

## Suggested Solutions for Homework Assignment #1

We assume the binding powers of the logical connectives and the entailment symbol decrease in this order:  $\neg$ ,  $\{\wedge, \vee\}$ ,  $\rightarrow$ ,  $\leftrightarrow$ ,  $\vdash$ .

1. Prove, using Gentzen's System *LK*, the validity of the following sequents:

$$(a) \vdash (\neg p \vee q) \rightarrow (p \rightarrow q) \quad (10 \text{ points})$$

$$(b) p \vee q \rightarrow r \vdash (p \rightarrow r) \wedge (q \rightarrow r) \quad (10 \text{ points})$$

$$(c) \vdash (p \rightarrow (q \rightarrow r)) \rightarrow (p \wedge q \rightarrow r) \quad (10 \text{ points})$$

$$(d) \vdash ((p \rightarrow q) \rightarrow p) \rightarrow p \quad (10 \text{ points})$$

*Solution.* In the following, the right label *L* stands for *Left* and *R* stands for *Right*. We use the set semantics in the proofs of the following sequents.

(a)

$$\frac{\frac{\frac{\overline{p \vdash p, q} \text{ Axiom}}{\neg p, p \vdash q} \neg : L \quad \frac{\overline{q, p \vdash q} \text{ Axiom}}{\vee : L}}{\neg p \vee q, p \vdash q} \rightarrow : R}{\neg p \vee q \vdash p \rightarrow q} \rightarrow : R}{\vdash (\neg p \vee q) \rightarrow (p \rightarrow q)} \rightarrow : R$$

(b)

$$\frac{\frac{\frac{\overline{p \vdash p} \text{ Axiom}}{p \vdash p \vee q} \vee : R_1 \quad \frac{\overline{r, p \vdash r} \text{ Axiom}}{\rightarrow : L} \quad \frac{\frac{\overline{q \vdash q} \text{ Axiom}}{q \vdash p \vee q} \vee : R_2 \quad \frac{\overline{r, q \vdash r} \text{ Axiom}}{\rightarrow : L}}{p \vee q \rightarrow r, q \vdash r} \rightarrow : R}{p \vee q \rightarrow r, p \vdash r} \rightarrow : R}{p \vee q \rightarrow r \vdash p \rightarrow r} \rightarrow : R}{p \vee q \rightarrow r \vdash (p \rightarrow r) \wedge (q \rightarrow r)} \wedge : R$$

(c)

$$\frac{\frac{\overline{p, q \vdash p} \text{ Axiom} \quad \frac{\frac{\overline{p, q \vdash q} \text{ Axiom} \quad \frac{\overline{p, q, r \vdash r} \text{ Axiom}}{\rightarrow : L}}{p, q, q \rightarrow r \vdash r} \rightarrow : L}}{p \rightarrow (q \rightarrow r), p, q \vdash r} \wedge : L}{p \rightarrow (q \rightarrow r), p \wedge q \vdash r} \rightarrow : R}{p \rightarrow (q \rightarrow r) \vdash (p \wedge q) \rightarrow r} \rightarrow : R}{\vdash (p \rightarrow (q \rightarrow r)) \rightarrow ((p \wedge q) \rightarrow r)} \rightarrow : R$$

(d)

$$\frac{\frac{\frac{\overline{p \vdash p, q} \text{ Axiom}}{\vdash p \rightarrow q, p} \rightarrow : R \quad \frac{\overline{p \vdash p} \text{ Axiom}}{\rightarrow : L}}{(p \rightarrow q) \rightarrow p \vdash p} \rightarrow : R}{\vdash ((p \rightarrow q) \rightarrow p) \rightarrow p} \rightarrow : R$$

□

2. We presented in class the  $\wedge : Left_1$  and  $\wedge : Left_2$  rules of Gentzen's System  $LK$  as follows.

$$\frac{\Gamma, A \vdash \Delta}{\Gamma, A \wedge B \vdash \Delta} (\wedge : Left_1) \qquad \frac{\Gamma, B \vdash \Delta}{\Gamma, A \wedge B \vdash \Delta} (\wedge : Left_2)$$

We also suggested that, for convenience, the following  $\wedge : Left$  rule can be used instead:

$$\frac{\Gamma, A, B \vdash \Delta}{\Gamma, A \wedge B \vdash \Delta} (\wedge : Left)$$

Show that every proof using the  $\wedge : Left_1$  or  $\wedge : Left_2$  rules may be converted into one using the  $\wedge : Left$  rule and vice versa. (20 points)

*Solution.* In the following, the set semantics is used.

- Convert the proof using the  $\wedge : Left_1$  rule into one using  $\wedge : Left$  rule.

$$\text{Original: } \frac{\frac{\dots}{\Sigma, A, \vdash \Delta} \wedge Left_1}{\Sigma, A \wedge B \vdash \Delta} \qquad \text{Converted: } \frac{\frac{\dots}{\Sigma, A \vdash \Delta} \text{Weakening : Left}}{\Sigma, A, B \vdash \Delta} \wedge : Left$$

Convert the proof using the  $\wedge : Left_2$  rule into one using  $\wedge : Left$  rule.

$$\text{Original: } \frac{\frac{\dots}{\Sigma, B, \vdash \Delta} \wedge Left_2}{\Sigma, A \wedge B \vdash \Delta} \qquad \text{Converted: } \frac{\frac{\dots}{\Sigma, B \vdash \Delta} \text{Weakening : Left}}{\Sigma, A, B \vdash \Delta} \wedge : Left$$

- Convert the proof using the  $\wedge : Left$  rule into one using  $\wedge : Left_1$  and  $\wedge : Left_2$  rules.

$$\text{Original: } \frac{\frac{\dots}{\Sigma, A, B, \vdash \Delta} \wedge Left}{\Sigma, A \wedge B \vdash \Delta} \qquad \text{Converted: } \frac{\frac{\frac{\dots}{\Sigma, A, B \vdash \Delta} \wedge : Left_2}{\Sigma, A, A \wedge B \vdash \Delta} \wedge : Left_1}{\Sigma, A \wedge B, A \wedge B \vdash \Delta} \text{Contraction : Left}$$

□

3. Prove, using *Natural Deduction*, the validity of the following sequents:

- (a)  $\vdash (p \rightarrow q) \rightarrow (\neg p \vee q)$  (10 points)
- (b)  $(p \rightarrow r) \wedge (q \rightarrow r) \vdash p \vee q \rightarrow r$  (10 points)
- (c)  $\vdash (p \wedge q \rightarrow r) \rightarrow (p \rightarrow (q \rightarrow r))$  (10 points)
- (d)  $\vdash ((p \rightarrow q) \rightarrow p) \rightarrow p$  (10 points)

*Solution.*

(a)

$$\begin{array}{c}
\alpha \quad \frac{\frac{\frac{}{p \rightarrow q, \neg(\neg p \vee q), p \vdash \neg(\neg p \vee q)}{Hyp}}{p \rightarrow q, \neg(\neg p \vee q), p \vdash (\neg p \vee q) \wedge \neg(\neg p \vee q)}{\wedge I}}{p \rightarrow q, \neg(\neg p \vee q) \vdash \neg p}{\neg I}}{\frac{\frac{\frac{}{p \rightarrow q, \neg(\neg p \vee q) \vdash \neg p \vee q}}{\vee I_1}}{p \rightarrow q, \neg(\neg p \vee q) \vdash \neg(\neg p \vee q)}{Hyp}}{\frac{\frac{\frac{}{p \rightarrow q, \neg(\neg p \vee q) \vdash (\neg p \vee q) \wedge \neg(\neg p \vee q)}{\wedge I}}{p \rightarrow q \vdash \neg\neg(\neg p \vee q)}{\neg\neg E}}{p \rightarrow q \vdash \neg p \vee q}}{\rightarrow I}}{\vdash (p \rightarrow q) \rightarrow (\neg p \vee q)}{\rightarrow I}}
\end{array}$$

 $\alpha :$ 

$$\frac{\frac{\frac{}{p \rightarrow q, \neg(\neg p \vee q), p \vdash p \rightarrow q}}{Hyp} \quad \frac{\frac{}{p \rightarrow q, \neg(\neg p \vee q), p \vdash p}}{Hyp}}{\rightarrow E}}{\frac{\frac{\frac{}{p \rightarrow q, \neg(\neg p \vee q), p \vdash q}}{\vee I_2}}{p \rightarrow q, \neg(\neg p \vee q), p \vdash \neg p \vee q}}{\vee I_2}}$$

(b)

$$\frac{\frac{\frac{}{(p \rightarrow r) \wedge (q \rightarrow r), p \vee q \vdash p \vee q}}{Hyp} \quad \alpha \quad \beta}{\vee E}}{\frac{\frac{\frac{}{(p \rightarrow r) \wedge (q \rightarrow r), p \vee q \vdash r}}{\rightarrow I}}{(p \rightarrow r) \wedge (q \rightarrow r) \vdash p \vee q \rightarrow r}}{\rightarrow I}}$$

 $\alpha :$ 

$$\frac{\frac{\frac{\frac{}{(p \rightarrow r) \wedge (q \rightarrow r), p \vee q, p \vdash (p \rightarrow r) \wedge (q \rightarrow r)}}{Hyp}}{\frac{\frac{}{(p \rightarrow r) \wedge (q \rightarrow r), p \vee q, p \vdash p \rightarrow r}}{\wedge E_1}}}{\frac{\frac{}{(p \rightarrow r) \wedge (q \rightarrow r), p \vee q, p \vdash p}}{\rightarrow E}}{Hyp}}$$

 $\beta :$ 

$$\frac{\frac{\frac{\frac{}{(p \rightarrow r) \wedge (q \rightarrow r), p \vee q, q \vdash (p \rightarrow r) \wedge (q \rightarrow r)}}{Hyp}}{\frac{\frac{}{(p \rightarrow r) \wedge (q \rightarrow r), p \vee q, q \vdash q \rightarrow r}}{\wedge E_2}}}{\frac{\frac{}{(p \rightarrow r) \wedge (q \rightarrow r), p \vee q, q \vdash r}}{\rightarrow E}}{Hyp}}$$

(c)

$$\frac{\frac{\frac{}{p \wedge q \rightarrow r, p, q \vdash p \wedge q \rightarrow r}}{Hyp} \quad \frac{\frac{\frac{}{p \wedge q \rightarrow r, p, q \vdash p}}{Hyp} \quad \frac{\frac{}{p \wedge q \rightarrow r, p, q \vdash q}}{Hyp}}{\wedge I}}{\frac{\frac{\frac{}{p \wedge q \rightarrow r, p, q \vdash p \wedge q}}{\rightarrow E}}{p \wedge q \rightarrow r, p, q \vdash p \wedge q \rightarrow r}}{\rightarrow E}}{\frac{\frac{\frac{}{p \wedge q \rightarrow r, p, q \vdash r}}{\rightarrow I}}{p \wedge q \rightarrow r, p \vdash q \rightarrow r}}{\rightarrow I}}{\frac{\frac{\frac{}{p \wedge q \rightarrow r \vdash p \rightarrow (q \rightarrow r)}}{\rightarrow I}}{\vdash (p \wedge q \rightarrow r) \rightarrow (p \rightarrow (q \rightarrow r))}}{\rightarrow I}}$$

(d)

