## Homework Assignment #8

## Due Time/Date

 $2:10\mathrm{PM}$  Tuesday, December 8, 2020. Late submission will be penalized by 20% for each working day overdue.

## Note

Please write or type your answers on A4 (or similar size) paper. Drop your homework by the due time in Yih-Kuen Tsay's mail box on the first floor of Management College Building 2. You may discuss the problems with others, but copying answers is strictly forbidden.

## Problems

There are five problems in this assignment, each accounting for 20 points. (Note: problems marked with "(X.XX)" are taken from [Manber 1989] with probable adaptation.)

- 1. (7.38) Given a directed acyclic graph G = (V, E), find a simple (directed) path in G that has the maximum number of edges among all simple paths in G. The algorithm should run in linear time.
- 2. Dijkstra's algorithm for single-source shortest paths assumes that every edge of the input graph has a nonnegative weight. Suppose we are given a graph with negative weights on some edges, where the minimum weight of the edges is -c for some c > 0. If we add c to the weight of every edge, then we obtain a new graph with nonnegative edge weights. We could then apply Dijkstra's algorithm to find the shortest paths for the new graph and thereafter subtract c from each edge of a path. Does this give us the shortest paths for the original graph? Please explain your answer.
- 3. (7.9) Prove that if the costs of all edges in a given connected graph are distinct, then the graph has exactly one unique minimum-cost spanning tree.
- 4. What is wrong with the following algorithm for computing the minimum-cost spanning tree of a given weighted undirected graph (assumed to be connected)?

If the input is just a single-node graph, return the single node. Otherwise, divide the graph into two subgraphs, recursively compute their minimum-cost spanning trees, and then connect the two spanning trees with an edge between the two subgraphs that has the minimum weight.

5. (7.61) Let G = (V, E) be a connected weighted undirected graph and T be a minimum-cost spanning tree (MCST) of G. Suppose that the cost of one edge  $\{u, v\}$  in G is changed (*increased* or *decreased*);  $\{u, v\}$  may or may not belong to T. Design an algorithm to either find a new MCST or to determine that T is still an MCST. The more efficient your algorithm is, the more points you will be credited for this problem. Explain why your algorithm is correct and analyze its time complexity.