

## Homework Assignment #4

**Due Time/Date**

This assignment is due 2:10PM Tuesday, April 7, 2020. Late submission will be penalized by 20% for each working day overdue.

**How to Submit**

Please use a word processor or scan hand-written answers to produce a single PDF file. Name your file according to this pattern: "b057050xx-hw4". Upload the PDF file to the Ceiba course site for Theory of Computing 2020: <https://ceiba.ntu.edu.tw/1082theory2020>. You may discuss the problems with others, but copying answers is strictly forbidden.

**Problems**

(Note: problems marked with "Exercise X.XX" or "Problem X.XX" are taken from [Sipser 2013] with probable adaptation.)

1. (Problem 1.32; 20 points) For languages  $A$  and  $B$ , let the *shuffle* of  $A$  and  $B$  be the language  $\{w \mid w = a_1b_1 \cdots a_kb_k, \text{ where } a_1 \cdots a_k \in A \text{ and } b_1 \cdots b_k \in B, \text{ each } a_i, b_i \in \Sigma^*\}$ . Show that the class of regular languages is closed under shuffle.
2. (Problem 1.38; 20 points) Let

$$\Sigma_2 = \left\{ \begin{bmatrix} 0 \\ 0 \end{bmatrix}, \begin{bmatrix} 0 \\ 1 \end{bmatrix}, \begin{bmatrix} 1 \\ 0 \end{bmatrix}, \begin{bmatrix} 1 \\ 1 \end{bmatrix} \right\}.$$

Here,  $\Sigma_2$  contains all columns of 0s and 1s of length two. A string of symbols in  $\Sigma_2$  gives two rows of 0s and 1s. Consider each row to be a binary number and let

$$C = \{w \in \Sigma_2^* \mid \text{the bottom row of } w \text{ is three times the top row}\}.$$

For example,  $\begin{bmatrix} 0 \\ 0 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \end{bmatrix} \begin{bmatrix} 0 \\ 0 \end{bmatrix} \in C$ , but  $\begin{bmatrix} 0 \\ 1 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} \notin C$ . Show that  $C$  is regular. (You may assume the result claimed in Problem 5 of HW#2.)

3. (Problem 1.40; 10 points) Let  $\Sigma_2$  be the same as in Problem 2. Consider the top and bottom rows to be strings of 0s and 1s and let

$$E = \{w \in \Sigma_2^* \mid \text{the bottom row of } w \text{ is the reverse of the top row of } w\}.$$

Show that  $E$  is not regular.

4. (Problem 1.42; 20 points) Let  $C_n = \{x \mid x \text{ is a binary number that is a multiple of } n\}$ . Show that for each  $n \geq 1$ , the language  $C_n$  is regular.

5. (Problem 1.51; 10 points) Prove that the language  $\{w \in \{0, 1\}^* \mid w \text{ is not a palindrome}\}$  is not regular. You may use the pumping lemma and the closedness of the class of regular languages under union, intersection, and complement. (Note: a *palindrome* is a string that reads the same forward and backward.)
6. (Problem 1.66; 20 points) Let  $M = (Q, \Sigma, \delta, q_0, F)$  be a DFA and let  $h$  be a state of  $M$  called its “home”. A *synchronizing sequence* for  $M$  and  $h$  is a string  $s \in \Sigma^*$  where  $\delta(q, s) = h$  for every  $q \in Q$ . Say that  $M$  is *synchronizable* if it has a synchronizing sequence for some state  $h$ . Prove that, if  $M$  is a  $k$ -state synchronizable DFA, then it has a synchronizing sequence of length at most  $k^3$ . (Note:  $\delta(q, s)$  equals the state where  $M$  ends up, when  $M$  starts from state  $q$  and reads input  $s$ .)