

The Sequent Calculus: System LK

- Axioms

$$\frac{}{A \vdash A} (Axiom)$$

For convenience, you may extend this to the following:

$$\frac{}{\Gamma, A \vdash A, \Delta} (Axiom)$$

- Logical Rules (I)

$$\begin{array}{c} \frac{\Gamma, A \vdash \Delta}{\Gamma, A \wedge B \vdash \Delta} (\wedge : Left_1) \quad \frac{\Gamma \vdash A, \Delta \quad \Gamma \vdash B, \Delta}{\Gamma \vdash A \wedge B, \Delta} (\wedge : Right) \\ \frac{\Gamma, B \vdash \Delta}{\Gamma, A \wedge B \vdash \Delta} (\wedge : Left_2) \\ \\ \frac{\Gamma, A \vdash \Delta \quad \Gamma, B \vdash \Delta}{\Gamma, A \vee B \vdash \Delta} (\vee : Left) \quad \frac{\Gamma \vdash A, \Delta}{\Gamma \vdash A \vee B, \Delta} (\vee : Right_1) \\ \frac{\Gamma \vdash B, \Delta}{\Gamma \vdash A \vee B, \Delta} (\vee : Right_2) \end{array}$$

For convenience, you may take the following alternatives:

$$\begin{array}{cc} \frac{\Gamma, A, B \vdash \Delta}{\Gamma, A \wedge B \vdash \Delta} (\wedge : Left) & \frac{\Gamma \vdash A, \Delta \quad \Gamma \vdash B, \Delta}{\Gamma \vdash A \wedge B, \Delta} (\wedge : Right) \\ \\ \frac{\Gamma, A \vdash \Delta \quad \Gamma, B \vdash \Delta}{\Gamma, A \vee B \vdash \Delta} (\vee : Left) & \frac{\Gamma \vdash A, B, \Delta}{\Gamma \vdash A \vee B, \Delta} (\vee : Right) \end{array}$$

- Logical Rules (II)

$$\begin{array}{cc} \frac{\Gamma \vdash A, \Delta_1 \quad \Gamma, B \vdash \Delta_2}{\Gamma, A \rightarrow B \vdash \Delta_1, \Delta_2} (\rightarrow : Left) & \frac{\Gamma, A \vdash B, \Delta}{\Gamma \vdash A \rightarrow B, \Delta} (\rightarrow : Right) \\ \\ \frac{\Gamma \vdash A, \Delta}{\Gamma, \neg A \vdash \Delta} (\neg : Left) & \frac{\Gamma, A \vdash \Delta}{\Gamma \vdash \neg A, \Delta} (\neg : Right) \end{array}$$

- Structural Rules

$$\frac{\Gamma \vdash \Delta}{\Gamma, A \vdash \Delta} (\text{Weakening : Left})$$

$$\frac{\Gamma \vdash \Delta}{\Gamma \vdash A, \Delta} (\text{Weakening : Right})$$

$$\frac{\Gamma, A, A \vdash \Delta}{\Gamma, A \vdash \Delta} (\text{Contraction : Left})$$

$$\frac{\Gamma \vdash A, A, \Delta}{\Gamma \vdash A, \Delta} (\text{Contraction : Right})$$

$$\frac{\Gamma_1, A, B, \Gamma_2 \vdash \Delta}{\Gamma_1, B, A, \Gamma_2 \vdash \Delta} (\text{Exchange : Left})$$

$$\frac{\Gamma \vdash \Delta_1, A, B, \Delta_2}{\Gamma \vdash \Delta_1, B, A, \Delta_2} (\text{Exchange : Right})$$

For convenience, you may treat Γ , Δ , etc. as *sets*, A , B as $\{A\}$, $\{B\}$, and the comma (in “ Γ, A ” etc.) as set union, so that you can do without these structural rules; use this alternative with the extended notion of an axiom.

- The Cut Rule

$$\frac{\Gamma_1 \vdash A, \Delta_1 \quad \Gamma_2, A \vdash \Delta_2}{\Gamma_1, \Gamma_2 \vdash \Delta_1, \Delta_2} (\text{Cut})$$

- Quantifier Rules

$$\frac{\Gamma, A[t/x] \vdash \Delta}{\Gamma, \forall x A \vdash \Delta} (\forall : \text{Left})$$

$$\frac{\Gamma \vdash A[y/x], \Delta}{\Gamma \vdash \forall x A, \Delta} (\forall : \text{Right})$$

$$\frac{\Gamma, A[y/x] \vdash \Delta}{\Gamma, \exists x A \vdash \Delta} (\exists : \text{Left})$$

$$\frac{\Gamma \vdash A[t/x], \Delta}{\Gamma \vdash \exists x A, \Delta} (\exists : \text{Right})$$

In these quantifier rules, we assume that all substitutions are admissible, y is not free in A , and y does not occur free in the lower sequent.

Axioms for Equality (an extension for languages with $=$):

Let $t, s_1, \dots, s_n, t_1, \dots, t_n$ be arbitrary terms.

$$\overline{\Gamma \vdash t = t, \Delta}$$

For every n -ary function f ,

$$\overline{\Gamma, s_1 = t_1, \dots, s_n = t_n \vdash f(s_1, \dots, s_n) = f(t_1, \dots, t_n), \Delta}$$

For every n -ary predicate P (including $=$),

$$\overline{\Gamma, s_1 = t_1, \dots, s_n = t_n, P(s_1, \dots, s_n) \vdash P(t_1, \dots, t_n), \Delta}$$