## Midterm: Part I

## Note

This is a closed-book exam. Part I contains five problems, each accounting for 10 points.

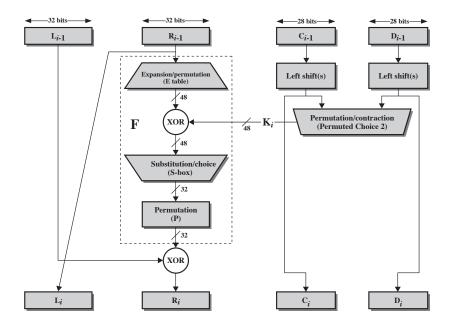
## Problems

- 1. Consider known-plaintext attacks on the  $3 \times 3$  Hill cipher. What are the criteria on the known plaintext-ciphertext pairs for an attack to succeed? Please be as precise as possible and explain the reasons for such criteria.
- 2. (a) Why in DES the round function  $(\mathbf{F})$  need not be invertible?
  - (b) How does three-key triple DES achieve backward compatibility with DES? Please describe all alternatives.
- 3. Consider the AES algorithm, where the irreducible polynomial modulus is  $x^8 + x^4 + x^3 + x + 1$ .
  - (a) What is the result of  $(0110\ 1101) \cdot (0000\ 0110)$ ? Show the steps of your calculation.
  - (b) What is the value of  $(0101\ 1010)^{-1}$ ? Show the steps of your calculation.
- 4. Using AES, decryption takes a slightly longer time than encryption.
  - (a) Which operation and its inverse are most responsible for this difference? Why does the inverse takes a longer time than the original operation?
  - (b) Why is this difference not reflected in the encryption and decryption with some modes of operation?
- 5. Consider pseudorandom number generation with the OFB mode of operation using 128-bit encryption. Suppose, as an observer (not knowing the seed value), you have observed so far n different blocks  $C_1, C_2, \ldots, C_n$  of pseudorandom bits on the output.
  - (a) If the next block  $C_{n+1}$  would be equal to any of the previous blocks, it must be  $C_1$ . Why?
  - (b) What is the probability that the stream of blocks will start to repeat itself from  $C_{n+1}$ ?

Please justify your answers.

## Appendix

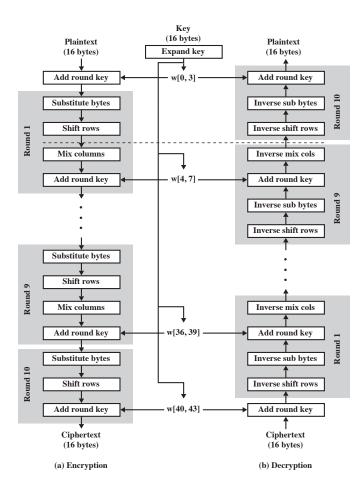
• Single round of the DES Algorithm:



• Extended Euclid's algorithm for polynomials:

EXTENDED EUCLID(a(x), b(x)):

- 1.  $[V_1(x), W_1(x), R_1(x)] \leftarrow [1, 0, a(x)]; [V_2(x), W_2(x), R_2(x)] \leftarrow [0, 1, b(x)]$
- 2. if  $R_2(x) = 0$  then return  $R_1(x) = \gcd(a(x), b(x))$ ; no inverse
- 3. if  $R_2(x) = 1$  then return  $R_2(x) = \gcd(a(x), b(x)); W_2(x) = b^{-1}(x) \pmod{a(x)}$
- 4. Q(x) = the quotient of  $R_1(x)/R_2(x)$
- 5. [V(x), W(x), R(x)]
  - $\leftarrow [V_1(x) Q(x)V_2(x), W_1(x) Q(x)W_2(x), R_1(x) Q(x)R_2(x)]$
- 6.  $[V_1(x), W_1(x), R_1(x)] \leftarrow [V_2(x), W_2(x), R_2(x)]$
- 7.  $[V_2(x), W_2(x), R_2(x)] \leftarrow [V(x), W(x), R(x)]$
- $8. \quad {\rm go to} \ 2$
- AES encryption and decryption:



• Pseudorandom number generation with the OFB mode:

