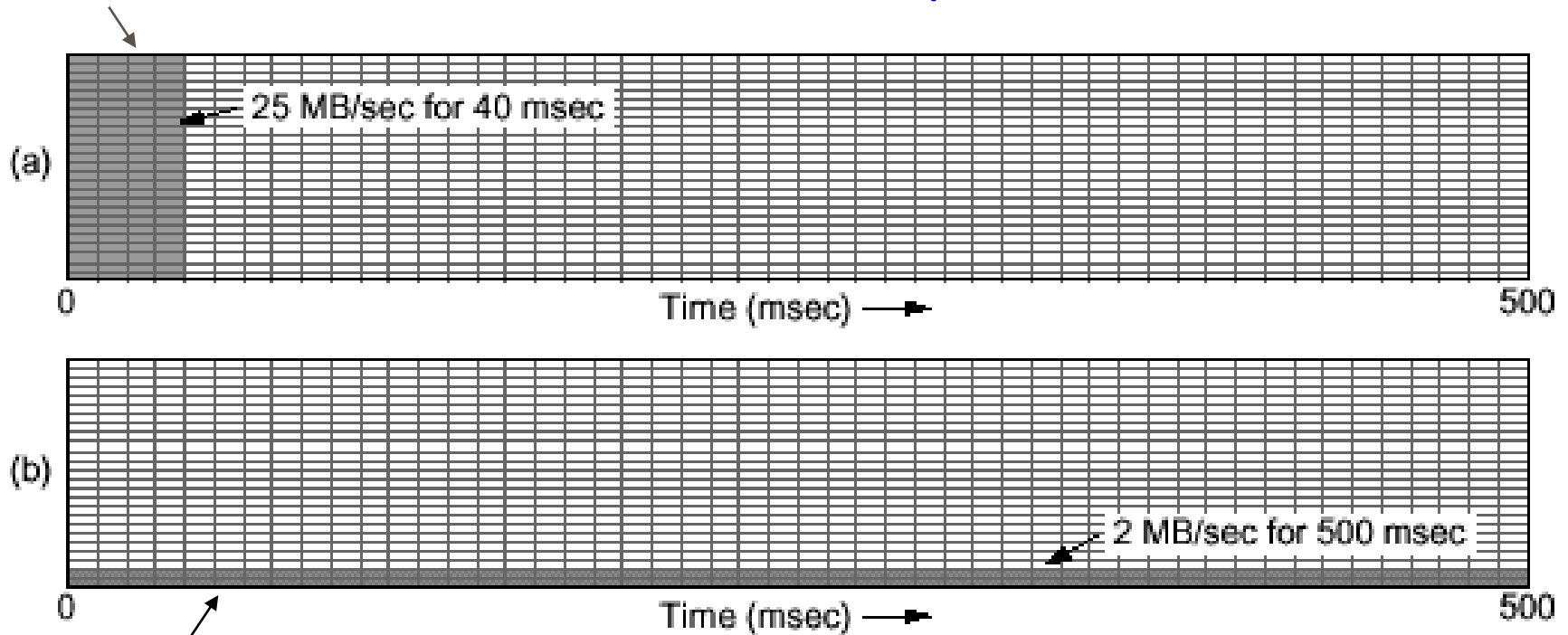
Traffic Control Example: using Leaky Bucket

- A burst of up to 1MBytes can be handled without data loss



Input : 25MB/40ms

Leaky Bucket



Output: constant 2MBps

Legend:

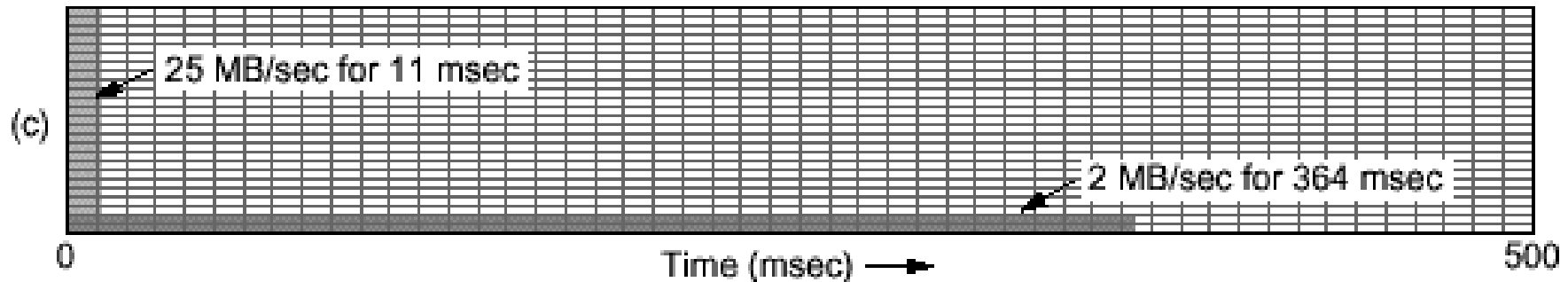
- 10ms per column

$\rho = 2\text{MBps}$ (drain-out rate)
 $C = 1\text{MB}$ (leaky bucket size)

$25\text{M} \times 40\text{ms} = 2\text{M} \times Y$
 $Y = 500\text{ms}$

Traffic Control Example: using Token Bucket (1/3)

- Assume
 - Token bucket is full (of token) when the $N=1\text{MB}$ burst arrives
 - Token bucket size $C= 250\text{KBytes}$;
 - The bucket drains at the full $M=25\text{MB/sec}$ for 11msec
 - ρ : token arrival rate (bytes/sec) = 2MB/sec
- S : maximum allowable burst length (sec)
 - > $25\text{MB} \times S = 250\text{KB} + 2\text{MB} \times S$
 - > $S = 250\text{KB}/(25\text{MB}-2\text{MB}) = 10.8\text{ms}$
- Drain out duration Y (sec)
 - > $MS + \rho Y = N$
 - > $25\text{MB} \times 10.8\text{ms} + 2\text{MB} \times Y = 1\text{MB}$
 - > $Y = 364.5\text{ms}$



Traffic Control Example: using Token Bucket (2/3)

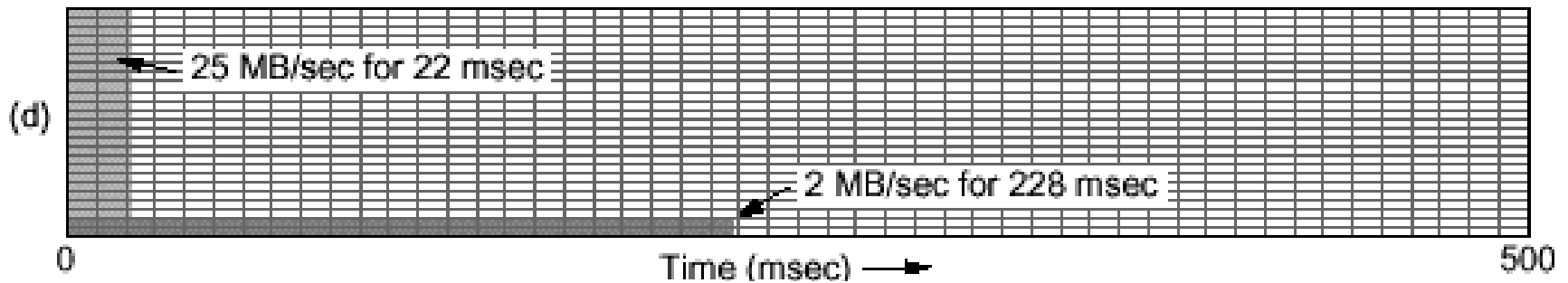


Figure 5.25 (d) - Find the duration with maximum output rate 25Mps

- Given token bucket size = 500KB
- The maximum allowable burst size (sec) (i.e., $S=C/(M-\rho)$)
-> $S=500KB/(25MB-2MB)=21.739ms\sim 22ms$
- The duration with average output rate 2Mps
-> $25MB \times 22ms + 2MB \times Y = 1MB$
-> $Y=228.26ms$

Traffic Control Example: using Token Bucket (3/3)

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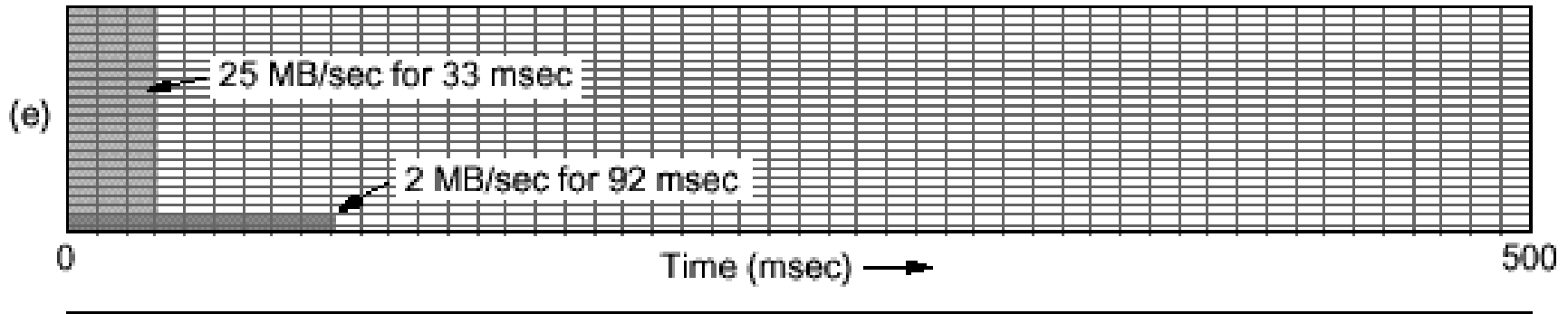


Figure 5.25 (e) - Find the duration with maximum output rate 25Mps

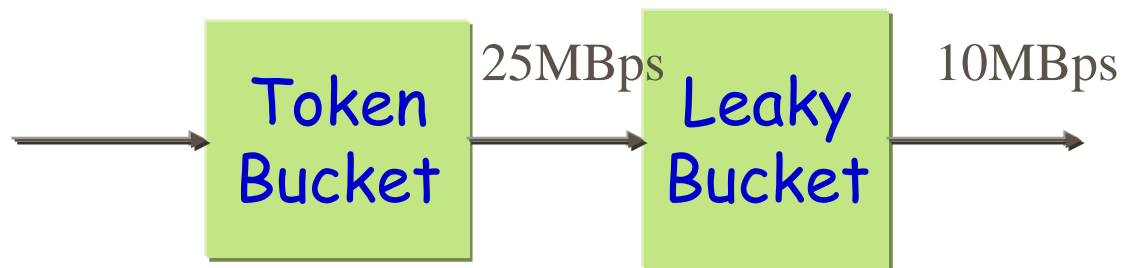
- Given token bucket size = 750KB
- The maximum allowable burst size (sec) (i.e., $S=C/(M-\rho)$)
-> $S=750KB/(25MB-2MB)=32.609\sim 33ms$
- The duration with average output rate 2Mps
-> $25MB \times 33ms + 2MB \times Y = 1MB$
-> $Y=92.391ms$

Discussion

- 250KB → 25MB x 10.8ms, Y = 364.5ms
 - 500KB → 25MB x 22ms, Y=228.26ms
 - 750KB → 25MB x 33ms, Y=92.391ms
-
- Which input arrivals are more bursty to the network?
 - Allow “bursts”!

Concatenate a token bucket with a leaky bucket

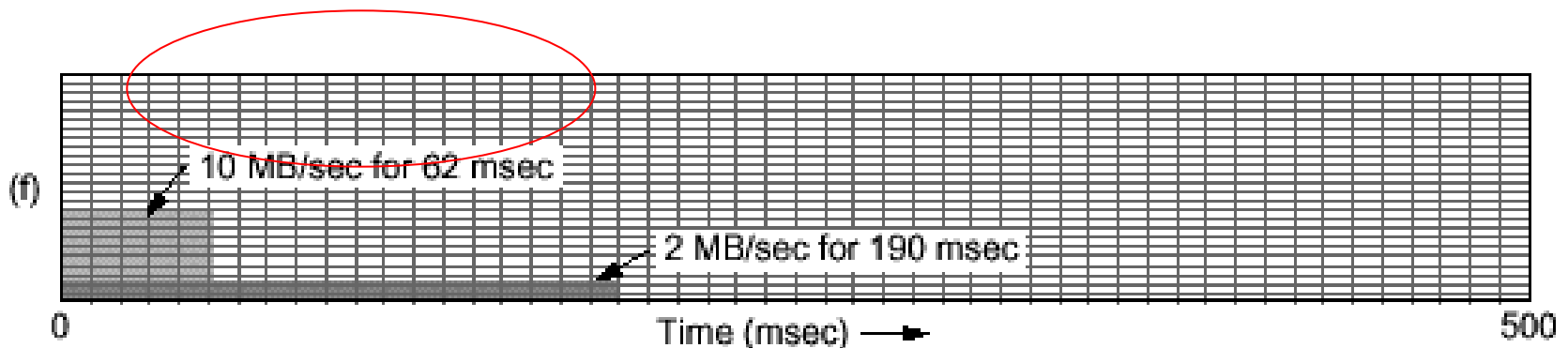
- Figure 5.25 (f)
 - To further “smooth” the traffic (less bursty output)
 - To concatenate a token bucket with a leaky bucket



- $\rho = 2\text{MBps}$ (token generation rate)
- $C = 1\text{MB}$ (token bucket size)

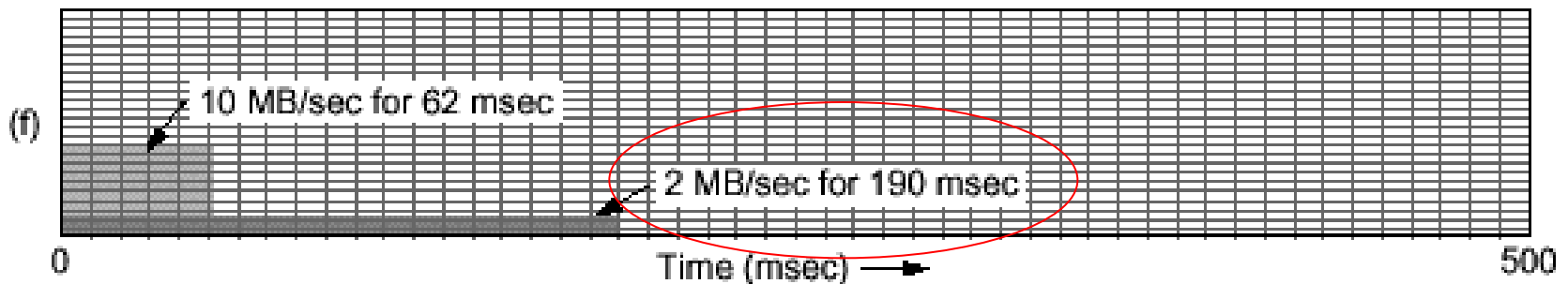
Concatenate a token bucket with a leaky bucket (cont'd)

- $\text{TokenBucket_max_output_rate} \times \text{max_rate_duration} + \text{token_gen_rate} \times \text{avg_rate_duration}$
 $= \text{LeakyBucket_max_output_rate} \times (\text{max_rate_duration} + \text{avg_rate_duration})$
-> $MS + \rho * X = L * (S + X)$
- Average rate duration X
 - From figure 5.25 (d), we have $S=22\text{ms}$
 - > $X = (25\text{MB} - 10\text{MB}) * 22\text{ms} / (10\text{MB} - 2\text{MB})$
 $= 40.761\text{ms}$
- The duration of maximum output rate from leaky bucket is $S+X$
-> $40.71 + 22 = 62.71\text{ms}$



Concatenate a token bucket with a leaky bucket (cont'd)

- The duration of having average rate output after token+leaky buckets traffic shaping
 - Total_traffic_at_leaky_bucket_max_rate + total_traffic_at_avg_rate = total_traffic_amount
 - > $L*(S+X) + \rho*Y = 1\text{MB}$
 - > $Y = (1\text{MB} - 10\text{MB} * 62.5\text{ms}) / 2\text{MB} = 187.5\text{ms}$



Compare Leaky bucket and Token bucket

■ Leaky bucket

- does *not* allow idle hosts to *save* up permits or tokens
- does *not* allow *bursts of outputs*
- enforces a rigid output pattern at the average rate no matter how bursty the traffic is.
- Packet lost when the bucket is full

■ Token bucket

- *allows burst arrivals and bursty output*
- Sustainable rate and peak rate
- No packet loss

Implementations

- **Leaky bucket algorithm**
 - conformance check and traffic shaping
- **Token bucket algorithm**
 - transmission control
 - allows burst arrivals
- **Traffic smoothing *between* routers**
 - Self-similar traffic
 - Individual flows and traffic aggregate
- **Traffic regulation of host traffic output**

Summary

- Policing is necessary to ensure that traffic sources conform to the committed rates.
- A powerful mechanism in traffic regulation
- Leaky bucket is often used for peak rate regulation.
- Token bucket is often used to allow instantaneous bursty arrival while regulating the traffic on their average behaviors.

The end. 😊

Reference

- S. Tanenbaum, “Computer Networks,” 4th edition, Prentice Hall, 2003.