## Homework Assignment \#2

## Note

This assignment is due 1:20PM Thursday, October 23, 2014. Please write or type your answers on A4 (or similar size) paper. Put your completed homework on the instructor's desk before the class starts. For late submissions, please drop them in Yih-Kuen Tsay's mail box on the first floor of Management Building II. A late submission will be penalized by $20 \%$ for each working day overdue. You may discuss the problems with others, but copying answers is strictly forbidden.

## Problems

We assume the binding powers of the logical connectives and the entailment symbol decrease in this order: $\neg,\{\forall, \exists\},\{\wedge, \vee\}, \rightarrow, \leftrightarrow, \vdash$.

1. Prove, using Natural Deduction, the validity of the following sequents:
(a) $\forall x(P(x) \rightarrow Q(x)) \vdash \forall x P(x) \rightarrow \forall x Q(x)$
(10 points)
(b) $\vdash \exists x \forall y P(x, y) \rightarrow \forall y \exists x P(x, y)$
(c) $\forall x(A(x) \rightarrow B) \vdash \exists x A(x) \rightarrow B$, assuming $x$ does not occur free in $B$.
2. Prove, using Natural Deduction for the first-order logic with equality ( $=$ ), that $=$ is an equivalence relation between terms, i.e., the following are valid sequents, in addition to the obvious " $\vdash=t$ " (Reflexivity), which follows from the $=$-Introduction rule.
(a) $t_{2}=t_{1} \vdash t_{1}=t_{2}$ (Symmetry)
(10 points)
(b) $t_{1}=t_{2}, t_{2}=t_{3} \vdash t_{1}=t_{3}$ (Transitivity)
3. Taking the preceding valid sequents as axioms, prove using Natural Deduction the following derived rules for equality.
(a) $\frac{\Gamma \vdash t_{2}=t_{1}}{\Gamma \vdash t_{1}=t_{2}}(=$ Symmetry $)$
(10 points)
(b) $\frac{\Gamma \vdash t_{1}=t_{2} \quad \Gamma \vdash t_{2}=t_{3}}{\Gamma \vdash t_{1}=t_{3}}(=$ Transitivity $)$
(10 points)
4. Prove, using Coq, the validity of the following sequents. You may either print out the Coq proofs or email them to the instructor.
(a) $\vdash(p \vee q \rightarrow r) \rightarrow(p \rightarrow r) \wedge(q \rightarrow r)$
(b) $\vdash(p \rightarrow(q \rightarrow r)) \rightarrow(p \wedge q \rightarrow r)$
(c) $\vdash(\neg p \vee q) \rightarrow(p \rightarrow q)$
