## Homework Assignment \#3

## Note

This assignment is due $2: 10 \mathrm{PM}$ Tuesday, March 21, 2017. Please write or type your answers on A4 (or similar size) paper. Drop your homework by the due time in Yih-Kuen Tsay's mail box on the first floor of Management College Building 2. Late submission will be penalized by $20 \%$ for each working day overdue. You may discuss the problems with others, but copying answers is strictly forbidden.

## Problems

There are five problems in this assignment, each accounting for 20 points. (Note: problems marked with "(X.XX)" are taken from [Manber 1989] with probable adaptation.)

1. Consider the following algorithm for computing the square of a given non-negative integer.
```
Algorithm mySquare( \(n\) );
begin
    // assume that \(n \geq 0\)
    \(x:=n\);
    \(y:=0\);
    while \(x>0\) do
        \(y:=y+2 x-1 ;\)
        \(x:=x-1\);
    od
    ...
end
```

Let $\operatorname{Inv}(n, x, y)$ denote the assertion:

$$
x \geq 0 \wedge y \geq 0 \wedge y=n^{2}-x^{2}
$$

Claim: $\operatorname{Inv}(n, x, y)$ is a loop invariant of the while loop, assuming that $n \geq 0$.
Prove the claim by induction. What does the loop invariant implicate when the while loop terminates?
2. (3.5) For each of the following pairs of functions, say whether $f(n)=O(g(n))$ and/or $f(n)=\Omega(g(n))$. Justify your answers.

|  | $f(n)$ | $g(n)$ |
| :--- | :--- | :--- |
| (a) $\frac{n}{\log n}$ | $(\log n)^{2}$ |  |
| (b) | $n^{4} 2^{n}$ | $4^{n}$ |

3. Solve the following recurrence relation using generating functions. This is a very simple recurrence relation, but for the purpose of practicing you must use generating functions in your solution.

$$
\left\{\begin{array}{l}
T(1)=1 \\
T(2)=3 \\
T(n)=2 T(n-1)-T(n-2), \quad n \geq 3
\end{array}\right.
$$

4. (3.26) Find the asymptotic behavior of the function $T(n)$ defined by the recurrence relation

$$
T(n)=T(n / 2)+\sqrt{n}, T(1)=1 .
$$

You can consider only values of $n$ that are powers of 2 .
5. (3.30) Use Equation 1, shown below, to prove that $S(n)=\sum_{i=1}^{n}\left\lceil\log _{2}(n / i)\right\rceil$ satisfies $S(n)=O(n)$.

## Bounding a summation by an integral

If $f(x)$ is a monotonically increasing continuous function, then

$$
\begin{equation*}
\sum_{i=1}^{n} f(i) \leq \int_{x=1}^{x=n+1} f(x) d x . \tag{1}
\end{equation*}
$$

