

Proofs in the Sequent Calculus (System LK)

- Below is a proof of $p \rightarrow q \vdash \neg p \vee q$. Note that the converse, i.e., $\neg p \vee q \vdash p \rightarrow q$, is also provable, affirming the equivalence of $p \rightarrow q$ and $\neg p \vee q$.

$$\begin{array}{c}
 \frac{}{p \vdash p} \text{ (Axiom)} \\
 \frac{p \vdash p}{\vdash \neg p, p} (\neg : \text{Right}) \\
 \frac{\vdash \neg p, p}{\vdash \neg p \vee q, p} (\vee : \text{Right}_1) \\
 \frac{\vdash \neg p \vee q, p}{\vdash p, \neg p \vee q} (\text{Exchange} : \text{Right}) \\
 \frac{}{q \vdash q} \text{ (Axiom)} \\
 \frac{q \vdash q}{q \vdash \neg p \vee q} (\vee : \text{Right}_2) \\
 \frac{\vdash p, \neg p \vee q \quad q \vdash \neg p \vee q}{p \rightarrow q \vdash \neg p \vee q, \neg p \vee q} (\rightarrow : \text{Left}) \\
 \frac{p \rightarrow q \vdash \neg p \vee q, \neg p \vee q}{p \rightarrow q \vdash \neg p \vee q} (\text{Contraction} : \text{Right})
 \end{array}$$

- If we treat Γ and Δ in a sequent $\Gamma \vdash \Delta$ as *sets* of formulae, then we can do without the structural rules. The preceding proof can be simplified as follows.

$$\begin{array}{c}
 \frac{}{p \vdash p} \text{ (Axiom)} \\
 \frac{p \vdash p}{\vdash p, \neg p} (\neg : \text{Right}) \\
 \frac{\vdash p, \neg p}{\vdash p, \neg p \vee q} (\vee : \text{Right}_1) \\
 \frac{}{q \vdash q} \text{ (Axiom)} \\
 \frac{q \vdash q}{q \vdash \neg p \vee q} (\vee : \text{Right}_2) \\
 \frac{\vdash p, \neg p \vee q \quad q \vdash \neg p \vee q}{p \rightarrow q \vdash \neg p \vee q} (\rightarrow : \text{Left})
 \end{array}$$

- Below is a proof of $\vdash \exists x(A \wedge B) \rightarrow \exists xA \wedge \exists xB$. Its converse does not necessarily hold and hence is not provable.

$$\begin{array}{c}
 \frac{}{A[y/x] \vdash A[y/x]} \text{ (Axiom)} \\
 \frac{A[y/x] \vdash A[y/x]}{A[y/x] \wedge B[y/x] \vdash A[y/x]} (\wedge : \text{Left}_1) \\
 \frac{A[y/x] \wedge B[y/x] \vdash A[y/x]}{A[y/x] \wedge B[y/x] \vdash A[y/x] \{y \text{ as } t \text{ in } A[t/x]\}} (\exists : \text{Right}) \\
 \frac{A[y/x] \wedge B[y/x] \vdash A[y/x] \{y \text{ as } t \text{ in } A[t/x]\}}{A[y/x] \wedge B[y/x] \{= (A \wedge B)[y/x]\} \vdash \exists xA} (\exists : \text{Left}) \\
 \frac{\exists x(A \wedge B) \vdash \exists xA}{\exists x(A \wedge B) \vdash \exists xA} \text{ (analogous to the left)} \\
 \frac{\exists x(A \wedge B) \vdash \exists xA \quad \exists x(A \wedge B) \vdash \exists xB}{\exists x(A \wedge B) \vdash \exists xA \wedge \exists xB} (\wedge : \text{Right}) \\
 \frac{\exists x(A \wedge B) \vdash \exists xA \wedge \exists xB}{\vdash \exists x(A \wedge B) \rightarrow \exists xA \wedge \exists xB} (\rightarrow : \text{Right})
 \end{array}$$

- A proof of $\forall x(x \geq 0) \vdash \forall x(\forall y(x + y \geq 0))$:

$$\begin{array}{c}
 \frac{}{(w + z \geq 0) \{= (x \geq 0)[w + z/x]\} \vdash (w + z \geq 0)} \text{ (Axiom)} \\
 \frac{(w + z \geq 0) \{= (x \geq 0)[w + z/x]\} \vdash (w + z \geq 0)}{\forall x(x \geq 0) \vdash (w + z \geq 0) \{= (w + y \geq 0)[z/y]\}} (\forall : \text{Left}) \\
 \frac{\forall x(x \geq 0) \vdash (w + z \geq 0) \{= (w + y \geq 0)[z/y]\}}{\forall x(x \geq 0) \vdash \forall y(w + y \geq 0) \{= (\forall y(x + y \geq 0))[w/x]\}} (\forall : \text{Right}) \\
 \frac{\forall x(x \geq 0) \vdash \forall y(w + y \geq 0) \{= (\forall y(x + y \geq 0))[w/x]\}}{\forall x(x \geq 0) \vdash \forall x(\forall y(x + y \geq 0))} (\forall : \text{Right})
 \end{array}$$