Analysis of Algorithms
(Based on [Manber 1989])

Yih-Kuen Tsay

Department of Information Management
National Taiwan University

Introduction

- The purpose of algorithm analysis is to predict the behavior (running time, space requirement, etc.) of an algorithm *without implementing it* on a specific computer. (Why?)

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- The purpose of algorithm analysis is to predict the behavior (running time, space requirement, etc.) of an algorithm *without implementing it* on a specific computer. (Why?)
- As the exact behavior of an algorithm is hard to predict, the analysis is usually an *approximation*:
  - Relative to the input size (usually denoted by n): input possibilities too enormous to elaborate
  - Asymptotic: should care more about larger inputs
  - Worst-Case: easier to do, often representative (Why not average-case?)
- Such an approximation is usually good enough for comparing different algorithms for the same problem.

Complexity

- Theoretically, "complexity of an algorithm" is a more precise term for "approximate behavior of an algorithm".
- Two most important measures of complexity:
  - Time Complexity
    an upper bound on the number of steps that the algorithm performs.
  - Space Complexity
    an upper bound on the amount of temporary storage required for running the algorithm (excluding the input, the output, and the program itself).
- We will focus on time complexity.

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  1. $100n$
  2. $2n^2 + 50$
  3. $100n^{1.8}$

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- We will study an approach (the O notation) that allows us to ignore constant factors and concentrate on the behavior as n goes to infinity.
- For most algorithms, the constants in the expressions of their running times tend to be small.

The O Notation

- A function $g(n)$ is $O(f(n))$ for another function $f(n)$ if there exist constants c and N such that, for all $n \geq N$, $g(n) \leq cf(n)$.
- The function $g(n)$ may be substantially less than $cf(n)$; the O notation bounds it *only from above*.
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- The function $g(n)$ may be substantially less than $cf(n)$; the O notation bounds it *only from above*.
- The O notation allows us to ignore constants conveniently.
- Examples:
  - $5n^2 + 15 = O(n^2)$.
    (cf. $5n^2 + 15 \leq O(n^2)$ or $5n^2 + 15 \in O(n^2)$)
  - $5n^2 + 15 = O(n^3)$.
    (cf. $5n^2 + 15 \leq O(n^3)$ or $5n^2 + 15 \in O(n^3)$)
  - In an expression, $T(n) = 3n^2 + O(n)$.

The O Notation (cont.)

- No need to specify the base of a logarithm.
  - $\log_2 n = \frac{\log_{10} n}{\log_{10} 2} = \frac{1}{\log_{10} 2} \log_{10} n.$
  - For example, we can just write $O(\log n)$.
- $O(1)$ denotes a constant.

Properties of O

- We can add and multiply with O.

Lemma (3.2)

1. If $f(n) = O(s(n))$ and $g(n) = O(r(n))$, then $f(n) + g(n) = O(s(n) + r(n))$.
2. If $f(n) = O(s(n))$ and $g(n) = O(r(n))$, then $f(n) \cdot g(n) = O(s(n) \cdot r(n))$.

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- However, we cannot subtract or divide with O. (Why?)
  - $2n = O(n)$, $n = O(n)$, and $2n - n = n \neq O(n - n)$.
  - $n^2 = O(n^2)$, $n = O(n^2)$, and $n^2/n = n \neq O(n^2/n^2)$.

Polynomial vs. Exponential

- A function $f(n)$ is *monotonically growing* if $n_1 \geq n_2$ implies that $f(n_1) \geq f(n_2)$.
- An exponential function grows *at least* as fast as a polynomial function does.

Theorem (3.1)

For all constants $c > 0$ and $a > 1$, and for all monotonically growing functions $f(n)$, $(f(n))^c = O(a^{f(n)})$.

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- Examples:
  - Take n as $f(n)$, $n^c = O(a^n)$.
  - Take $\log_a n$ as $f(n)$, $(\log_a n)^c = O(a^{\log_a n}) = O(n)$.

Speed of Growth

$\log n$	n	$n \log n$	n^2	n^3	2^n
0	1	0	1	1	2
1	2	2	4	8	4
2	4	8	16	64	16
3	8	24	64	512	256
4	16	64	256	4,096	65,536
5	32	160	1,024	32,768	4,294,967,296

Table 1.7 Function values

Source: [E. Horowitz *et al.* 1998].

Speed of Growth (cont.)

running times	*time* 1 1000 steps/sec	*time* 2 2000 steps/sec	*time* 3 4000 steps/sec	*time* 4 8000 steps/sec
$\log_2 n$	0.010	0.005	0.003	0.001
n	1	0.5	0.25	0.125
$n \log_2 n$	10	5	2.5	1.25
$n^{1.5}$	32	16	8	4
n^2	1,000	500	250	125
n^3	1,000,000	500,000	250,000	125,000
1.1^n	10^{39}	10^{39}	10^{38}	10^{38}

Table 3.1 Running times (int seconds) under different assumptions (n = 1000)

Source: [Manber 1989].

O, o, Ω, and Θ

- Let $T(n)$ be the number of steps required to solve a given problem for input size n.
- We say that $T(n) = \Omega(g(n))$ or the problem has a lower bound of $\Omega(g(n))$ if there exist constants c and N such that, for all $n \geq N$, $T(n) \geq cg(n)$.
- If a certain function $f(n)$ satisfies both $f(n) = O(g(n))$ and $f(n) = \Omega(g(n))$, then we say that $f(n) = \Theta(g(n))$.

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- We say that $f(n) = o(g(n))$ if $\lim_{n\to\infty} \frac{f(n)}{g(n)} = 0$.

Polynomial vs. Exponential (cont.)

- An exponential function grows *faster* than a polynomial function does.

Theorem (3.3)

For all constants $c > 0$ and $a > 1$, and for all monotonically growing functions $f(n)$, we have

$$(f(n))^c = o(a^{f(n)}).$$

- Consider a previous example again:
  Take $\log_a n$ as $f(n)$. For all $c > 0$ and $a > 1$,

$$(\log_a n)^c = o(a^{\log_a n}) = o(n).$$

Sums

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- Techniques for summing expressions are essential for complexity analysis.
- For example, given that we know

$$S_0(n) = \sum_{i=1}^{n} 1 = n$$

and

$$S_1(n) = \sum_{i=1}^{n} i = 1 + 2 + 3 + \cdots + n = \frac{n(n+1)}{2},$$

we want to compute the sum

$$S_2(n) = \sum_{i=1}^{n} i^2 = 1^2 + 2^2 + 3^2 + \cdots + n^2.$$

Sums (cont.)

From

$$(i+1)^3 = i^3 + 3i^2 + 3i + 1,$$

we have

$$(i+1)^3 - i^3 = 3i^2 + 3i + 1.$$

$$
\begin{array}{rcl}
2^3 - 1^3 & = & 3 \times 1^2 + 3 \times 1 + 1 \\
3^3 - 2^3 & = & 3 \times 2^2 + 3 \times 2 + 1 \\
4^3 - 3^3 & = & 3 \times 3^2 + 3 \times 3 + 1 \\
\cdots & \cdots & \cdots \\
(n+1)^3 - n^3 & = & 3 \times n^2 + 3 \times n + 1 \\
\hline
(n+1)^3 - 1 & = & 3 \times S_2(n) + 3 \times S_1(n) + S_0(n) \\
(S_3(n+1) - S_3(1)) - S_3(n) & = & 3 \times S_2(n) + 3 \times S_1(n) + S_0(n)
\end{array}
$$

Sums (cont.)

- So, we have

$$(n+1)^3 - 1 = 3 \times S_2(n) + 3 \times S_1(n) + S_0(n).$$

- Given $S_0(n)$ and $S_1(n)$, the sum $S_2(n)$ can be computed by straightforward algebra.
- Recall that the left-hand side $(n+1)^3 - 1$ equals $(S_3(n+1) - S_3(1)) - S_3(n)$, a result from "shifting and canceling" terms of two sums.

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- This generalizes to $S_k(n)$, for $k > 2$.
- Similar shifting and canceling techniques apply to other kinds of sums.

Recurrence Relations

- A *recurrence relation* is a way to define a function by an expression involving the same function.
- The Fibonacci numbers can be defined as follows:

$$\begin{cases} F(1) = 1 \\ F(2) = 1 \\ F(n) = F(n-2) + F(n-1) \end{cases}$$

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- It is more convenient to have an explicit (or closed-form) expression.
- To obtain the explicit expression is called *solving* the recurrence relation.

Guessing and Proving an Upper Bound

- Recurrence relation: $\begin{cases} T(2) = 1 \\ T(2n) \leq 2T(n) + 2n - 1 \end{cases}$
- Guess: $T(n) = O(n \log n)$.

Guessing and Proving an Upper Bound

- Recurrence relation: $\begin{cases} T(2) = 1 \\ T(2n) \leq 2T(n) + 2n - 1 \end{cases}$
- Guess: $T(n) = O(n \log n)$.
- Proof:
  1. Base case: $T(2) \leq 2 \log 2$.
  2. Inductive step: $$\begin{aligned} T(2n) &\leq 2T(n) + 2n - 1 \\ &\leq 2(n \log n) + 2n - 1 \\ &= 2n \log n + 2n \log 2 - 1 \\ &\leq 2n(\log n + \log 2) \\ &= 2n \log 2n \end{aligned}$$

Recurrent Relations with Full History

- Example:

$$T(n) = c + \sum_{i=1}^{n-1} T(i),$$

  where c is a constant and $T(1)$ is given separately.
- $T(n) - T(n-1) = (c + \sum_{i=1}^{n-1} T(i)) - (c + \sum_{i=1}^{n-2} T(i)) = T(n-1)$; hence, $T(n) = 2T(n-1)$. (This holds only for $n \geq 3$.)
- So, we get

$$\begin{cases} T(2) = c + T(1) \\ T(n) = 2T(n-1) & \text{if } n \geq 3 \end{cases}$$

  which is easier to solve.
- Other examples as a reading assignment ...

Divide and Conquer Relations

- The running time $T(n)$ of a divide-and-conquer algorithm satisfies

$$T(n) = aT(n/b) + cn^k$$

  where
  - a is the number of subproblems,
  - n/b is the size of each subproblem, and
  - cn^k is the running time of the solutions-combining algorithm.

Divide and Conquer Relations (cont.)

Assume, for simplicity, $n = b^m$ ($\frac{n}{b^m} = 1$, $\frac{n}{b^{m-1}} = b$, etc.).

$$\begin{aligned} T(n) &= aT(\tfrac{n}{b}) + cn^k \\ &= a(aT(\tfrac{n}{b^2}) + c(\tfrac{n}{b})^k) + cn^k \\ &= a(a(aT(\tfrac{n}{b^3}) + c(\tfrac{n}{b^2})^k) + c(\tfrac{n}{b})^k) + cn^k \\ &\cdots \\ &= a(a(\cdots(aT(\tfrac{n}{b^m}) + c(\tfrac{n}{b^{m-1}})^k) + \cdots) + c(\tfrac{n}{b})^k) + cn^k \end{aligned}$$

Assuming $T(1) = c$,

$$T(n) = c\sum_{i=0}^{m} a^{m-i} b^{ik} = ca^m \sum_{i=0}^{m} (\frac{b^k}{a})^i.$$

Three cases: $\frac{b^k}{a} < 1$, $\frac{b^k}{a} = 1$, and $\frac{b^k}{a} > 1$.

Divide and Conquer Relations (cont.)

Theorem (3.4)

The solution of the recurrence relation $T(n) = aT(n/b) + cn^k$*, where* a *and* b *are integer constants,* $a \geq 1$*,* $b \geq 2$*, and* c *and* k *are positive constants, is*

$$T(n) = \begin{cases} O(n^{\log_b a}) & \text{if } a > b^k \\ O(n^k \log n) & \text{if } a = b^k \\ O(n^k) & \text{if } a < b^k \end{cases}$$

Useful Facts

- Harmonic series

$$H_n = \sum_{k=1}^{n} \frac{1}{k} = \ln n + \gamma + O(1/n),$$

where $\gamma = 0.577\ldots$ is Euler's constant.

- Sum of logarithms

$$\begin{aligned} \textstyle\sum_{i=1}^{n} \lfloor \log_2 i \rfloor &= (n+1)\lfloor \log_2 n \rfloor - 2^{\lfloor \log_2 n \rfloor + 1} + 2 \\ &= \Theta(n \log n). \end{aligned}$$

Useful Facts (cont.)

- Bounding a summation by an integral:
  If $f(x)$ is monotonically *increasing*, then

$$\sum_{i=1}^{n} f(i) \leq \int_{1}^{n+1} f(x)dx.$$

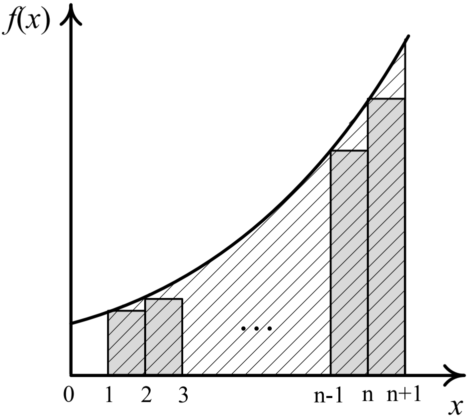
  If $f(x)$ is monotonically *decreasing*, then

$$\sum_{i=1}^{n} f(i) \leq f(1) + \int_{1}^{n} f(x)dx.$$

- Stirling's approximation

$$n! = \sqrt{2\pi n}\left(\frac{n}{e}\right)^n (1 + O(1/n)).$$

Bounding a Summation by an Integral



$$\sum_{i=1}^{n} f(i) \leq \int_{1}^{n+1} f(x)dx.$$