Reduction

(Based on [Manber 1989])

Yih-Kuen Tsay

Department of Information Management
National Taiwan University

Introduction

- The basic idea of *reduction* is to solve a problem with the solution to another "similar" problem.
- When Problem A can be reduced (efficiently) to Problem B, there are two consequences:
  - A solution to Problem B may be used to solve Problem A.
  - If A is known to be "hard", then B is also necessarily "hard".
- One should avoid the pitfall of reducing a problem to another that is too general or too hard.

Matching

- Given an undirected graph $G = (V, E)$, a **matching** is a set of edges that do not share a common vertex.
- A **maximum** matching is one with the maximum number of edges.
- A **maximal** matching is one that cannot be extended by adding any other edge.

Bipartite Matching

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- A bipartite graph $G = (V, E, U)$ is a graph with $V \cup U$ as the set of vertices and E as the set of edges such that
  - V and U are disjoint and
  - The edges in E connect vertices from V to vertices in U.

Problem

Given a bipartite graph $G = (V, E, U)$, find a maximum matching in G.

Networks

- Consider a directed graph, or network, $G = (V, E)$ with two distinguished vertices: s (the source) with indegree 0 and t (the sink) with outdegree 0.
- Each edge e in E has an associated positive weight $c(e)$, called the *capacity* of e.

The Network Flow Problem

- A **flow** is a function f on E that satisfies the following two conditions:
  1. $0 \leq f(e) \leq c(e)$.
  2. $\sum_{u} f(u,v) = \sum_{w} f(v,w)$, for all $v \in V - \{s,t\}$.
- The **network flow problem** is to maximize the flow f for a given network G.

Bipartite Matching to Network Flow

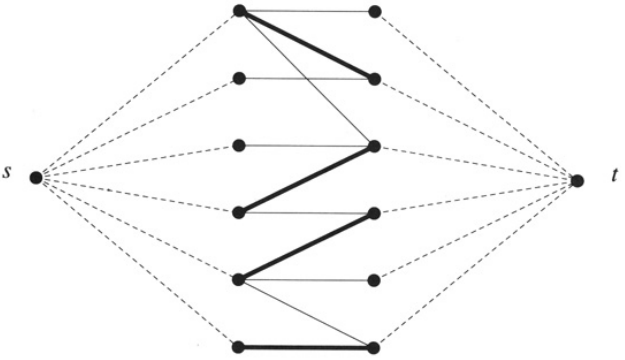

Figure 7.39 Reducing bipartite matching to network flow (the directions of all the edges are from left to right).

Source: [Manber 1989].

Bipartite Matching to Network Flow (cont.)

- Mapping from the input $G = (V, E, U)$ of the bipartite matching problem to the input $G' = (V', E')$ and c of the network flow problem:
  - The network is $G' = (V', E')$ where
    - $V' = \{s\} \cup V \cup U \cup \{t\}$
    - $E' = \{(s, v) \mid v \in V\} \cup E \cup \{(u, t) \mid u \in U\}$
  - The capacity for every $e \in E'$ is 1, i.e., $\forall e \in E', c(e) = 1$.
- Correspondence between the two solutions
  - A maximum flow f in G' defines a maximum matching M_f in G.
  - A maximum matching M in G induces a maximum flow f_M in G'.

Notations

- Let $\overline{v}$ denote a vector $(v_1, v_2, \ldots, v_n)$ of n constants or n variables.
- In the following, $\overline{a}$, $\overline{b}$, $\overline{c}$, and $\overline{e}$ are vectors of n constants.
- And, $\overline{x}$ and $\overline{y}$ are vectors of n variables.
- The (inner or dot) product $\overline{a} \cdot \overline{x}$ of two vectors $\overline{a}$ and $\overline{x}$ is defined as follows:

$$\overline{a} \cdot \overline{x} = \sum_{i=1}^{n} a_i \cdot x_i$$

Linear Programming

- Objective function:

$$\overline{c} \cdot \overline{x}$$

- Equality constraints:

$$\begin{array}{rcl} \overline{e}_1 \cdot \overline{x} & = & d_1 \\ \overline{e}_2 \cdot \overline{x} & = & d_2 \\ & \vdots & \\ \overline{e}_m \cdot \overline{x} & = & d_m \end{array}$$

- Inequality constraints may be turned into equality constraints by introducing *slack* variables.
- Non-negative constraints: $x_j \geq 0$, for all j in P, where P is a subset of $\{1, 2, \ldots, n\}$.
- The goal is to *maximize* (or *minimize*) the value of the objective function, subject to the equality constraints.

Network Flow to Linear Programming

- From the input $G = (V, E)$ and c of the network flow problem to the objective function and constraints of linear programming:
  - Let $x_1, x_2, \ldots, x_n$ represent the flow values of the n edges.
  - Objective function:
    $$\sum_{i \in S} x_i$$
    where S is the set of edges leaving the source.
  - Inequality constraints:
    $$x_i \leq c_i, \text{for all } i,\ 1 \leq i \leq n$$
    where c_i is the capacity of edge i.
  - Equality constraints:
    $$\sum_{i \text{ leaves } v} x_i - \sum_{j \text{ enters } v} x_j = 0, \text{ for every } v \in V \setminus \{s, t\}$$
  - Non-negative constraints: $x_i \geq 0, \text{for all } i,\ 1 \leq i \leq n$.