

Suggested Solutions for Homework #1

(prepared by Jen-Feng Shih)

1. Determine, using
- Truth Tables*
- , the validity of the following propositions: (30 points)

(a) $(p \vee q \rightarrow r) \rightarrow (p \rightarrow r) \wedge (q \rightarrow r)$

Solution. Let Φ_a denote $(p \vee q \rightarrow r) \rightarrow (p \rightarrow r) \wedge (q \rightarrow r)$.

p	q	r	$p \vee q$	$p \vee q \rightarrow r$	$p \rightarrow r$	$q \rightarrow r$	$(p \rightarrow r) \wedge (q \rightarrow r)$	Φ_a
T	T	T	T	T	T	T	T	T
T	T	F	T	F	F	F	F	T
T	F	T	T	T	T	T	T	T
T	F	F	T	F	F	T	F	T
F	T	T	T	T	T	T	T	T
F	T	F	T	F	T	F	F	T
F	F	T	F	T	T	T	T	T
F	F	F	F	T	T	T	T	T

□

(b) $(p \rightarrow (q \rightarrow r)) \rightarrow (p \wedge q \rightarrow r)$

Solution. Let Φ_b denote $(p \rightarrow (q \rightarrow r)) \rightarrow (p \wedge q \rightarrow r)$.

p	q	r	$q \rightarrow r$	$p \rightarrow (q \rightarrow r)$	$p \wedge q$	$p \wedge q \rightarrow r$	Φ_b
T	T	T	T	T	T	T	T
T	T	F	F	F	T	F	T
T	F	T	T	T	F	T	T
T	F	F	T	T	F	T	T
F	T	T	T	T	F	T	T
F	T	F	F	T	F	T	T
F	F	T	T	T	F	T	T
F	F	F	T	T	F	T	T

□

(c) $((p \rightarrow q) \rightarrow p) \rightarrow p$

Solution. Let Φ_c denote $((p \rightarrow q) \rightarrow p) \rightarrow p$.

p	q	$p \rightarrow q$	$(p \rightarrow q) \rightarrow p$	Φ_c
T	T	T	T	T
T	F	F	T	T
F	T	T	F	T
F	F	T	F	T

□

2. Prove, using
- Natural Deduction*
- , the validity of the following sequents: (30 points)

(a) $(p \rightarrow r) \wedge (q \rightarrow r) \vdash p \vee q \rightarrow r$

Solution.

$$\frac{\frac{\frac{\frac{}{(p \rightarrow r) \wedge (q \rightarrow r), p \vee q \vdash p \vee q} (Hyp)}{\frac{}{(p \rightarrow r) \wedge (q \rightarrow r), p \vee q \vdash r} (\rightarrow I)}{\frac{}{(p \rightarrow r) \wedge (q \rightarrow r) \vdash p \vee q \rightarrow r} (\rightarrow I)} \alpha \quad \beta (\vee E)}{\frac{}{(p \rightarrow r) \wedge (q \rightarrow r), p \vee q, p \vdash (p \rightarrow r) \wedge (q \rightarrow r)} (\wedge E_1) \quad \frac{}{(p \rightarrow r) \wedge (q \rightarrow r), p \vee q, p \vdash p} (Hyp)}{\frac{}{(p \rightarrow r) \wedge (q \rightarrow r), p \vee q, p \vdash r} (\rightarrow E)} (\rightarrow E)$$

α :

$$\frac{\frac{\frac{}{(p \rightarrow r) \wedge (q \rightarrow r), p \vee q, p \vdash (p \rightarrow r) \wedge (q \rightarrow r)} (Hyp)}{\frac{}{(p \rightarrow r) \wedge (q \rightarrow r), p \vee q, p \vdash p \rightarrow r} (\wedge E_1) \quad \frac{}{(p \rightarrow r) \wedge (q \rightarrow r), p \vee q, p \vdash p} (Hyp)}{\frac{}{(p \rightarrow r) \wedge (q \rightarrow r), p \vee q, p \vdash r} (\rightarrow E)} (\rightarrow E)$$

β :

$$\frac{\frac{\frac{}{(p \rightarrow r) \wedge (q \rightarrow r), p \vee q, p \vdash (p \rightarrow r) \wedge (q \rightarrow r)} (Hyp)}{\frac{}{(p \rightarrow r) \wedge (q \rightarrow r), p \vee q, q \vdash q \rightarrow r} (\wedge E_2) \quad \frac{}{(p \rightarrow r) \wedge (q \rightarrow r), p \vee q, q \vdash q} (Hyp)}{\frac{}{(p \rightarrow r) \wedge (q \rightarrow r), p \vee q, q \vdash r} (\rightarrow E)} (\rightarrow E)$$

□

(b) $p \wedge q \rightarrow r \vdash p \rightarrow (q \rightarrow r)$

Solution.

$$\frac{\frac{\frac{}{p \wedge q \rightarrow r, p, q \vdash p \wedge q \rightarrow r} (Hyp) \quad \frac{\frac{}{p \wedge q \rightarrow r, p, q \vdash p} (Hyp) \quad \frac{}{p \wedge q \rightarrow r, p, q \vdash q} (Hyp)}{\frac{}{p \wedge q \rightarrow r, p, q \vdash p \wedge q} (\wedge I)} (\rightarrow E) \quad \frac{\frac{}{p \wedge q \rightarrow r, p, q \vdash r} (\rightarrow I) \quad \frac{}{p \wedge q \rightarrow r, p \vdash q \rightarrow r} (\rightarrow I)}{\frac{}{p \wedge q \rightarrow r \vdash p \rightarrow (q \rightarrow r)} (\rightarrow I)} (\rightarrow I)$$

□

(c) $\vdash ((p \rightarrow q) \rightarrow p) \rightarrow p$

Solution. (Jinn-Shu Chang)

$$\frac{\frac{\frac{}{(p \rightarrow q) \rightarrow p, \neg p \vdash (p \rightarrow q) \rightarrow p} (Hyp)}{\frac{}{(p \rightarrow q) \rightarrow p, \neg p \vdash p} (\rightarrow E) \quad \frac{}{(p \rightarrow q) \rightarrow p, \neg p \vdash \neg p} (Hyp)}{\frac{}{(p \rightarrow q) \rightarrow p, \neg p \vdash p \wedge \neg p} (\wedge I)} (\wedge I) \quad \frac{\frac{}{(p \rightarrow q) \rightarrow p, \neg p \vdash p \wedge \neg p} (\wedge I) \quad \frac{}{(p \rightarrow q) \rightarrow p \vdash \neg \neg p} (\neg \neg E)}{\frac{}{(p \rightarrow q) \rightarrow p \vdash p} (\rightarrow I)} (\rightarrow I) \quad \frac{}{\vdash ((p \rightarrow q) \rightarrow p) \rightarrow p} (\rightarrow I)$$

α :

$$\frac{\frac{\frac{}{(p \rightarrow q) \rightarrow q, \neg p, p, \neg q \vdash p} (Hyp) \quad \frac{}{(p \rightarrow q) \rightarrow q, \neg p, p, \neg q \vdash \neg p} (Hyp)}{\frac{}{(p \rightarrow q) \rightarrow q, \neg p, p, \neg q \vdash p \wedge \neg p} (\wedge I)} (\wedge I) \quad \frac{\frac{}{(p \rightarrow q) \rightarrow p, \neg p, p, \neg q \vdash p \wedge \neg p} (\wedge I) \quad \frac{}{(p \rightarrow q) \rightarrow p, \neg p, p \vdash \neg \neg q} (\neg \neg E)}{\frac{}{(p \rightarrow q) \rightarrow p, \neg p, p \vdash q} (\rightarrow I)} (\rightarrow I) \quad \frac{}{(p \rightarrow q) \rightarrow p, \neg p \vdash p \rightarrow q} (\rightarrow I)$$

□

3. Express each of the following requirements as a first-order formula. (40 points)

- (a) R is a functional binary relation. (Note: a binary relation R is *functional* if no element is related by R to more than one elements.)

Solution.

$$\forall a, b, c (R(a, b) \wedge R(a, c) \rightarrow b = c)$$

□

- (b) In graph $G(V, E)$ a node, named u , has exactly two neighbors.

Solution.

$$\exists a, b \in V (E(u, a) \wedge E(u, b) \wedge \neg(a = b) \wedge \forall c \in V (E(u, c) \rightarrow c = a \vee c = b))$$

□

- (c) a is the smallest prime number larger than b .

Solution. Let div denote a binary predicate such that $div(x, y)$ is *true* if and only if x divides y .

$$\forall i (b < i < a \rightarrow \exists j (div(j, i) \wedge \neg(j = 1) \wedge \neg(j = i))) \wedge \forall k (div(k, a) \rightarrow k = 1 \vee k = a) \wedge (b < a)$$

□

- (d) Array A , with n elements, is cyclically sorted. (Note: an array is *cyclically sorted* if, by circular-shifting its elements, one can turn it into a sorted array.)

Solution. We assume that the n elements are indexed from 0 through $n - 1$.

$$\begin{aligned} \exists k (& (k = n - 1 \vee (0 \leq k < n - 1 \wedge A[n - 1] < A[0])) \\ & \wedge \forall i (0 \leq i < k \rightarrow A[i] \leq A[i + 1]) \\ & \wedge \forall i (k \leq i < n - 1 \rightarrow A[i] \leq A[i + 1]) \end{aligned})$$

Alternatively,

$$\exists k (\forall i (0 \leq i < n \rightarrow A[(k + i) \bmod n] \leq A[(k + i + 1) \bmod n]))$$

where “ $a \bmod n$ ” gives the remainder of a divided by n .

Note: in pure logic terms, an array may be modeled as a function, so $A[i]$ would be written as $A(i)$. □