## Homework Assignment \#3

## Note

This assignment is due 2:10PM Monday, March 19, 2012. Please write or type your answers on A4 (or similar size) paper. Drop your homework by the due time in Yih-Kuen Tsay's mail box on the first floor of Management College Building II. Late submission will be penalized by $20 \%$ for each working day overdue. You may discuss the problems with others, but copying answers is strictly forbidden.

## Problems

There are five problems in this assignment, each accounting for 20 points. (Note: problems marked with "(X.XX)" are taken from [Manber 1989] with probable adaptation.)

1. (3.4) Below is a theorem from Manber's book:

For all constants $c>0$ and $a>1$, and for all monotonically increasing functions $f(n)$, we have $(f(n))^{c}=O\left(a^{f(n)}\right)$.

Prove, by using the above theorem, that for all constants $a, b>0,\left(\log _{2} n\right)^{a}=O\left(n^{b}\right)$.
2. (3.5) For each of the following pairs of functions, say whether $f(n)=O(g(n))$ and/or $f(n)=\Omega(g(n))$. Justify your answers.

|  | $f(n)$ | $g(n)$ |
| :--- | :--- | :--- |
| (a) | $(\log n)^{\log n}$ | $\frac{n}{\log n}$ |
| (b) | $\sqrt{n}$ | $(\log n)^{2}$ |

3. (3.12) Solve the following recurrence relation:

$$
T(n)=n+\sum_{i=1}^{n-1} T(i),
$$

where $T(1)=1$.
4. Solve the following recurrence relation using generating functions. This is a very simple recurrence relation, but you must use generating functions in your solution.

$$
\left\{\begin{array}{l}
T(1)=2 \\
T(2)=4 \\
T(n)=3 T(n-1)-2 T(n-2), \quad n \geq 3
\end{array}\right.
$$

5. (3.30) Use Equation 1, shown below, to prove that $S(n)=\sum_{i=1}^{n}\left\lceil\log _{2}(n / i)\right\rceil$ satisfies $S(n)=O(n)$.

Bounding a summation by an integral
If $f(x)$ is a monotonically increasing continuous function, then

$$
\begin{equation*}
\sum_{i=1}^{n} f(i) \leq \int_{x=1}^{x=n+1} f(x) d x \tag{1}
\end{equation*}
$$

