

## Suggested Solutions for Homework Assignment #2

We assume the binding powers of the logical connectives and the entailment symbol decrease in this order:  $\neg$ ,  $\{\forall, \exists\}$ ,  $\{\wedge, \vee\}$ ,  $\rightarrow, \leftrightarrow, \vdash$ .

1. (30 points) Prove, using *Natural Deduction*, the validity of the following sequents:

$$(a) \forall x(P(x) \rightarrow Q(x)) \vdash \forall xP(x) \rightarrow \forall xQ(x)$$

*Solution.* Assume  $w$  does not occur free either in  $P(x)$  or in  $Q(x)$ .

$$\alpha : \frac{\frac{\frac{\frac{\forall x(P(x) \rightarrow Q(x)), \forall xP(x) \vdash \forall xP(x)}{(\forall E)}}{\frac{\forall x(P(x) \rightarrow Q(x)), \forall xP(x) \vdash P(w)}{(\rightarrow E)}}}{\frac{\forall x(P(x) \rightarrow Q(x)), \forall xP(x) \vdash Q(w)}{(\forall I)}}}{\frac{\forall x(P(x) \rightarrow Q(x)), \forall xP(x) \vdash \forall xQ(x)}{(\rightarrow I)}}$$

$\alpha :$

$$\frac{\forall x(P(x) \rightarrow Q(x)), \forall xP(x) \vdash \forall x(P(x) \rightarrow Q(x))}{\forall x(P(x) \rightarrow Q(x)), \forall xP(x) \vdash P(w) \rightarrow Q(w)}$$

□

$$(b) \vdash \exists x \forall y P(x, y) \rightarrow \forall y \exists x P(x, y)$$

*Solution.* Assume both  $w$  and  $z$  do not occur free in  $P(x, y)$ .

$$\frac{\frac{\frac{\exists x \forall y P(x, y) \vdash \exists x \forall y P(x, y)}{(\exists E)}}{\frac{\frac{\frac{\exists x \forall y P(x, y), \forall y P(z, y) \vdash \forall y P(z, y)}{(\forall E)}}{\frac{\exists x \forall y P(x, y), \forall y P(z, y) \vdash P(z, w)}{(\exists I)}}}{\frac{\exists x \forall y P(x, y), \forall y P(z, y) \vdash \exists x P(x, w)}{(\exists E)}}}}{\frac{\exists x \forall y P(x, y) \vdash \exists x P(x, w)}{(\forall I)}}}{\vdash \exists x \forall y P(x, y) \rightarrow \forall y \exists x P(x, y)}$$

□

$$(c) \forall x(A(x) \rightarrow B) \vdash \exists x A(x) \rightarrow B, \text{ assuming } x \text{ does not occur free in } B.$$

*Solution.* Assume  $w$  does not occur free either in  $A(x)$  or in  $B$ .

$$\alpha : \frac{\frac{\frac{\forall x(A(x) \rightarrow B), \exists x A(x) \vdash \exists x A(x)}{(\exists E)}}{\frac{\forall x(A(x) \rightarrow B), \exists x A(x) \vdash B}{(\rightarrow I)}}}{\forall x(A(x) \rightarrow B) \vdash \exists x A(x) \rightarrow B}$$

$\alpha :$

$$\beta : \frac{\frac{\frac{\forall x(A(x) \rightarrow B), \exists x A(x), A(w) \vdash \forall x(A(x) \rightarrow B)}{(\forall E)}}{\frac{\forall x(A(x) \rightarrow B), \exists x A(x), A(w) \vdash A(w) \rightarrow B}{(\rightarrow E)}}}{\forall x(A(x) \rightarrow B), \exists x A(x), A(w) \vdash B}$$

$\beta :$

$$\frac{\forall x(A(x) \rightarrow B), \exists x A(x), A(w) \vdash A(w)}{(\forall E)}$$

□

2. Prove, using *Natural Deduction* for the first-order logic with equality ( $=$ ), that  $=$  is an equivalence relation between terms, i.e., the following are valid sequents, in addition to the obvious “ $\vdash t = t$ ” (Reflexivity), which follows from the  $=$ -Introduction rule.

(a)  $t_2 = t_1 \vdash t_1 = t_2$  (Symmetry)

*Solution.*

$$\frac{\frac{t_2 = t_1 \vdash t_2 = t_1}{(Hyp)} \quad \frac{t_2 = t_1 \vdash t_2 = t_2}{(=I)}}{t_2 = t_1 \vdash t_1 = t_2} (=E)$$

□

(b)  $t_1 = t_2, t_2 = t_3 \vdash t_1 = t_3$  (Transitivity)

*Solution.*

$$\frac{\frac{t_1 = t_2, t_2 = t_3 \vdash t_2 = t_3}{(Hyp)} \quad \frac{t_1 = t_2, t_2 = t_3 \vdash t_1 = t_2}{(=E)}}{t_1 = t_2, t_2 = t_3 \vdash t_1 = t_3}$$

□

3. (20 points) Taking the preceding valid sequents as axioms, prove using *Natural Deduction* the following derived rules for equality.

(a)  $\frac{\Gamma \vdash t_2 = t_1}{\Gamma \vdash t_1 = t_2} (= Symmetry)$

*Solution.*

$$\frac{\frac{\frac{\Gamma, t_2 = t_1 \vdash t_1 = t_2}{(Axiom(Symmetry))}}{\Gamma \vdash t_2 = t_1 \rightarrow t_1 = t_2} \quad \frac{\Gamma \vdash t_2 = t_1}{(\rightarrow E)}}{\Gamma \vdash t_1 = t_2}$$

□

(b)  $\frac{\Gamma \vdash t_1 = t_2 \quad \Gamma \vdash t_2 = t_3}{\Gamma \vdash t_1 = t_3} (= Transitivity)$

*Solution.*

$$\frac{\alpha \quad \frac{\Gamma \vdash t_1 = t_2}{(\rightarrow E)} \quad \Gamma \vdash t_2 = t_3}{\Gamma \vdash t_1 = t_3} \quad \Gamma \vdash t_2 = t_3 \quad (\rightarrow E)$$

$\alpha :$

$$\frac{\frac{\frac{\Gamma, t_1 = t_2, t_2 = t_3 \vdash t_1 = t_3}{(Axiom(Transitivity))}}{\frac{\Gamma, t_1 = t_2 \vdash t_2 = t_3 \rightarrow t_1 = t_3}{(\rightarrow I)}} \quad \frac{\Gamma \vdash t_2 = t_3}{(\rightarrow E)}}{\Gamma \vdash t_1 = t_2 \rightarrow (t_2 = t_3 \rightarrow t_1 = t_3)}$$

□

4. (30 points) A first-order theory for *groups* contains the following three axioms:

- $\forall a \forall b \forall c (a \cdot (b \cdot c) = (a \cdot b) \cdot c)$ . (Associativity)
- $\forall a ((a \cdot e = a) \wedge (e \cdot a = a))$ . (Identity)
- $\forall a ((a \cdot a^{-1} = e) \wedge (a^{-1} \cdot a = e))$ . (Inverse)

Here  $\cdot$  is the binary operation,  $e$  is a constant, called the identity, and  $(\cdot)^{-1}$  is the inverse function which gives the inverse of an element. Let  $M$  denote the set of the three axioms. Prove, using *Natural Deduction* plus the derived rules in the preceding problem, the validity of the following sequents:

- (a)  $M \vdash \forall a \forall b \forall c((a \cdot b = a \cdot c) \rightarrow b = c)$ . (Hint: a typical proof in algebra books is the following:  $b = e \cdot b = (a^{-1} \cdot a) \cdot b = a^{-1} \cdot (a \cdot b) = a^{-1} \cdot (a \cdot c) = (a^{-1} \cdot a) \cdot c = e \cdot c = c$ .)

*Solution.*

$$\frac{\alpha \quad \delta}{M, x \cdot y = x \cdot z \vdash y = z} (= E) \\ \frac{}{M \vdash (x \cdot y = x \cdot z) \rightarrow y = z} (\rightarrow I) \\ \frac{}{M \vdash \forall c((x \cdot y = x \cdot c) \rightarrow y = c)} (\forall I) \\ \frac{}{M \vdash \forall b \forall c((x \cdot b = x \cdot c) \rightarrow b = c)} (\forall I) \\ \frac{}{M \vdash \forall a \forall b \forall c((a \cdot b = a \cdot c) \rightarrow b = c)} (\forall I)$$

$\alpha :$

$$\frac{\beta \quad \gamma}{M, x \cdot y = x \cdot z \vdash (x^{-1} \cdot x) \cdot y = y} (= E) \quad \frac{\overline{M, x \cdot y = x \cdot z \vdash \forall a \forall b \forall c(a \cdot (b \cdot c) = (a \cdot b) \cdot c)} (Hyp) \quad M, x \cdot y = x \cdot z \vdash \forall b \forall c(x^{-1} \cdot (b \cdot c) = (x^{-1} \cdot b) \cdot c)} (\forall E) \\ \frac{}{M, x \cdot y = x \cdot z \vdash \forall c(x^{-1} \cdot (x \cdot c) = (x^{-1} \cdot x) \cdot c)} (\forall E) \\ \frac{}{M, x \cdot y = x \cdot z \vdash x^{-1} \cdot (x \cdot y) = (x^{-1} \cdot x) \cdot y} (\forall E) \\ M, x \cdot y = x \cdot z \vdash x^{-1} \cdot (x \cdot y) = y$$

$\beta :$

$$\frac{}{M, x \cdot y = x \cdot z \vdash \forall a(a \cdot a^{-1} = e \wedge a^{-1} \cdot a = e)} (Hyp) \\ \frac{}{M, x \cdot y = x \cdot z \vdash x \cdot x^{-1} = e \wedge x^{-1} \cdot x = e} (\forall E) \\ \frac{}{M, x \cdot y = x \cdot z \vdash x^{-1} \cdot x = e} (\wedge E_2) \\ \frac{}{M, x \cdot y = x \cdot z \vdash e = x^{-1} \cdot x} (= Symmetry)$$

$\gamma :$

$$\frac{}{M, x \cdot y = x \cdot z \vdash \forall a(a \cdot e = a \wedge e \cdot a = a)} (Hyp) \\ \frac{}{M, x \cdot y = x \cdot z \vdash y \cdot e = y \wedge e \cdot y = y} (\forall E) \\ \frac{}{M, x \cdot y = x \cdot z \vdash e \cdot y = y} (\wedge E_2)$$

$\delta :$

$$\frac{\overline{M, x \cdot y = x \cdot z \vdash x \cdot y = x \cdot z} (Hyp) \quad M, x \cdot y = x \cdot z \vdash x \cdot z = x \cdot y (= Symmetry) \quad \text{the proof tree is similar to } \alpha}{\overline{M, x \cdot y = x \cdot z \vdash x^{-1} \cdot (x \cdot z) = z} (= E)} \\ M, x \cdot y = x \cdot z \vdash x^{-1} \cdot (x \cdot y) = z$$

□

- (b)  $M \vdash \forall a \forall b \forall c(((a \cdot b = e) \wedge (b \cdot a = e) \wedge (a \cdot c = e) \wedge (c \cdot a = e)) \rightarrow b = c)$ , which says that every element has a unique inverse. (Hint: a typical proof in algebra books is the following:  $b = b \cdot e = b \cdot (a \cdot c) = (b \cdot a) \cdot c = e \cdot c = c$ .)

*Solution.* We use  $N$  to denote  $x \cdot y = e \wedge y \cdot x = e \wedge x \cdot z = e \wedge z \cdot x = e$ .

$$\begin{array}{c}
\frac{(1)\alpha \quad (1)\delta}{M, N, x \cdot y = x \cdot z \vdash y = z} (=E) \\
\frac{}{M, N \vdash x \cdot y = x \cdot z \rightarrow y = z} (\rightarrow I) \quad \frac{\alpha \quad \beta}{M, N \vdash x \cdot y = x \cdot z} (=E) \\
\frac{}{M, N \vdash y = z} (\rightarrow I) \\
\frac{}{M \vdash (x \cdot y = e \wedge y \cdot x = e \wedge x \cdot z = e \wedge z \cdot x = e) \rightarrow y = z} (\rightarrow I) \\
\frac{}{M \vdash \forall c((x \cdot y = e \wedge y \cdot x = e \wedge x \cdot c = e \wedge c \cdot x = e) \rightarrow y = c)} (\forall I) \\
\frac{}{M \vdash \forall b \forall c((x \cdot b = e \wedge b \cdot x = e \wedge x \cdot c = e \wedge c \cdot x = e) \rightarrow b = c)} (\forall I) \\
\frac{}{M \vdash \forall a \forall b \forall c((a \cdot b = e \wedge b \cdot a = e \wedge a \cdot c = e \wedge c \cdot a = e) \rightarrow b = c)} (\forall I)
\end{array}$$

$\alpha :$

$$\begin{array}{c}
\frac{}{M, N \vdash x \cdot y = e \wedge y \cdot x = e \wedge x \cdot z = e \wedge z \cdot x = e} (Hyp) \\
\frac{}{M, N \vdash x \cdot z = e \wedge z \cdot x = e} (\wedge E_2) \\
\frac{}{M, N \vdash x \cdot z = e} (\wedge E_1) \\
\frac{}{M, N \vdash e = x \cdot z} (=Symmetry)
\end{array}$$

$\beta :$

$$\frac{}{M, N \vdash x \cdot y = e \wedge y \cdot x = e \wedge x \cdot z = e \wedge z \cdot x = e} (Hyp) \\
\frac{}{M, N \vdash x \cdot y = e} (\wedge E_1)$$

□