

# Ordered Sets and Fixpoints

(Based on [Davey and Priestley 2002])

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# Partial Orders

- Let  $P$  be a set.
- A *partial order*, or simply *order*, on  $P$  is a binary relation  $\leq$  on  $P$  such that:
  - $\forall x \in P, x \leq x$ , (**reflexivity**)
  - $\forall x, y, z \in P, x \leq y \wedge y \leq z \rightarrow x \leq z$ , (**transitivity**)
  - $\forall x, y \in P, x \leq y \wedge y \leq x \rightarrow x = y$ . (**antisymmetry**)
- A set  $P$  equipped with a partial order  $\leq$ , often written as  $\langle P, \leq \rangle$ , is called a *partially ordered set*, or simply *ordered set*, sometimes abbreviated as *poset*.
- A binary relation that is reflexive and transitive is called a *pre-order* or *quasi-order*.
- We write  $x < y$  to mean  $x \leq y$  and  $x \neq y$ .

# Examples of Ordered Sets


## $\langle \mathcal{N}, \leq \rangle$


  $\mathcal{N} = \{1, 2, 3, \dots\}$ , the set of natural numbers.

  $\leq$  is the usual “less than or equal to” relation.


Variant:  $\langle \mathcal{N}_0, \leq \rangle$  with  $\mathcal{N}_0 = \mathcal{N} \cup \{0\} = \{0, 1, 2, 3, \dots\}$ .


## $\langle \mathcal{P}(X), \subseteq \rangle$

  $\mathcal{P}(X)$  is the powerset of  $X$ , consisting of all subsets of  $X$ .

  $\subseteq$  is the set inclusion relation.

## $\langle \Sigma^*, \leq \rangle$

  $\Sigma^*$  is the set of all finite strings over the alphabet  $\Sigma$ .

  $\leq$  is the “is a prefix of” relation.

# Order-Isomorphisms

- 🌐 We want to be able to tell when two ordered sets are essentially the same.
- 🌐 Let  $\langle P, \leq_P \rangle$  and  $\langle Q, \leq_Q \rangle$  be two ordered sets.
- 🌐  $P$  and  $Q$  are said to be (*order-*)*isomorphic*, denoted  $P \cong Q$ , if there is a map  $\varphi$  from  $P$  onto  $Q$  such that  $x \leq_P y$  if and only if  $\varphi(x) \leq_Q \varphi(y)$ .
- 🌐 The map  $\varphi$  above is called an *order-isomorphism*.
- 🌐 For example,  $\mathcal{N}_0$  and  $\mathcal{N}$  are order-isomorphic with the successor function  $n \mapsto n + 1$  as the order-isomorphism.
- 🌐 An order-isomorphism is necessarily *bijjective* (one-to-one and onto). Therefore, an order-isomorphism  $\varphi : P \rightarrow Q$  has a well-defined inverse  $\varphi^{-1} : Q \rightarrow P$ .

# Chains and Antichains

- Let  $P$  be an ordered set.
- $P$  is called a *chain* if  $\forall x, y \in P, x \leq y \vee y \leq x$ , i.e., any two elements in  $P$  are comparable.
- For example,  $\langle \mathcal{N}, \leq \rangle$  is a chain.
- Alternative names for a chain are *totally ordered set* and *linearly ordered set*.
- $P$  is called an *antichain* if  $\forall x, y \in P, x \leq y \rightarrow x = y$ , i.e., no two distinct elements in  $P$  are ordered.
- Clearly, any subset of a chain (an antichain) is a chain (an antichain).
- We write  $\mathbf{n}$  to denote a chain of  $n$  elements and  $\bar{\mathbf{n}}$  an antichain of  $n$  elements.

# Sums of Ordered Sets

- 🌐 Let  $P$  and  $Q$  be two *disjoint* ordered sets.
- 🌐 The **disjoint union**  $P \uplus Q$  is defined by  $x \leq y$  in  $P \uplus Q$  if and only if
  1.  $x, y \in P$  and  $x \leq y$  in  $P$ , or
  2.  $x, y \in Q$  and  $x \leq y$  in  $Q$ .
- 🌐 The **linear sum**  $P \oplus Q$  is defined by  $x \leq y$  in  $P \oplus Q$  if and only if
  1.  $x, y \in P$  and  $x \leq y$  in  $P$ , or
  2.  $x, y \in Q$  and  $x \leq y$  in  $Q$ , or
  3.  $x \in P$  and  $y \in Q$ .

# Diagrams for Ordered Sets

🌐 All possible ordered sets with three elements:



**3**



**$\bar{2} \oplus 1$**



**$1 \oplus \bar{2}$**

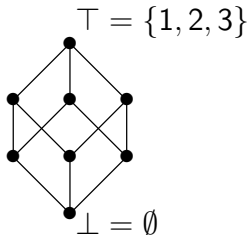


**$2 \uplus 1$**



**$\bar{3}$**

🌐  $\langle \mathcal{P}(\{1, 2, 3\}), \subseteq \rangle$ :



# Partial Maps

- 🌐 A (total) map or function  $f$  from  $X$  to  $Y$  is a binary relation on  $X$  and  $Y$  satisfying the following conditions:
  1. (**single-valued**) For every  $x \in X$ , there is **at most one**  $y \in Y$  such that  $(x, y)$  is related by  $f$ .  
In other words, if both  $(x, y_1)$  and  $(x, y_2)$  are related by  $f$ , then  $y_1$  and  $y_2$  must be equal.
  2. (**total**) For every  $x \in X$ , there is **at least one**  $y \in Y$  such that  $(x, y)$  is related by  $f$ .
- 🌐 A *partial map*  $f$  from  $X$  to  $Y$  is a *single-valued*, not necessarily total, binary relation on  $X$  and  $Y$ .
- 🌐 Representation of a total or partial map  $f$  from  $X$  to  $Y$  as a subset of  $X \times Y$ , or as an element of  $\mathcal{P}(X \times Y)$ , is called the *graph* of  $f$ , denoted  $\text{graph}(f)$ .



# Partial Maps as an Ordered Set

- 🌐 We write  $(X \dashrightarrow Y)$  to denote the set of all partial maps from  $X$  to  $Y$ .
- 🌐 For  $\sigma, \tau \in (X \dashrightarrow Y)$ , we define  $\sigma \leq \tau$  if and only if  $\text{graph}(\sigma) \subseteq \text{graph}(\tau)$ .  
In other words,  $\sigma \leq \tau$  if and only if whenever  $\sigma(x)$  is defined,  $\tau(x)$  is also defined and equals  $\sigma(x)$ .
- 🌐  $\langle (X \dashrightarrow Y), \leq \rangle$  is an ordered set.

# Programs as Partial Maps

- 🌐 Two programs  $P$  and  $Q$  with common sets  $X$  and  $Y$  respectively of *initial* states and *final* states may be seen as defining two partial maps  $\sigma_P, \sigma_Q : X \dashrightarrow Y$ .
- 🌐 The two programs might be related by  $\sigma_P \leq \sigma_Q$ , meaning that
  - ☀️ for any input state from which  $P$  terminates,  $Q$  also terminates, and
  - ☀️ for every case where  $P$  terminates,  $Q$  produces the same output as  $P$  does.
- 🌐 When  $\sigma_P \leq \sigma_Q$  does hold, we say  $P$  is refined by  $Q$  or  $Q$  refines  $P$ . (Some prefer the opposite.)
- 🌐 The refinement relation between two programs as defined is clearly a partial order.

# Order-Preserving Maps

- Let  $P$  and  $Q$  be ordered sets.
- A map  $\varphi : P \rightarrow Q$  is said to be **order-preserving** (or **monotone**) if  $x \leq y$  in  $P$  implies  $\varphi(x) \leq \varphi(y)$  in  $Q$ .
- The composition of two order-preserving maps is also order-preserving.
- A map  $\varphi : P \rightarrow Q$  is said to be an **order-embedding** (denoted  $P \hookrightarrow Q$ ) if  $x \leq y$  in  $P$  if and only if  $\varphi(x) \leq \varphi(y)$  in  $Q$ .

# Galois Connections and Insertions

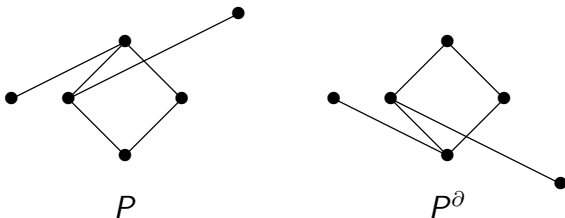
- Let  $P$  and  $Q$  be ordered sets.
- A pair  $(\alpha, \gamma)$  of maps  $\alpha : P \rightarrow Q$  and  $\gamma : Q \rightarrow P$  is a *Galois connection* between  $P$  and  $Q$  if, for all  $p \in P$  and  $q \in Q$ ,

$$\alpha(p) \leq q \leftrightarrow p \leq \gamma(q)$$

- Alternatively,  $(\alpha, \gamma)$  is a Galois connection between  $P$  and  $Q$  if, for all  $p, p_1, p_2 \in P$ ,  $q, q_1, q_2 \in Q$ ,
  - $p_1 \leq p_2 \rightarrow \alpha(p_1) \leq \alpha(p_2)$  and  $q_1 \leq q_2 \rightarrow \gamma(q_1) \leq \gamma(q_2)$   
(i.e.,  $\alpha$  and  $\gamma$  are monotone)
  - $p \leq \gamma(\alpha(p))$  and  $\alpha(\gamma(q)) \leq q$ .
- A *Galois insertion* is a Galois connection where  $\alpha \circ \gamma$  is the identity map, i.e.,  $\alpha(\gamma(q)) = q$ .

# Dual of an Ordered Set

- Given an ordered set  $P$ , we can form a new ordered set  $P^\partial$  (the “dual of  $P$ ”) by defining  $x \leq y$  to hold in  $P^\partial$  if and only if  $y \leq x$  holds in  $P$ .
- For a finite  $P$ , a diagram for  $P^\partial$  can be obtained by turning upside down a diagram for  $P$ :



# The Duality Principle

- For a statement  $\Phi$  about ordered sets, its **dual statement**  $\Phi^\partial$  is obtained by replacing each occurrence of  $\leq$  with  $\geq$  and vice versa.
- The Duality Principle:** Given a statement  $\Phi$  about ordered sets that is true for all ordered sets, the dual statement  $\Phi^\partial$  is also true for all ordered sets.

# Bottom and Top

- Let  $P$  be an ordered set.
- $P$  has a bottom element if there exists  $\perp \in P$  (“bottom”) such that  $\perp \leq x$  for all  $x \in P$ .
- Dually,  $P$  has a top element if there exists  $\top \in P$  (“top”) such that  $x \leq \top$  for all  $x \in P$ .
- $\perp$  is unique when it exists; dually,  $\top$  is unique when it exists.
- In  $\langle \mathcal{P}(X), \subseteq \rangle$ , we have  $\perp = \emptyset$  and  $\top = X$ .
- A finite chain always has a bottom and a top elements; this may not hold for an infinite chain.
- Given a bottomless  $P$ , we may form  $P_{\perp}$  ( $P$  lifted or the lifting of  $P$ ) by  $P_{\perp} \triangleq \mathbf{1} \oplus P$ .

# Maximal and Minimal Elements

- Let  $P$  be an ordered set and  $S \subseteq P$ .
- An element  $a \in S$  is a *maximal element* of  $S$  if  $a \leq x$  and  $x \in S$  imply  $x = a$ .
- If  $Q$  has a top element  $\top_Q$ , it is called the *greatest element* (or *maximum*) of  $Q$ .
- A *minimal element* of  $S$  and the *least element* (or *minimum*) of  $S$  (if it exists) are defined dually.



# Down-sets and Up-sets

- 🌍 Let  $P$  be an ordered set and  $S \subseteq P$ .
- 🌍  $S$  is a *down-set* (order ideal) if, whenever  $x \in S$ ,  $y \in P$ , and  $y \leq x$ , we have  $y \in S$ .
- 🌍 Dually,  $S$  is a *up-set* (order filter) if, whenever  $x \in S$ ,  $y \in P$ , and  $y \geq x$ , we have  $y \in S$ .
- 🌍 Given an arbitrary  $Q \subseteq P$  and  $x \in P$ , we define
  - ☀️  $\downarrow Q \triangleq \{y \in P \mid \exists x \in Q, y \leq x\}$  (“down  $Q$ ”),
  - ☀️  $\uparrow Q \triangleq \{y \in P \mid \exists x \in Q, y \geq x\}$  (“up  $Q$ ”),
  - ☀️  $\downarrow x \triangleq \{y \in P \mid y \leq x\}$ , and
  - ☀️  $\uparrow x \triangleq \{y \in P \mid y \geq x\}$ .
- 🌍  $\downarrow Q$  is the smallest down-set containing  $Q$  and  $Q$  is a down-set if and only if  $Q = \downarrow Q$ ; dually for  $\uparrow Q$ .

# Upper and Lower Bounds

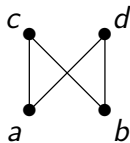
- Let  $P$  be an ordered set and  $S \subseteq P$ .
- An element  $x \in P$  is an *upper bound* of  $S$  if, for all  $s \in S$ ,  $s \leq x$ .
- Dually, an element  $x \in P$  is a *lower bound* of  $S$  if, for all  $s \in S$ ,  $s \geq x$  (or  $x \leq s$ ).
- The set of all upper bounds of  $S$  is denoted by  $S^u$  (“S upper”);  
 $S^u = \{x \in P \mid \forall s \in S, s \leq x\}$ .
- The set of all lower bounds of  $S$  is denoted by  $S^l$  (“S lower”);  
 $S^l = \{x \in P \mid \forall s \in S, s \geq x\}$ .
- By convention,  $\emptyset^u = P$  and  $\emptyset^l = P$ .
- Since  $\leq$  is transitive,  $S^u$  is an up-set and  $S^l$  a down-set.

# Least Upper and Greatest Lower Bounds

- 🌐 Let  $P$  be an ordered set and  $S \subseteq P$ .
- 🌐 If  $S^u$  has a least element, it is called the *least upper bound* (*supremum*) of  $S$ , denoted  $\sup(S)$ .
- 🌐 Equivalently,  $x$  is the least upper bound of  $S$  if
  - ☀  $x$  is an upper bound of  $S$ , and
  - ☀ for every upper bound  $y$  of  $S$ ,  $x \leq y$ .
- 🌐 Dually, if  $S^l$  has a greatest element, it is called the *greatest lower bound* (*infimum*) of  $S$ , denoted  $\inf(S)$ .
- 🌐 When  $P$  has a top element,  $P^u = \{\top\}$  and  $\sup(P) = \top$ . Dually, if  $P$  has a bottom element,  $P^l = \{\perp\}$  and  $\inf(P) = \perp$ .
- 🌐 Since  $\emptyset^u = \emptyset^l = P$ ,  $\sup(\emptyset)$  exists if  $P$  has a bottom element; dually,  $\inf(\emptyset)$  exists if  $P$  has a top element.

# Join and Meet

- 🌐 We write  $x \vee y$  (“ $x$  join  $y$ ”) in place of  $\sup(\{x, y\})$  when it exists and  $x \wedge y$  (“ $x$  meet  $y$ ”) in place of  $\inf(\{x, y\})$  when it exists.
- 🌐 Let  $P$  be an ordered set. If  $x, y \in P$  and  $x \leq y$ ,  $x \vee y = y$  and  $x \wedge y = x$ .
- 🌐 In the following two cases,  $a \vee b$  does not exist.



- 🌐 Analogously, we write  $\bigvee S$  (the “join of  $S$ ”) and  $\bigwedge S$  (the “meet of  $S$ ”).

# Lattices and Complete Lattices

- Let  $P$  be a *non-empty* ordered set.
- $P$  is called a *lattice* if  $x \vee y$  and  $x \wedge y$  exist for all  $x, y \in P$ .
- $P$  is called a *complete lattice* if  $\bigvee S$  and  $\bigwedge S$  exist for all  $S \subseteq P$ .  
Note: as  $S$  may be empty, the definition implies that every complete lattice is *bounded*, i.e., it has *top* and *bottom* elements.
- Every finite lattice is complete.

# Fixpoints

- Given an ordered set  $P$  and a map  $F : P \rightarrow P$ , an element  $x \in P$  is called a *fixpoint* of  $F$  if  $F(x) = x$ .
- The set of fixpoints of  $F$  is denoted  $\text{fix}(F)$ .
- The least element of  $\text{fix}(F)$ , when it exists, is denoted  $\mu(F)$ , and the greatest by  $\nu(F)$  if it exists.

# A Fixpoint Theorem for Complete Lattices

## Theorem (Knaster-Tarski Fixpoint Theorem)

Let  $L$  be a complete lattice and  $F : L \rightarrow L$  an order-preserving map. Then,

$$\mu(F) = \bigwedge \{x \in L \mid F(x) \leq x\}.$$

Dually,  $\nu(F) = \bigvee \{x \in L \mid x \leq F(x)\}.$

- 🌍 Let  $M = \{x \in L \mid F(x) \leq x\}$  and  $\alpha = \bigwedge M$ . We need to show (1)  $F(\alpha) = \alpha$  and (2) for every  $\beta \in \text{fix}(F)$ ,  $\alpha \leq \beta$ .
- 🌍 For all  $x \in M$ ,  $\alpha \leq x$  and so  $F(\alpha) \leq F(x) \leq x$ . Thus,  $F(\alpha) \in M'$  and hence  $F(\alpha) \leq \alpha (= \bigwedge M)$ .
- 🌍  $F(F(\alpha)) \leq F(\alpha)$ , implying  $F(\alpha) \in M$  and so  $\alpha \leq F(\alpha)$ .
- 🌍 For every  $\beta \in \text{fix}(F)$ ,  $\beta \in M$  and hence  $\alpha \leq \beta$ .

# Chain Conditions

- Let  $P$  be an ordered set.
- $P$  satisfies the **ascending chain condition** (ACC), if given any sequence  $x_1 \leq x_2 \leq \cdots \leq x_n \leq \cdots$  of elements in  $P$ , there exists  $k \in \mathbb{N}$  such that  $x_k = x_{k+1} = \cdots$ .
- Dually,  $P$  satisfies the **descending chain condition** (DCC), if given any sequence  $x_1 \geq x_2 \geq \cdots \geq x_n \geq \cdots$  of elements in  $P$ , there exists  $k \in \mathbb{N}$  such that  $x_k = x_{k+1} = \cdots$ .



# Directed Sets

- Let  $S$  be a *non-empty* subset of an ordered set.
- $S$  is said to be *directed* if, for every pair of elements  $x, y \in S$  there exists  $z \in S$  such that  $z \in \{x, y\}^u$ .
- $S$  is directed if and only if, for every finite subset  $F$  of  $S$ , there exists  $z \in S$  such that  $z \in F^u$ .
- In an ordered set with the ACC, a set is directed if and only if it has a greatest element.
- When  $D$  is directed for which  $\bigvee D$  exists, we write  $\bigsqcup D$  in place of  $\bigvee D$ .

# Complete Partial Orders (CPO)

- 🌐 An ordered set  $P$  is called a *Complete Partial Order (CPO)* if
  1.  $P$  has a bottom element  $\perp$  and
  2.  $\bigsqcup D$  exists for each directed subset  $D$  of  $P$ .
- 🌐 Alternatively,  $P$  is a CPO if **each chain of  $P$  has a least upper bound in  $P$ .**
- 🌐 Any complete lattice is a CPO.
- 🌐 For an ordered  $P$  satisfying Condition 2 above (called a pre-CPO), its lifting  $P_{\perp}$  is a CPO.

# Continuous Maps

- Let  $P$  and  $Q$  be CPOs.
- A map  $\varphi : P \rightarrow Q$  is said to be **continuous** if, for every directed set  $D$  in  $P$ ,
  - the subset  $\varphi(D)$  of  $Q$  is directed and
  - $\varphi(\bigsqcup D) = \bigsqcup \varphi(D)$ .
- A continuous map need not preserve bottoms, since by definition the empty set is not directed.
- A map  $\varphi : P \rightarrow Q$  such that  $\varphi(\perp) = \perp$  is called **strict**.

# A Fixpoint Theorem for CPOs


- 🌍 The  $n$ -fold composite  $F^n$  of  $F : P \rightarrow P$  is defined as follows.
  1.  $F^0$  is the identity.
  2.  $F^n = F \circ F^{n-1}$  for  $n \geq 1$ .
- 🌍 If  $F$  is order-preserving, so is  $F^n$ .

## Theorem (CPO Fixpoint Theorem I)






Let  $P$  be a CPO and  $F : P \rightarrow P$  an order-preserving map. Define  $\alpha \triangleq \bigsqcup_{n \geq 0} F^n(\perp)$ .

1. If  $\alpha \in \text{fix}(F)$ , then  $\alpha = \mu(F)$ .
2. If  $F$  is continuous, then  $\mu(F)$  exists and equals  $\alpha$ .

# Proof of CPO Fixpoint Theorem I (1)

-   $\perp \leq F(\perp)$ . So,  $F^n(\perp) \leq F^{n+1}(\perp)$ , for all  $n$ , inducing a chain in  $P$ :

$$\perp \leq F(\perp) \leq F^2(\perp) \leq \dots \leq F^n(\perp) \leq F^{n+1}(\perp) \leq \dots$$

-  Since  $P$  is a CPO,  $\alpha \triangleq \bigsqcup_{n \geq 0} F^n(\perp)$  exists.
-  Let  $\beta$  be any fixpoint of  $F$ ; we need to show that  $\alpha \leq \beta$ .
-  By induction,  $F^n(\beta) = \beta$ , for all  $n$ .
-  We have  $\perp \leq \beta$ , hence  $F^n(\perp) \leq F^n(\beta) = \beta$ .
-  The definition of  $\alpha$  then ensures  $\alpha \leq \beta$ .

# Proof of CPO Fixpoint Theorem I (2)

It suffices to show that  $\alpha \in \text{fix}(F)$ .

We have

$$\begin{aligned} F(\bigsqcup_{n \geq 0} F^n(\perp)) &= \bigsqcup_{n \geq 0} F(F^n(\perp)) && (F \text{ continuous}) \\ &= \bigsqcup_{n \geq 1} F^n(\perp) \\ &= \bigsqcup_{n \geq 0} F^n(\perp) && (\perp \leq F^n(\perp) \text{ for all } n) \end{aligned}$$

# Another Fixpoint Theorem for CPOs

## Theorem (CPO Fixpoint Theorem II)

*Let  $P$  be a CPO and  $F : P \rightarrow P$  an order-preserving map. Then  $F$  has a least fixpoint.*