

Suggested Solutions for Homework Assignment #2

We assume the binding powers of the logical connectives and the entailment symbol decrease in this order: \neg , $\{\forall, \exists\}$, $\{\wedge, \vee\}$, $\rightarrow, \leftrightarrow, \vdash$.

1. (30 points) Prove, using *Natural Deduction*, the validity of the following sequents:

$$(a) \forall x(P(x) \rightarrow Q(x)) \vdash \forall xP(x) \rightarrow \forall xQ(x)$$

Solution. Assume w does not occur free either in $P(x)$ or in $Q(x)$.

$$\alpha : \frac{\frac{\frac{\frac{\forall x(P(x) \rightarrow Q(x)), \forall xP(x) \vdash \forall xP(x)}{(\forall E)} \quad (\neg I)}{\forall x(P(x) \rightarrow Q(x)), \forall xP(x) \vdash P(w)} \quad (\neg E)}{\forall x(P(x) \rightarrow Q(x)), \forall xP(x) \vdash Q(w)} \quad (\forall I)}{\forall x(P(x) \rightarrow Q(x)), \forall xP(x) \vdash \forall xQ(x)} \quad (\neg I)$$

$\alpha :$

$$\frac{\frac{\forall x(P(x) \rightarrow Q(x)), \forall xP(x) \vdash \forall x(P(x) \rightarrow Q(x))}{(\forall E)}}{\forall x(P(x) \rightarrow Q(x)), \forall xP(x) \vdash P(w) \rightarrow Q(w)}$$

□

$$(b) \vdash \exists x \forall y P(x, y) \rightarrow \forall y \exists x P(x, y)$$

Solution. Assume both w and z do not occur free in $P(x, y)$.

$$\frac{\frac{\frac{\exists x \forall y P(x, y) \vdash \exists x \forall y P(x, y)}{(\exists E)} \quad \frac{\frac{\exists x \forall y P(x, y), \forall y P(z, y) \vdash \forall y P(z, y)}{(\forall E)} \quad (\neg I)}{\exists x \forall y P(x, y), \forall y P(z, y) \vdash P(z, w)} \quad (\neg E)}{\exists x \forall y P(x, y), \forall y P(z, y) \vdash \exists x P(x, w)} \quad (\exists I)}{\exists x \forall y P(x, y) \vdash \exists x P(x, w)} \quad (\forall I)$$

□

$$(c) \forall x(A(x) \rightarrow B) \vdash \exists x A(x) \rightarrow B, \text{ assuming } x \text{ does not occur free in } B.$$

Solution. Assume w does not occur free either in $A(x)$ or in B .

$$\frac{\frac{\forall x(A(x) \rightarrow B), \exists x A(x) \vdash \exists x A(x)}{(\exists E)} \quad \alpha}{\frac{\forall x(A(x) \rightarrow B), \exists x A(x) \vdash B}{\forall x(A(x) \rightarrow B) \vdash \exists x A(x) \rightarrow B}} \quad (\neg I)$$

$\alpha :$

$$\beta : \frac{\frac{\frac{\forall x(A(x) \rightarrow B), \exists x A(x), A(w) \vdash \forall x(A(x) \rightarrow B)}{(\forall E)} \quad (\neg I)}{\forall x(A(x) \rightarrow B), \exists x A(x), A(w) \vdash A(w) \rightarrow B} \quad (\neg E)}{\forall x(A(x) \rightarrow B), \exists x A(x), A(w) \vdash B}$$

$\beta :$

$$\frac{}{\forall x(A(x) \rightarrow B), \exists x A(x), A(w) \vdash A(w)} \text{ (Hyp)}$$

□

2. (20 points) Prove, using *Natural Deduction* for the first-order logic with equality ($=$), that $=$ is an equivalence relation between terms, i.e., the following are valid sequents, in addition to the obvious “ $\vdash t = t$ ” (Reflexivity), which follows from the $=$ -Introduction rule.

(a) $t_2 = t_1 \vdash t_1 = t_2$ (Symmetry)

Solution.

$$\frac{\frac{}{t_2 = t_1 \vdash t_2 = t_1} \text{ (Hyp)} \quad \frac{}{t_2 = t_1 \vdash t_2 = t_2} \text{ (= I)}}{t_2 = t_1 \vdash t_1 = t_2} \text{ (= E)}$$

□

(b) $t_1 = t_2, t_2 = t_3 \vdash t_1 = t_3$ (Transitivity)

Solution.

$$\frac{\frac{}{t_1 = t_2, t_2 = t_3 \vdash t_2 = t_3} \text{ (Hyp)} \quad \frac{}{t_1 = t_2, t_2 = t_3 \vdash t_1 = t_2} \text{ (= E)}}{t_1 = t_2, t_2 = t_3 \vdash t_1 = t_3} \text{ (Hyp)}$$

□

3. (20 points) Taking the preceding valid sequents as axioms, prove using *Natural Deduction* the following derived rules for equality.

(a) $\frac{\Gamma \vdash t_2 = t_1}{\Gamma \vdash t_1 = t_2} \text{ (= Symmetry)}$

Solution.

$$\frac{\frac{\frac{}{\Gamma, t_2 = t_1 \vdash t_1 = t_2} \text{ (Axiom(Symmetry))}}{\Gamma \vdash t_2 = t_1 \rightarrow t_1 = t_2} \text{ (→I)}}{\Gamma \vdash t_1 = t_2} \text{ (→E)}$$

□

(b) $\frac{\Gamma \vdash t_1 = t_2 \quad \Gamma \vdash t_2 = t_3}{\Gamma \vdash t_1 = t_3} \text{ (= Transitivity)}$

Solution.

$$\frac{\frac{\alpha \quad \Gamma \vdash t_1 = t_2}{\Gamma \vdash t_2 = t_3 \rightarrow t_1 = t_3} \text{ (→E)}}{\Gamma \vdash t_1 = t_3} \text{ (→E)}$$

$\alpha :$

$$\frac{\frac{\frac{}{\Gamma, t_1 = t_2, t_2 = t_3 \vdash t_1 = t_3} \text{ (Axiom(Transitivity))}}{\frac{\frac{}{\Gamma, t_1 = t_2 \vdash t_2 = t_3 \rightarrow t_1 = t_3} \text{ (→I)}}{\Gamma \vdash t_1 = t_2 \rightarrow (t_2 = t_3 \rightarrow t_1 = t_3)} \text{ (→I)}} \text{ (→I)}$$

□

4. (30 points) A first-order theory for *groups* contains the following three axioms:

- $\forall a \forall b \forall c (a \cdot (b \cdot c) = (a \cdot b) \cdot c)$. (Associativity)
- $\forall a ((a \cdot e = a) \wedge (e \cdot a = a))$. (Identity)

- $\forall a((a \cdot a^{-1} = e) \wedge (a^{-1} \cdot a = e))$. (Inverse)

Here \cdot is the binary operation, e is a constant, called the identity, and $(\cdot)^{-1}$ is the inverse function which gives the inverse of an element. Let M denote the set of the three axioms. Prove, using *Natural Deduction* plus the derived rules in the preceding problem, the validity of the following sequents:

- (a) $M \vdash \forall a \forall b \forall c((a \cdot b = a \cdot c) \rightarrow b = c)$. (Hint: a typical proof in algebra books is the following: $b = e \cdot b = (a^{-1} \cdot a) \cdot b = a^{-1} \cdot (a \cdot b) = a^{-1} \cdot (a \cdot c) = (a^{-1} \cdot a) \cdot c = e \cdot c = c$.)
Solution.

$$\frac{\alpha \quad \delta}{M, x \cdot y = x \cdot z \vdash y = z} (= E) \\ \frac{}{M \vdash (x \cdot y = x \cdot z) \rightarrow y = z} (\rightarrow I) \\ \frac{}{M \vdash \forall c((x \cdot y = x \cdot c) \rightarrow y = c)} (\forall I) \\ \frac{}{M \vdash \forall b \forall c((x \cdot b = x \cdot c) \rightarrow b = c)} (\forall I) \\ \frac{}{M \vdash \forall a \forall b \forall c((a \cdot b = a \cdot c) \rightarrow b = c)} (\forall I)$$

$\alpha :$

$$\frac{\beta \quad \gamma}{M, x \cdot y = x \cdot z \vdash (x^{-1} \cdot x) \cdot y = y} (= E) \quad \frac{\begin{array}{c} M, x \cdot y = x \cdot z \vdash \forall a \forall b \forall c(a \cdot (b \cdot c) = (a \cdot b) \cdot c) \\ M, x \cdot y = x \cdot z \vdash \forall b \forall c(x^{-1} \cdot (b \cdot c) = (x^{-1} \cdot b) \cdot c) \end{array}}{\begin{array}{c} M, x \cdot y = x \cdot z \vdash \forall c(x^{-1} \cdot (x \cdot c) = (x^{-1} \cdot x) \cdot c) \\ M, x \cdot y = x \cdot z \vdash x^{-1} \cdot (x \cdot y) = (x^{-1} \cdot x) \cdot y \end{array}} \begin{array}{l} (= Hyp) \\ (\forall E) \\ (\forall E) \\ (\forall E) \\ (= E) \end{array}$$

$$M, x \cdot y = x \cdot z \vdash x^{-1} \cdot (x \cdot y) = y$$

$\beta :$

$$\frac{}{M, x \cdot y = x \cdot z \vdash \forall a(a \cdot a^{-1} = e \wedge a^{-1} \cdot a = e)} (= Hyp) \\ \frac{}{M, x \cdot y = x \cdot z \vdash x \cdot x^{-1} = e \wedge x^{-1} \cdot x = e} (\forall E) \\ \frac{}{M, x \cdot y = x \cdot z \vdash x^{-1} \cdot x = e} (\wedge E_2) \\ M, x \cdot y = x \cdot z \vdash e = x^{-1} \cdot x (= Symmetry)$$

$\gamma :$

$$\frac{}{M, x \cdot y = x \cdot z \vdash \forall a(a \cdot e = a \wedge e \cdot a = a)} (= Hyp) \\ \frac{}{M, x \cdot y = x \cdot z \vdash y \cdot e = y \wedge e \cdot y = y} (\forall E) \\ M, x \cdot y = x \cdot z \vdash e \cdot y = y (\wedge E_2)$$

$\delta :$

$$\frac{\begin{array}{c} M, x \cdot y = x \cdot z \vdash x \cdot y = x \cdot z \\ M, x \cdot y = x \cdot z \vdash x \cdot z = x \cdot y \end{array}}{M, x \cdot y = x \cdot z \vdash x^{-1} \cdot (x \cdot y) = z} (= Symmetry) \quad \frac{\text{the proof tree is similar to } \alpha}{M, x \cdot y = x \cdot z \vdash x^{-1} \cdot (x \cdot z) = z} (= E)$$

□

- (b) $M \vdash \forall a \forall b \forall c(((a \cdot b = e) \wedge (b \cdot a = e) \wedge (a \cdot c = e) \wedge (c \cdot a = e)) \rightarrow b = c)$, which says that every element has a unique inverse. (Hint: a typical proof in algebra books is the following: $b = b \cdot e = b \cdot (a \cdot c) = (b \cdot a) \cdot c = e \cdot c = c$.)

Solution. We use N to denote $x \cdot y = e \wedge y \cdot x = e \wedge x \cdot z = e \wedge z \cdot x = e$.

$$\begin{array}{c}
\frac{(1)\alpha \quad (1)\delta}{M, N, x \cdot y = x \cdot z \vdash y = z \stackrel{(\equiv E)}{\quad} M, N \vdash x \cdot y = x \cdot z \rightarrow y = z \stackrel{(\rightarrow I)}{\quad} \frac{\alpha \quad \beta}{M, N \vdash x \cdot y = x \cdot z \stackrel{(\equiv E)}{\quad} M, N \vdash y = z \stackrel{(\rightarrow I)}{\quad} M \vdash (x \cdot y = e \wedge y \cdot x = e \wedge x \cdot z = e \wedge z \cdot x = e) \rightarrow y = z \stackrel{(\rightarrow I)}{\quad} M \vdash \forall c((x \cdot y = e \wedge y \cdot x = e \wedge x \cdot c = e \wedge c \cdot x = e) \rightarrow y = c) \stackrel{(\forall I)}{\quad} M \vdash \forall b \forall c((x \cdot b = e \wedge b \cdot x = e \wedge x \cdot c = e \wedge c \cdot x = e) \rightarrow b = c) \stackrel{(\forall I)}{\quad} M \vdash \forall a \forall b \forall c((a \cdot b = e \wedge b \cdot a = e \wedge a \cdot c = e \wedge c \cdot a = e) \rightarrow b = c) \stackrel{(\forall I)}{\quad}} \\
\alpha : \\
\frac{}{M, N \vdash x \cdot y = e \wedge y \cdot x = e \wedge x \cdot z = e \wedge z \cdot x = e \stackrel{(\text{Hyp})}{\quad} M, N \vdash x \cdot z = e \wedge z \cdot x = e \stackrel{(\wedge E_2)}{\quad} M, N \vdash x \cdot z = e \stackrel{(\wedge E_1)}{\quad} M, N \vdash e = x \cdot z \stackrel{(\equiv Symmetry)}{\quad}} \\
\beta : \\
\frac{}{M, N \vdash x \cdot y = e \wedge y \cdot x = e \wedge x \cdot z = e \wedge z \cdot x = e \stackrel{(\text{Hyp})}{\quad} M, N \vdash x \cdot y = e \stackrel{(\wedge E_1)}{\quad}}
\end{array}$$

□