

Suggested Solutions for Homework Assignment #4

We assume the binding powers of the logical connectives and the entailment symbol decrease in this order: \neg , $\{\forall, \exists\}$, $\{\wedge, \vee\}$, \rightarrow , \leftrightarrow , \vdash .

1. Prove that the following annotated program segments are correct:

(a) (10 points)

$\{true\}$
if $x < y$ **then** $x, y := y, x$ **fi**
 $\{x \geq y\}$

Solution.

$$\frac{\frac{\text{pred. calculus + algebra}}{true \wedge x < y \rightarrow y \geq x} \quad \frac{\{y \geq x\} x, y := y, x \{x \geq y\}}{\{true \wedge x < y\} x, y := y, x \{x \geq y\}} \text{ (Assign) (SP)}}{\{true\} \text{ if } x < y \text{ then } x, y := y, x \text{ fi } \{x \geq y\}} \text{ (If-Then) (pred. calculus + algebra)}$$

□

(b) (10 points)

$\{g = 0 \wedge p = n \wedge n \geq 1\}$
while $p \geq 2$ **do**
 $g, p := g + 1, p - 1$
od
 $\{g = n - 1\}$

Solution.

$$\frac{\frac{\text{pred. calculus + algebra}}{g = 0 \wedge p = n \wedge n = 1 \rightarrow p > 0 \wedge p + g = n} \quad \alpha \quad \frac{\text{pred. calculus + algebra}}{p > 0 \wedge p + g = n \wedge \neg(p \geq 2) \rightarrow g = n - 1}}{\{g = 0 \wedge p = n \wedge n = 1\} \text{ while } p \geq 2 \text{ do } g, p := g - 1, p + 1 \text{ od } \{g = n - 1\}} \text{ (Consequence)}$$

α :

$$\frac{\beta \quad \frac{\{p + 1 > 0 \wedge (p + 1) + (g - 1) = n\} g, p := g - 1, p + 1 \{p > 0 \wedge p + g = n\}}{\{p > 0 \wedge p + g = n \wedge p \geq 2\} g, p := g - 1, p + 1 \{p > 0 \wedge p + g = n\}} \text{ (Assign) (SP)}}{\{p > 0 \wedge p + g = n\} \text{ while } p \geq 2 \text{ do } g, p := g - 1, p + 1 \text{ od } \{p > 0 \wedge p + g = n \wedge \neg(p \geq 2)\}} \text{ (while)}$$

β :

$$\frac{\text{pred. calculus + algebra}}{p > 0 \wedge p + g = n \wedge p \geq 2 \rightarrow p + 1 > 0 \wedge (p + 1) + (g - 1) = n}$$

□

(c) (20 points) For this program, prove its total correctness.

$\{y > 0 \wedge (x \equiv m \pmod{y})\}$
while $x \geq y$ **do**
 $x := x - y$
od
 $\{(x \equiv m \pmod{y}) \wedge x < y\}$

Solution.

$$\frac{\alpha \quad \frac{\text{pred. calculus + algebra}}{y > 0 \wedge (x \equiv m \pmod{y}) \wedge \neg(x \geq y) \rightarrow (x \equiv m \pmod{y}) \wedge x < y}}{\{ y > 0 \wedge (x \equiv m \pmod{y}) \} \mathbf{while} \ x \geq y \ \mathbf{do} \ x := x - y \ \mathbf{od} \ \{ (x \equiv m \pmod{y}) \wedge x < y \}} \text{(SP)}$$

$\alpha :$

$$\frac{\beta \quad \gamma \quad \frac{\text{pred. calculus + algebra}}{y > 0 \wedge (x \equiv m \pmod{y}) \wedge x \geq y \rightarrow x \geq 0}}{\{ y > 0 \wedge (x \equiv m \pmod{y}) \}} \text{(while: simply total)}$$

$$\mathbf{while} \ x \geq y \ \mathbf{do} \ x := x - y \ \mathbf{od}$$

$$\{ y > 0 \wedge (x \equiv m \pmod{y}) \wedge \neg(x \geq y) \}$$

$\beta :$

$$\frac{\frac{\text{pred. calculus + algebra}}{y > 0 \wedge (x \equiv m \pmod{y}) \wedge x \geq y \rightarrow} \quad \frac{\text{(Assign)}}{\{ y > 0 \wedge ((x - y) \equiv m \pmod{y}) \}}}{\frac{y > 0 \wedge ((x - y) \equiv m \pmod{y}) \quad \{ y > 0 \wedge (x \equiv m \pmod{y}) \}}{\{ y > 0 \wedge (x \equiv m \pmod{y}) \wedge x \geq y \} \ x := x - y \ \{ y > 0 \wedge (x \equiv m \pmod{y}) \}} \text{(SP)}}$$

$\gamma :$

$$\frac{\frac{\text{pred. calculus + algebra}}{y > 0 \wedge (x \equiv m \pmod{y}) \wedge x \geq y \wedge x = Z \rightarrow x - y < Z} \quad \frac{\text{(Assign)}}{\{ x - y < Z \} \ x := x - y \ \{ x < Z \}}}{\frac{\{ y > 0 \wedge (x \equiv m \pmod{y}) \wedge x \geq y \wedge x = Z \} \ x := x - y \ \{ x < Z \}}{\{ y > 0 \wedge (x \equiv m \pmod{y}) \wedge x \geq y \wedge x = Z \} \ x := x - y \ \{ x < Z \}} \text{(SP)}}$$

□