



# Data Structures TA Session #2

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# Preface

Few things before we start

# When doing you homework

1. Always write **in nice & clear format**.
2. Please **DO indent** your code, and use monospaced fonts.
3. Answer all requirements one by one. Don't skip them!
4. Read the problems carefully. **Answer what they want!**  
(Not all problems ask you to write code.)



# Assignment #2



# Problem 2-1

Recursive function: `getSum()`

# What does this problem want?

- List the criteria of a recursive function
- **State how the function meets the criteria**
- Refer to your textbook for more information
- This problem doesn't ask you to write code, nor does it ask to draw a demonstration graph

# Criterion A.

- Define the problem in terms of a smaller problem of the same type
- One action of `getSum()` is to *call itself*
- Calculation of sum is made by adding the first element to the sum of the remaining array, which is smaller than the current array

# Criterion B.

- How does each recursive call diminish the size of the problem?
- At each recursive call to `getSum()`, the size of the array you need to compute is *diminished by 1*

# Criterion C.

- What instance of the problem can serve as the base case?
- The function handles the sum of  $x$  differently from all the other ones: It does not generate a recursive call. Rather, you know that `getSum(x)` is the element itself (`x[lower]`). Thus, the *base case* occurs when *lower=upper*.

# Criterion D.

- As the problem size diminishes, will you reach this base case?
- Given that  $\theta \leq \textit{Lower} \leq \textit{upper}$ , criterion B assures you that you will always *reach the base case*.

# Grading policy

- There are 4 questions you need to ask when writing a recursive function
- Each accounts for 5 points



# Problem 2-9

Digit sum of a given positive integer

# Let's think about recursion

- What's the base case?
- When  $N < 10$ , or  $N == 0$
- What's the answer of that?
- The sum is the digit itself
- What's the recursion formula?
- $N \% 10 + \text{getSum}( N / 10 )$

```
int getSum( int n )
{
    if ( n < 10 ) // or if ( n == 0 )
        return n;
    else
        return n % 10 + getSum( n / 10 );
}
```

```
int getSum( int n )
{
    return n < 10? n : n % 10 + getSum( n / 10 );
}
```

# Grading policy

- The base case 5
- The answer for the base case 5
- Other cases 5
- The answer for other cases 5



# Problem 2-16

Box trace of binary search

# Box trace

- Please refer to your textbook
- Shows the information during each iteration of the process

## Part A. Box 1 & 2

- Target = 5
- First = 0
- Last = 7
- Mid = 3
- Target < x[ 3 ]
- Search the left part

- Target = 5
- First = 0
- Last = 2
- Mid = 1
- Target = x[ 1 ]
- Return 1

## Part B. Box 1 & 2

- Target = 13
- First = 0
- Last = 7
- Mid = 3
- Target > x[ 3 ]
- Search the right part

- Target = 13
- First = 4
- Last = 7
- Mid = 5
- Target < x[ 5 ]
- Search the left part

## Part B. Box 2 & 3

- Target = 13
- First = 4
- Last = 7
- Mid = 5
- Target < x[ 5 ]
- Search the left part

- Target = 13
- First = 4
- Last = 4
- Mid = 4
- Target < x[ 4 ]
- Search the left part!

## Part B. Box 3 & 4

- Target = 13
  - First = 4
  - Last = 4
  - Mid = 4
  - Target < x[ 4 ]
  - Search the left part!
- Target = 13
  - First = 4
  - Last = 3
  - Mid = 3
  - First > Last
  - Return -1 (not found)

# Grading policy

- Part A: 2 boxes, 5 points each
- Part B: 4 boxes, 2.5 points each (your score is rounded up to the nearest integer)



# Problem 2-19

Indent the rabbit function

# How to solve this problem?

- Keep the recursion depth, either as a parameter or as a global variable
- Print the information of each function call after tabs

```
int rabbit( int n, string prefix = "" )
{
    int child = n <= 2;

    cout << prefix << "Enter rabbit: n = " << n << endl;

    if( n > 2 )
        child += rabbit( n - 1, prefix + '\t' ) +
        rabbit( n - 2, prefix + '\t' );

    cout << prefix << "Leave rabbit: n = " << n << "
value = " << child << endl;

    return child;
}
```

# Grading policy

- Function prototype 2
- Indention 4
- The “enter” statement 3
- The “leave” statement 3
- Base case 1
- Recursive call 3
- Return the answer of rabbit() 3
- Syntax correctness 1



# Problem 2-23

Euclidean algorithm

# Part A. the proof

To prove  $\gcd(a, b) = \gcd(b, a \bmod b)$ , given  $ab \neq 0$

Let  $X = \gcd(a, b)$ , then let  $a = mX$ ,  $b = nX$ .

Let  $a = bq + r$ ,  $q$  and  $r$  are integers,  $0 \leq r < b$

$$a - bq = r$$

By Common Divisor Divides Integer Combination,

1. all common divisors of  $a$  and  $b$  divide  $r$  (from  $a - bq = r$ )
2. all common divisors of  $b$  and  $r$  divide  $a$  (from  $bq + r = a$ )

## Part A. the proof (cont.)

1. all common divisors of  $a$  and  $b$  divide  $r$  (from  $a - bq = r$ )
2. all common divisors of  $b$  and  $r$  divide  $a$  (from  $bq + r = a$ )

Every common factor of  $(a, b)$  will appear in common factors of  $(b, r)$ , or  $(b, a \bmod b)$ . (The reverse also holds.)

Therefore, these 2 sets are equal.

$$\gcd(a, b) = \gcd(b, a \bmod b)$$

## Part B. when $b > a$ ...

- Suppose  $b > a$  in  $\text{gcd}( a, b )$
- The next recursive call will be  
 $\text{gcd}( b, a \bmod b ) = \text{gcd}( b, a )$
- The recursion swaps these 2 numbers and continues as usual without a problem

# Part C-1. Will it end?

- If  $b \mid a$ , then  $b$  is their GCD, and the function ends immediately (though the problem excludes such cases)
- Otherwise, the parameters  $a$  and  $b$  will get smaller each time
- Also assume that  $a > b > 0$  (refer to part B when  $b > a$ )
- $a > b$ , and  $b > a \bmod b$  (by definition)
- But  $a$  and  $b$  are both always greater than 0, so termination of the process can be done in finite steps

## Part C-2: Why the base case is appropriate?

- When  $a \bmod b = 0$ ,  $b$  is the greatest common divisor of  $a$  and  $b$  (by definition.)
- Therefore, no further recursion calls are required, and the base case is appropriate.

# Grading policy

- The proof 4
- When  $b > a$  10
- Reach the base case 3
- Why is base case appropriate 3



# Assignment #3

# Exercise 1.9

a.

```
cout << p.coefficient( p.degree() );
```

b.

```
p.changeCoefficient( p.coefficient( 3 ) + 8, 3 );
```

# Exercise 1.9

C.

```
polynomial<double> add( polynomial<double> a,  
polynomial<double> b )  
{  
    polynomial<double> sum;  
    int high = max( a.degree(), b.degree() );  
  
    for ( int i = 0; i <= high; ++i )  
        sum.changeCoefficient( a.coefficient( i )  
+ b.coefficient( i ), i );  
    return sum;  
}
```

# Exercise 1.9 - Grading Policy

- a(5)
  - correctness 1, display result 2, using ADT operation 2
- b(5)
  - correctness 3, using ADT operation 2
- c(10)
  - correctness 4, syntax correctness 3, using ADT operation 3

## Exercise 2.24(a)

$$c(1) = 0$$

$$c(2) = 1$$

$$c(3) = 3$$

$$c(4) = 7$$

$$c(5) = 15$$

$$\begin{aligned} C(n) &= 0, \text{ when } n=1 \\ &= 1, \text{ when } n=2 \\ &= 2 * C(n-1) + 1, \text{ otherwise} \end{aligned}$$

```
int C(int n)
{
    if(n==1) return 0;
    else if(n==2) return 1;
    else return 2*C(n-1)+1;
}
```

## Exercise 2.24(b)

$$c(1) = 0$$

$$c(2) = 1$$

$$c(3) = 2$$

$$c(4) = 4$$

$$c(5) = 6$$

$$c(6) = 10$$

$$c(7) = 14$$

$$c(8) = 21$$

$$c(9) = 29$$

$$C(n) = b(n, n-1)$$

$$b(n, m) = 0, \text{ when } n < 1 \text{ or } m < 1$$

$$= 1, \text{ when } n = 1 \text{ or } m = 1$$

$$= 1 + b(n, n-1), \text{ when } n = m$$

$$= b(n-m, m) + b(n, m-1), \text{ otherwise}$$

```
int b(int n, int m)
{
    if(n < 1 || m < 1) return 0;
    else if(n == 1 || m == 1) return 1;
    else if(n == m) return 1 + b(n, n-1);
    else return b(n-m, m) + b(n, m-1);
}
```

# Exercise 2.24(b)

- Some reference for you
- [OEIS](#): online encyclopedia of integer sequences
- [A000041](#)          number of partitions of  $n$
- [A000065](#)           $-1 +$  number of partitions of  $n$ .

# Exercise 2.24 - Grading Policy

- Each has 10 points
- Recursion function, 4
- Definition correctness, 6

# Exercise 3.1

```
Class ArrayBag: public BagInterface<ItemType>
{
    ...
public:
    double getAvg() const;
}
```

```
double ArrayBag::getAvg() const
{
    double avg = 0, sum = 0;
    for ( int I = 0; I < itemCount; i++)
        sum += items[i];
    avg = sum / itemCount;
    return avg;
}
```

# Exercise 3.1

- Reference: [accumulate\(\)](#) in [#include<numeric>](#)
- You can use it to sum up the values

# Exercise 3.1 - Grading Policy

- using client function, 4
- function correctness, 4
- syntax correctness, 2

## Exercise 3.5

```
class Inventory {  
private:  
    string name;  
    int cost, quantity;  
public:  
    Inventory( const string Name, const int Cost,  
const int Quantity ): name( Name ), cost( Cost ),  
quantity( Quantity ) {}  
    Inventory(): name( "" ), cost( 0 ), quantity( 0 ) {}  
};
```

## Exercise 3.5 (cont'd)

```
string getName() const { return name; }  
void setName( const string val ) { name = val; }  
int getCost() const { return cost; }  
void setCost( const int val ) { cost = val; }  
int getQuantity() const { return quantity; }  
void setQuantity( const int val ) { quantity = val; }  
}
```

# Exercise 3.5 - Grading Policy

- class, 4
- syntax correctness, 5
- operation (look at, change value), 16
- attribute (product, price, quantity, date, rating...), 15

# Exercise 3.9

```
template<typename T>
ArrayBag<T>::ArrayBag( const int val[], int size )
{
    itemCount = std::min( size, DEFAULT_CAPACITY );
    maxItems = DEFAULT_CAPACITY;
    for ( int i = 0; i < itemCount; ++i )
        items[ i ] = val[ i ];
}
```

# Exercise 3.9

- Reference: [copy\(\)](#) in [#include<algorithm>](#)
- You can use it to copy the values into the bag

# Exercise 3.9 - Grading Policy

- template, 2
- constructor function, 2
- initializing itemCount and maxItems, 2
- create a bag, 3
- syntax correct, 1



# The end~

Hope you did a good job in this assignment.  
Average score for assignment #2 & #3 is 79.7

By the TAs

104/11/02