## Homework Assignment #3

## Note

This assignment is due 2:10PM Tuesday, March 22, 2016. Please write or type your answers on A4 (or similar size) paper. Drop your homework by the due time in Yih-Kuen Tsay's mail box on the first floor of Management College Building II. Late submission will be penalized by 20% for each working day overdue. You may discuss the problems with others, but copying answers is strictly forbidden.

## **Problems**

There are five problems in this assignment, each accounting for 20 points. (Note: problems marked with "(X.XX)" are taken from [Manber 1989] with probable adaptation.)

1. (3.4) Below is a theorem from Manber's book:

For all constants c > 0 and a > 1, and for all monotonically increasing functions f(n), we have  $(f(n))^c = O(a^{f(n)})$ .

Prove, by using the above theorem, that for all constants a, b > 0,  $(\log_2 n)^a = O(n^b)$ .

2. (3.5) For each of the following pairs of functions, say whether f(n) = O(g(n)) and/or  $f(n) = \Omega(g(n))$ . Justify your answers.

$$\frac{f(n) \quad g(n)}{\text{(a)} \quad \frac{n^2}{\log n} \quad n(\log n)^2}$$
(b) 
$$n^3 2^n \quad 3^n$$

3. Solve the following recurrence relation using *generating functions*. This is a very simple recurrence relation, but you must use generating functions in your solution.

$$\begin{cases} T(1) = 1 \\ T(2) = 3 \\ T(n) = 3T(n-1) - 2T(n-2), & n \ge 3 \end{cases}$$

4. (3.26) Find the asymptotic behavior of the function T(n) defined by the recurrence relation

$$T(n) = T(n/2) + \sqrt{n}, T(1) = 1.$$

You can consider only values of n that are powers of 2.

5. (3.30) Use Equation 1, shown below, to prove that  $S(n) = \sum_{i=1}^{n} \lceil \log_2(n/i) \rceil$  satisfies S(n) = O(n).

## Bounding a summation by an integral

If f(x) is a monotonically increasing continuous function, then

$$\sum_{i=1}^{n} f(i) \le \int_{x=1}^{x=n+1} f(x) dx. \tag{1}$$