## Suggested Solutions for Homework Assignment #2

We assume the binding powers of the logical connectives and the entailment symbol decrease in this order:  $\neg$ ,  $\{\forall, \exists\}, \{\land, \lor\}, \rightarrow, \leftrightarrow, \vdash$ .

- 1. Prove, using *Natural Deduction*, the validity of the following sequents:
  - (a)  $\forall x (P(x) \to Q(x)) \vdash \forall x P(x) \to \forall x Q(x)$ Solution. Assume w does not occur free either in P(x) or in Q(x).

$$\frac{\alpha}{\forall x(P(x) \to Q(x)), \forall xP(x) \vdash \forall xP(x)} (\forall yp)} (\forall x(P(x) \to Q(x)), \forall xP(x) \vdash P(w)} (\forall E) \\ \frac{\forall x(P(x) \to Q(x)), \forall xP(x) \vdash Q(w)}{\forall x(P(x) \to Q(x)), \forall xP(x) \vdash \forall xQ(x)} (\forall I)} \\ \frac{\forall x(P(x) \to Q(x)), \forall xP(x) \vdash \forall xQ(x)}{\forall x(P(x) \to Q(x)) \vdash \forall xP(x) \to \forall xQ(x)} (\to I)}$$

 $\alpha$ :

$$\frac{\overline{\forall x (P(x) \to Q(x)), \forall x P(x) \vdash \forall x (P(x) \to Q(x))}}{\forall x (P(x) \to Q(x)), \forall x P(x) \vdash P(w) \to Q(w)} \overset{(Hyp)}{\forall x (P(x) \to Q(x)), \forall x P(x) \vdash P(w) \to Q(w)}$$

(b)  $\vdash \exists x \forall y P(x, y) \rightarrow \forall y \exists x P(x, y)$ Solution. Assume both w and z do not occur free in P(x, y).

$$\frac{\exists x \forall y P(x,y), \forall y P(z,y) \vdash \forall y P(z,y)}{\exists x \forall y P(x,y), \forall y P(z,y) \vdash P(z,y)} \xrightarrow{(\forall E)} \frac{\exists x \forall y P(x,y), \forall y P(z,y) \vdash P(z,y)}{\exists x \forall y P(x,y), \forall y P(z,y) \vdash \exists x P(x,w)} \xrightarrow{(\exists I)} \frac{\exists x \forall y P(x,y) \vdash \exists x P(x,w)}{\exists x \forall y P(x,y) \vdash \forall y \exists x P(x,y)} \xrightarrow{(\forall I)} \frac{\exists x \forall y P(x,y) \vdash \forall y \exists x P(x,y)}{\vdash \exists x \forall y P(x,y) \rightarrow \forall y \exists x P(x,y)} \xrightarrow{(\to I)}$$

(c)  $\forall x (A(x) \to B) \vdash \exists x A(x) \to B$ , assuming x does not occur free in B. Solution. Assume w does not occur free either in A(x) or in B.

$$\frac{ \overline{\forall x (A(x) \to B), \exists x A(x) \vdash \exists x A(x)}^{(Hyp)} \quad \alpha}{ \overline{\forall x (A(x) \to B), \exists x A(x) \vdash B}^{(\to I)}} \xrightarrow{(\to I)}$$

 $\alpha$ :

$$\frac{ \forall x (A(x) \to B), \exists x A(x), A(w) \vdash \forall x (A(x) \to B)}{\forall x (A(x) \to B), \exists x A(x), A(w) \vdash A(w) \to B} \xrightarrow{(\forall E)} \frac{\beta}{\forall x (A(x) \to B), \exists x A(x), A(w) \vdash B} \xrightarrow{(\to E)}$$

 $\beta$ :

$$\forall x (A(x) \to B), \exists x A(x), A(w) \vdash A(w)$$
 (Hyp)

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- 2. Prove, using *Natural Deduction* for the first-order logic with equality (=), that = is an equivalence relation between terms, i.e., the following are valid sequents, in addition to the obvious " $\vdash t = t$ " (Reflexivity), which follows from the =-Introduction rule.
  - (a)  $t_2 = t_1 \vdash t_1 = t_2$  (Symmetry) Solution.

$$\frac{t_2 = t_1 \vdash t_2 = t_1}{t_2 = t_1 \vdash t_1 = t_2} \xrightarrow{(Hyp)} \frac{t_2 = t_1 \vdash t_2 = t_2}{(= E)}$$

(b)  $t_1 = t_2, t_2 = t_3 \vdash t_1 = t_3$  (Transitivity) Solution.

- 3. Taking the preceding valid sequents as axioms, prove using *Natural Deduction* the following derived rules for equality.
  - (a)  $\frac{\Gamma \vdash t_2 = t_1}{\Gamma \vdash t_1 = t_2} (= Symmetry)$ Solution.

$$\frac{\overline{\Gamma, t_2 = t_1 \vdash t_1 = t_2} \stackrel{(Axiom(Symmetry))}{(\rightarrow I)}}{\Gamma \vdash t_2 = t_1 \rightarrow t_1 = t_2} \stackrel{(\rightarrow I)}{(\rightarrow E)} \Gamma \vdash t_1 = t_2}$$

(b)  $\frac{\Gamma \vdash t_1 = t_2 \qquad \Gamma \vdash t_2 = t_3}{\Gamma \vdash t_1 = t_3} (= Transitivity)$ 

$$\frac{\alpha \qquad \Gamma \vdash t_1 = t_2}{\Gamma \vdash t_2 = t_3 \rightarrow t_1 = t_3} \stackrel{(\to E)}{} \qquad \Gamma \vdash t_2 = t_3 \qquad (\to E)$$

$$\Gamma \vdash t_1 = t_3$$

 $\alpha$ :

$$\frac{\Gamma, t_1 = t_2, t_2 = t_3 \vdash t_1 = t_3}{\Gamma, t_1 = t_2 \vdash t_2 = t_3 \rightarrow t_1 = t_3} \overset{(Axiom(Transitivity))}{(\rightarrow I)}{\Gamma \vdash t_1 = t_2 \rightarrow (t_2 = t_3 \rightarrow t_1 = t_3)} \overset{(\rightarrow I)}{(\rightarrow I)}$$

- 4. Prove, using Coq, the validity of the following sequents. You may either print out the Coq proofs or email them to the instructor.
  - (a)  $\vdash (p \lor q \to r) \to (p \to r) \land (q \to r)$
  - (b)  $\vdash (p \rightarrow (q \rightarrow r)) \rightarrow (p \land q \rightarrow r)$
  - (c)  $\vdash (\neg p \lor q) \to (p \to q)$

Solution. To be completed.