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2. Prove, using *Natural Deduction* for the first-order logic with equality ($=$), that $=$ is an equivalence relation between terms, i.e., the following are valid sequents, in addition to the obvious “ $\vdash t = t$ ” (Reflexivity), which follows from the $=$ -Introduction rule.

- (a) $t_2 = t_1 \vdash t_1 = t_2$ (Symmetry)

Solution.

$$\frac{\frac{}{t_2 = t_1 \vdash t_2 = t_1} \text{ (Hyp)} \quad \frac{}{t_2 = t_1 \vdash t_2 = t_2} \text{ (= I)}}{t_2 = t_1 \vdash t_1 = t_2} \text{ (= E)}$$

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- (b) $t_1 = t_2, t_2 = t_3 \vdash t_1 = t_3$ (Transitivity)

Solution.

$$\frac{\frac{}{t_1 = t_2, t_2 = t_3 \vdash t_2 = t_3} \text{ (Hyp)} \quad \frac{}{t_1 = t_2, t_2 = t_3 \vdash t_1 = t_2} \text{ (Hyp)}}{t_1 = t_2, t_2 = t_3 \vdash t_1 = t_3} \text{ (= E)}$$

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3. Taking the preceding valid sequents as axioms, prove using *Natural Deduction* the following derived rules for equality.

- (a) $\frac{\Gamma \vdash t_2 = t_1}{\Gamma \vdash t_1 = t_2} \text{ (= Symmetry)}$

Solution.

$$\frac{\frac{\frac{}{\Gamma, t_2 = t_1 \vdash t_1 = t_2} \text{ (Axiom(Symmetry))}}{\Gamma \vdash t_2 = t_1 \rightarrow t_1 = t_2} \text{ (}\rightarrow\text{I)}}{\Gamma \vdash t_1 = t_2} \text{ (}\rightarrow\text{E)}$$

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- (b) $\frac{\Gamma \vdash t_1 = t_2 \quad \Gamma \vdash t_2 = t_3}{\Gamma \vdash t_1 = t_3} \text{ (= Transitivity)}$

Solution.

$$\frac{\frac{\alpha \quad \Gamma \vdash t_1 = t_2}{\Gamma \vdash t_2 = t_3 \rightarrow t_1 = t_3} \text{ (}\rightarrow\text{E)} \quad \Gamma \vdash t_2 = t_3}{\Gamma \vdash t_1 = t_3} \text{ (}\rightarrow\text{E)}$$

$\alpha :$

$$\frac{\frac{\frac{}{\Gamma, t_1 = t_2, t_2 = t_3 \vdash t_1 = t_3} \text{ (Axiom(Transitivity))}}{\Gamma, t_1 = t_2 \vdash t_2 = t_3 \rightarrow t_1 = t_3} \text{ (}\rightarrow\text{I)}}{\Gamma \vdash t_1 = t_2 \rightarrow (t_2 = t_3 \rightarrow t_1 = t_3)} \text{ (}\rightarrow\text{I)}$$

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4. Prove, using Coq, the validity of the following sequents. You may either print out the Coq proofs or email them to the instructor.

- (a) $\vdash (p \vee q \rightarrow r) \rightarrow (p \rightarrow r) \wedge (q \rightarrow r)$

- (b) $\vdash (p \rightarrow (q \rightarrow r)) \rightarrow (p \wedge q \rightarrow r)$

- (c) $\vdash (\neg p \vee q) \rightarrow (p \rightarrow q)$

Solution. To be completed.

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